# Abstract geometrical computation for Black hole computation 

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## Outline

1. Black hole computation
2. Cellular automata to Abstract geometrical computation
3. Signal machine and restriction
4. Turing-computing power
5. Black hole effect
6. Conclusion and extension

## Black hole computation

## Black hole model



1. Observer at the "edge"

## Black hole model



1. Observer at the "edge"
2. Machine sent into the black hole infinitely accelerated

## Black hole model



1. Observer at the "edge"
2. Machine sent into the black hole infinitely accelerated
3. Message sent by the machine received by the observer within a bounded delay

## Malament-Hogarth space-time

## Observer's life-line



Message indicates the result of the computation
After the delay, the observer knows whether the computation stops
Any recursively enumerable problem can be decided!

## Related models

Main idea: infinitely many "iterations" on a sub-time-scale

Can be achieved with a transfinite ordinal scale as in:
Infinite time Turing machines
[Hamkins 02]

Or with a "Zeno" sub-scale as in:
Piecewise constant derivative systems
[Asarin \& Maler 95, Bournez 99]

We use the last approach

## Cellular automata to

Abstract geometrical computation

## Basis

Well-known Model for parallelism, biology, physics...

Discrete time and space


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Locally finitely many states


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Discrete time and space
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Local interaction


## Basis

Well-known Model for parallelism, biology, physics...

Discrete time and space
Locally finitely many states
Local interaction
Uniform in space and time

Turing-universal model


## Space-time diagrams understanding


(a)

(b)


FIG. 7. The four different (out of 14 possible) interaction products for the $\alpha+\beta$ interaction

FIG. 7. Rule 54. (a) Annibilation of the radiating partie. (b) The same as (a) with the mapping defined in Fig. 6.

## [Boccara, Nasser \& Roger 91]

Observation of discrete lines $\rightsquigarrow$ keys to dynamic

## Space-time diagrams designing


[Fischer 65]

[Varshavsky et al. 70]

[Varshavsky et al. 70]

Easily generated discrete lines $\rightsquigarrow$ special purposes CA design

## Continuous abstraction

Signal: important notion, often used in literature

- to describe
- to design


Space $(\mathbb{Z})$


## Signal machine and restriction

## Model definition


$\mathbb{R} \times \mathbb{R}^{+}$


Space

## Model definition


$\mathbb{R} \times \mathbb{R}^{+}$
Signal
Position
(Meta-signal, position)
Meta-signal
$\mu=(\iota, \nu)$


Space

## Model definition

## $\mathbb{R} \times \mathbb{R}^{+}$

Signal
Position
(Meta-signal, position)


Space

Rule

$$
\left\{\mu_{i}^{-}\right\}_{i} \rightarrow\left\{\mu_{j}^{-}\right\}_{j}
$$



1

## Properties and examples

- Finite number of values \& rules
- Light cone
- Local interaction
- Uniform in space and time
- Continuous space and time



## Strange space-time diagrams



Zeno artifact


Unwanted cases

Unwanted because

- The number of signals is bursting to infinity
(free creation of mater/energy)
- Difficulty (if not impossibility) to define continuation there


## Restriction

- Energy : $\mu \longrightarrow E(\mu) \in \mathbb{N}^{*}$
- $\forall \rho=\left\{\mu_{i}^{-}\right\}_{i} \rightarrow\left\{\mu_{j}^{+}\right\}_{j}, \quad \sum E\left(\mu_{i}^{-}\right) \geq \sum E\left(\mu_{i}^{+}\right)$
- $E($ configuration $)=\sum E($ existing signals $)$
- Total energy quantified and bounded
- The total number of signals is bounded

All energies equal
$\rightsquigarrow$ the number of signals is preserved by a collision

## Turing-computing power

## Simulating 2-counter automata

2 non-negative counters $\times 3$ operations
Encoding positions of counters


Encoding of configurations

|  | $a=0$ | $0<a$ |
| :---: | :---: | :---: |
| $b=0$ |  | \% |
| $0<b$ | ot ouba | ¢ |

## Implementation of A--



## Some Examples

be1:A--
IF A!=0 be
pa: B--
A++

$$
\mathrm{IF} \mathrm{~B}!=0 \mathrm{pa}
$$

$$
\text { IF } A!=0 \text { be }
$$

im: B--

$$
A++
$$

A++

$$
\text { IF B }!=0 \quad \mathrm{im} 1
$$

$$
\text { IF A }!=0 \text { be }
$$

im1: B--
A++
A++
A++

$$
\text { IF } B!=0 \quad \mathrm{im} 1
$$

$$
\text { IF } A!=0 \text { be }
$$

$$
\begin{aligned}
& \text { be: B++ } \\
& \text { A-- } \\
& \text { IF } A!=0 \text { be1 } \\
& \text { IF } B!=0 \text { im }
\end{aligned}
$$


$a=1 \quad b=0$

$a=3 \quad b=0$

$a=5 \quad b=0$

## Handling the halt

Restriction is always satisfied but. . .
what about halting?

The instruction turns into a yes/no signal leaving on the left

## Black hole effect

## Providing a strain



## Providing a shrinking

Two consecutive strains with the same directions coefficient $1 / 2$ on one direction then the other


Iterating possible if spatially bounded

## Iterating shrinking


(For a spatially bounded computation)

## Bounding delay

Simulation \& iterated shrinking construction satisfy the restriction


Bounding signals indicate when it is too late to get any answer

## Conclusion and extension

## Conclusion

- Turing computation power in a continuous space and time model
- Geometric model where geometric constructions allow Zeno effects
- Similarity with the Black hole model
- Rational numbers are enough to get all this (i.e. distinction lies in continuity and not in cardinality)


## Extension

- Second (and higher) accumulation could be generated by lifting the restriction (hierarchy climbing)
- Real values
- could be used as oracles
- analog model

