

Intrinsic Simulation between Cellular Automata

N. Ollinger

SEMINARIO MATEMÁTICAS DISCRETAS

CMM, Santiago, June 2001

(revised and commented version)

Cellular Automata

Definition. A d -CA \mathcal{A} is a 4-uple $(\mathbb{Z}^d, S, N, \delta)$ with:

- S the finite set of states of \mathcal{A} ,
- $N \subseteq \mathbb{Z}^d$ finite, the neighborhood of \mathcal{A} ,
- $\delta : S^{|N|} \rightarrow S$ the local rule of \mathcal{A} .

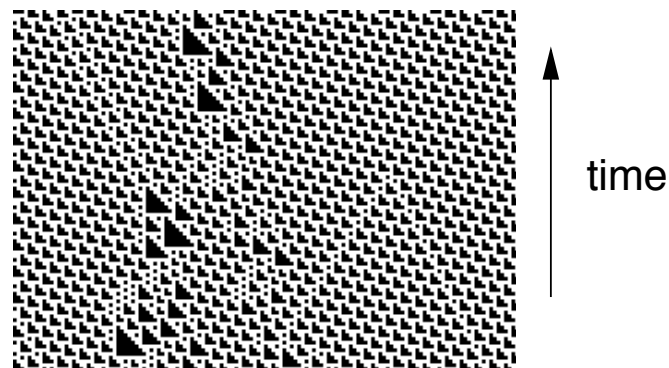
► A *configuration* C is a mapping from \mathbb{Z}^d to S .

Cellular Automata (2)

- ▶ The *global rule* applies δ uniformly according to N :

$$G_{\mathcal{A}}(C)_p = \delta (C_{p+N_1}, \dots, C_{p+N_n}).$$

- ▶ A *space-time diagram* is a graphical representation of an orbit.



Topological Characterization

- ▶ We endow S with the trivial topology.
- ▶ We endow $S^{\mathbb{Z}^d}$ with the induced product topology.
- ▶ The *shift* $\sigma_p : S^{\mathbb{Z}^d} \rightarrow S^{\mathbb{Z}^d}$ is defined as $\sigma_p(C)_i = C_{i+p}$.

Theorem[Hedlund 69]. A map $G : S^{\mathbb{Z}^d} \rightarrow S^{\mathbb{Z}^d}$ is the global rule of a d -CA if and only if it is continuous and commutes with the shifts.

Simulation

- ▶ Simulation = analysis of the computational power.
- ▶ *Extrinsic simulation*: the CA simulates another device.



Turing machine

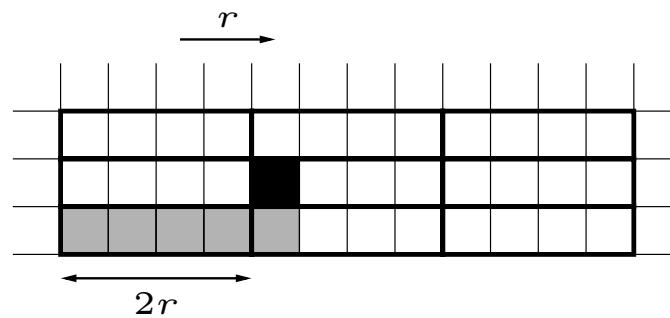


1-CA

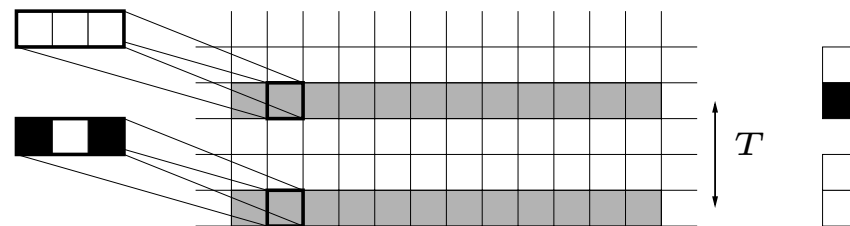
- ▶ *Intrinsic simulation*: the CA simulates another CA.

Classical Intrinsic Simulations

- ▶ Any CA can be simulated by an OCA (*Cole 69, Ibarra 85*)

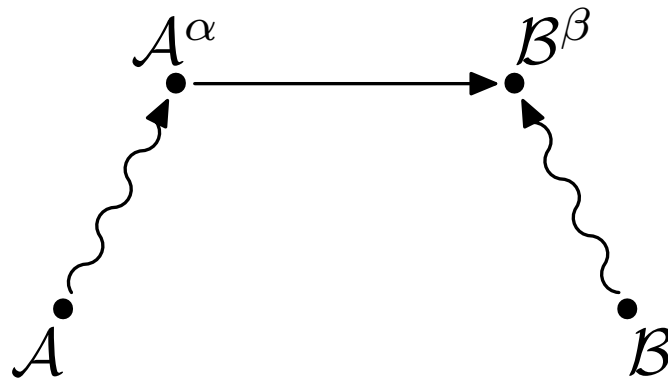


- ▶ Any nilpotent CA can be simulated by the trivial nilpotent CA



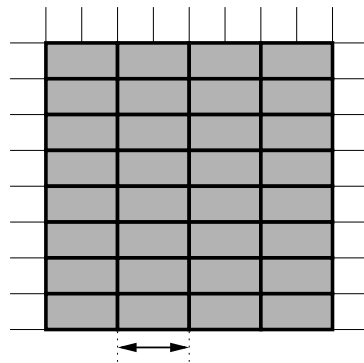
Simulation by Geometrical Transformation

Idea. A CA \mathcal{A} simulates another CA \mathcal{B} if, up to geometrical transformations, any space-time diagram from \mathcal{B} is a space-time diagram from \mathcal{A} .



5 good transformations

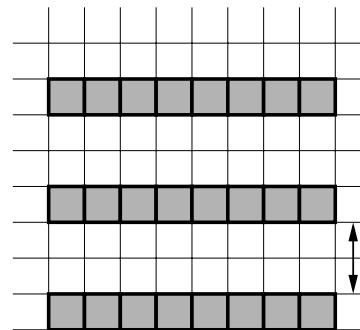
packing



$$o^T \circ \mathcal{A} \circ o^{-T}$$

spatial organization

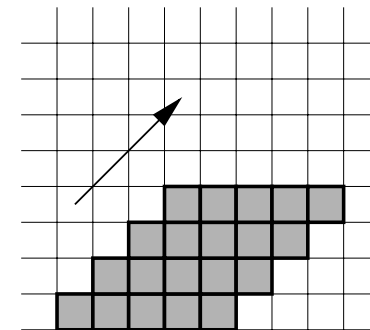
cutting



$$\mathcal{A}^n$$

temporal organization

shifting

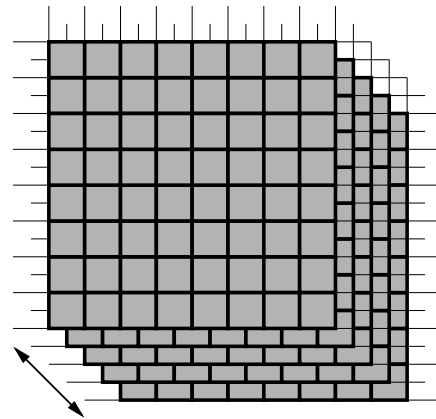


$$\mathcal{A} \circ \sigma_k$$

information mixing

5 good transformations (2)

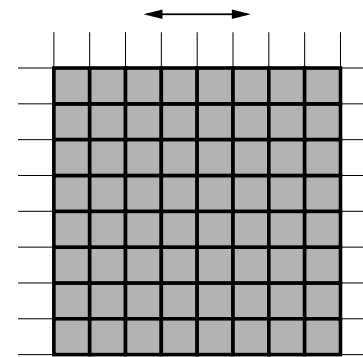
twisting



$$\prod_i \mathcal{A}_i$$

independent layers

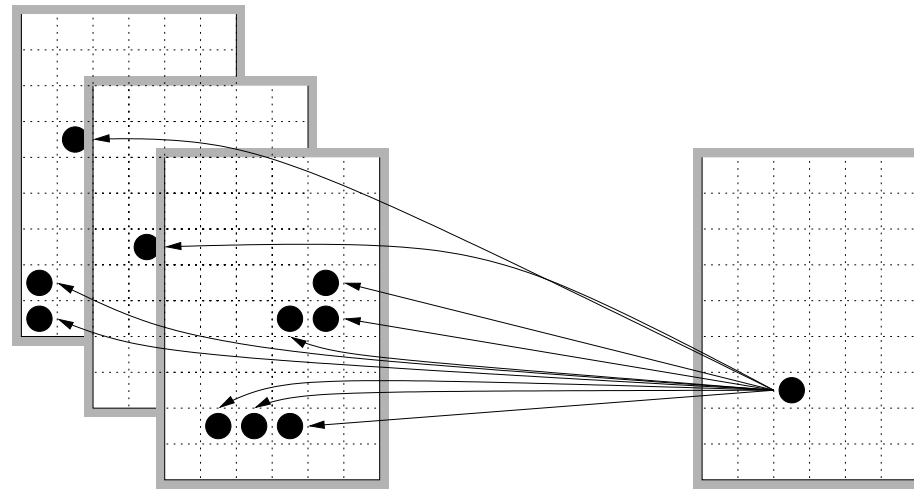
mirroring



$$\leftrightarrow \circ \mathcal{A} \circ \leftrightarrow$$

symmetry

Generalizing Geometrical Transformations



$$\varphi : \mathbb{N} \times \mathbb{Z} \rightarrow 2^{\{1, \dots, k\}} \times \mathbb{N} \times \mathbb{Z}$$

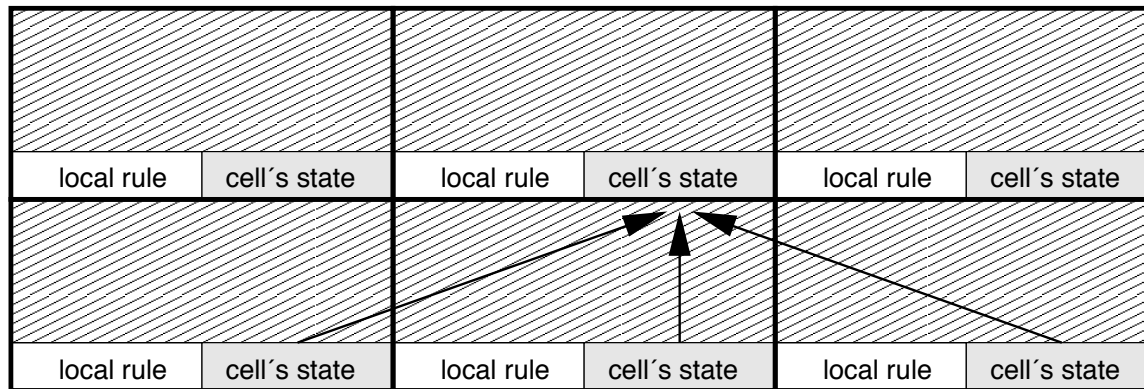
- The new CA must be *completely* defined for any initial CA.

Theorem. There exist no geometrical transformation but compositions of the 5 good previous ones.

Universality

Definition. A *universal CA* is a CA which can simulate any CA.

- ▶ There are universal CA for the P^-C simulation.



- ▶ Therefore, there are universal CA for the PCSTM simulation.

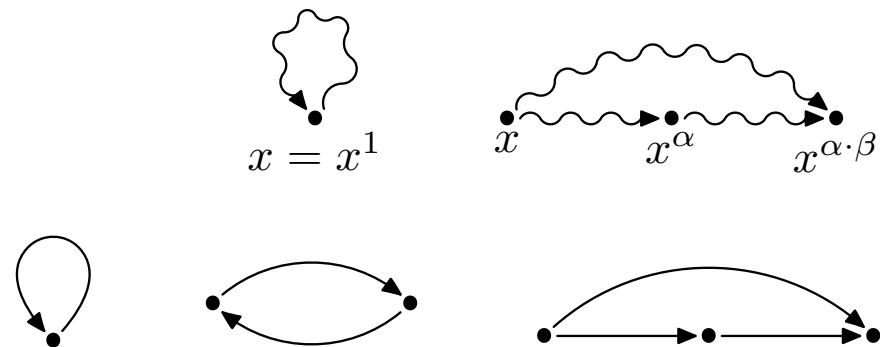
Theorem. The universal CA for PCSTM and P^-CS coincide.

Abstract Bulking

Definition. An *abstract bulking* is a two-sorted first-order structure

$$\mathfrak{A} = (\text{Obj}, \text{Pow}; f : \text{Obj} \times \text{Pow} \rightarrow \text{Obj}, \\ R \subseteq \text{Obj} \times \text{Obj}, \\ \cdot : \text{Pow} \times \text{Pow} \rightarrow \text{Pow})$$

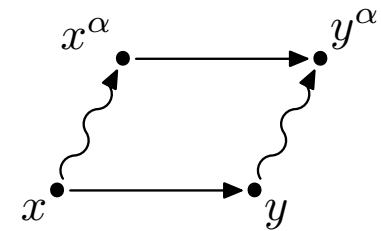
- ▶ (Pow, \cdot) is a monoid,
- ▶ f is compatible with \cdot ,
- ▶ R is a partial order.



More Axioms

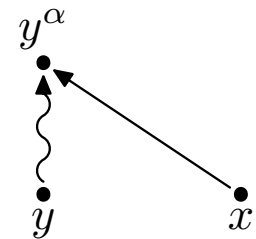
- R is compatible with f .

$$\forall x, y, \alpha, R(x, y) \Rightarrow R(f(x, \alpha), f(y, \alpha))$$



- f preserves richness.

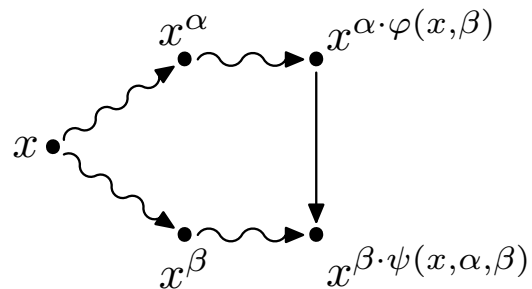
$$\forall x, \alpha, \exists y, R(x, f(y, \alpha))$$



The last axiom

► f keeps objects nearby.

$$\exists \varphi, \psi, \forall x, \alpha, \beta, R(f(x, \alpha \cdot \varphi(x, \beta)), f(x, \beta \cdot \psi(x, \alpha, \beta)))$$



The way to conduct proofs

- ▶ Let Φ denotes the set of axioms.
- ▶ An object y simulates an object x if $\phi(x, y)$ is satisfied.

$$\phi(x, y) = \exists\alpha\exists\beta, R(f(x, \alpha), f(y, \beta))$$

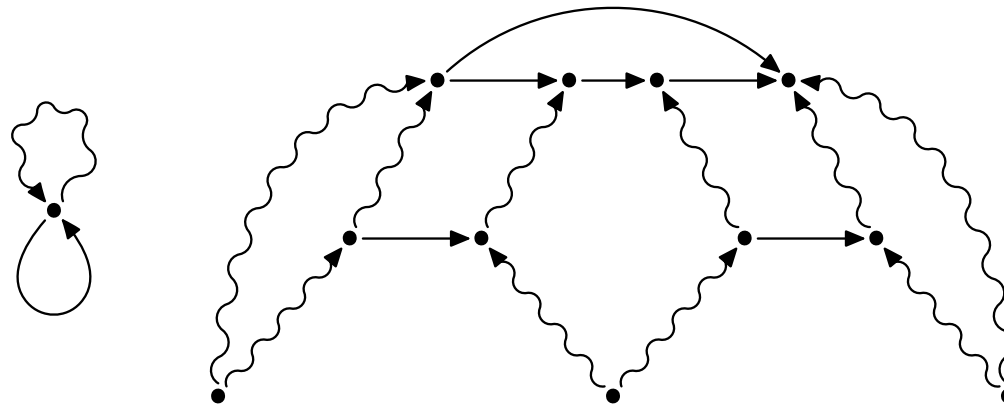
Definition. A *bulking property* is a property on ϕ expressed by a formula φ such that $\Phi \models \varphi$.

The quasi-order

Theorem. “ ϕ is a quasi-order” is a bulking property.

$$\varphi = \forall x, \phi(x, x) \wedge \forall x, y, z, (\phi(x, y) \wedge \phi(y, z)) \Rightarrow \phi(x, z)$$

Proof.



Universality

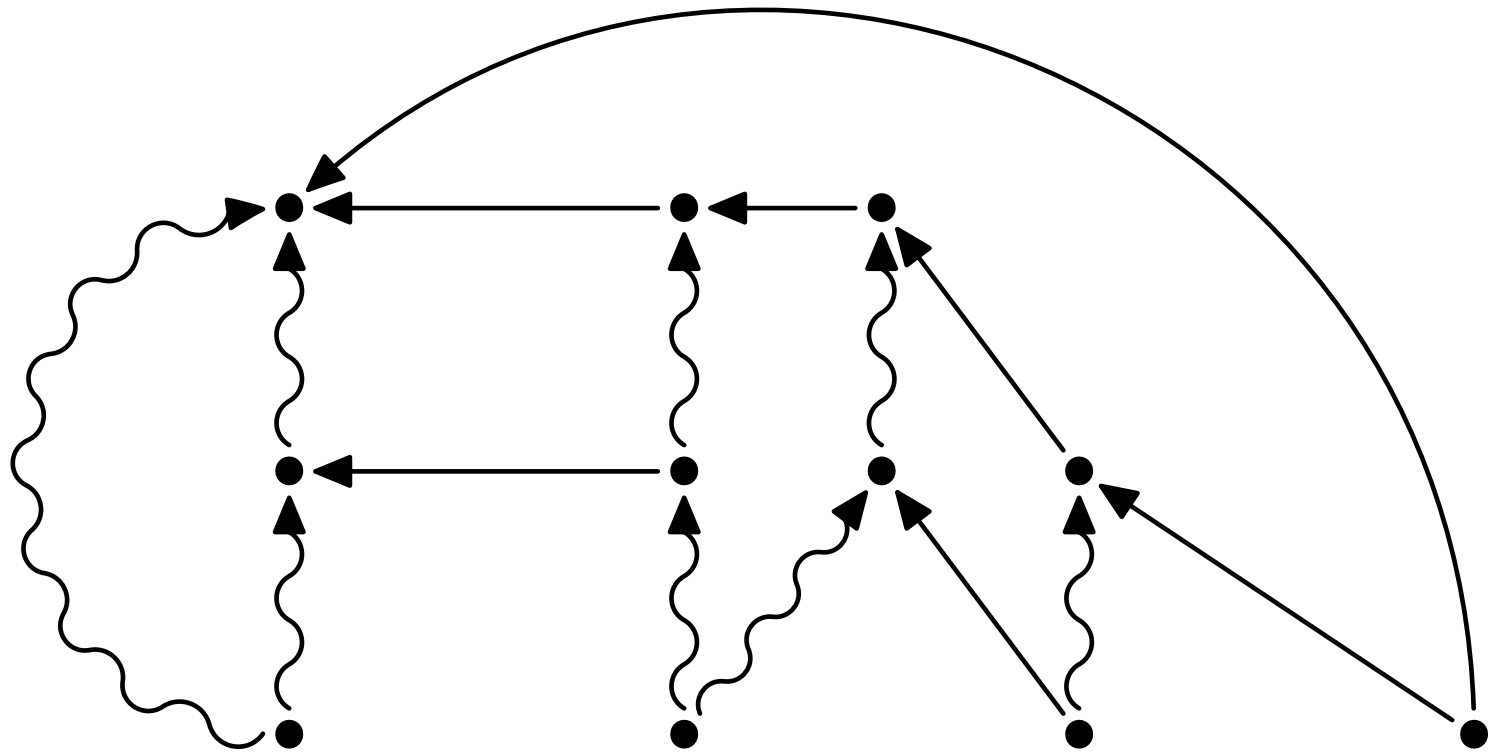
Definition. An object x is *strongly universal* if it can simulate any other object directly (ie without transformation).

$$\psi(x) = \forall y, \exists \alpha, R(y, f(x, \alpha))$$

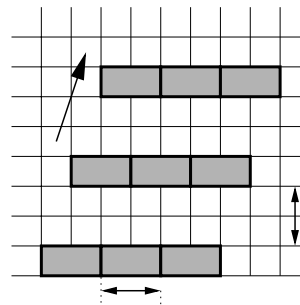
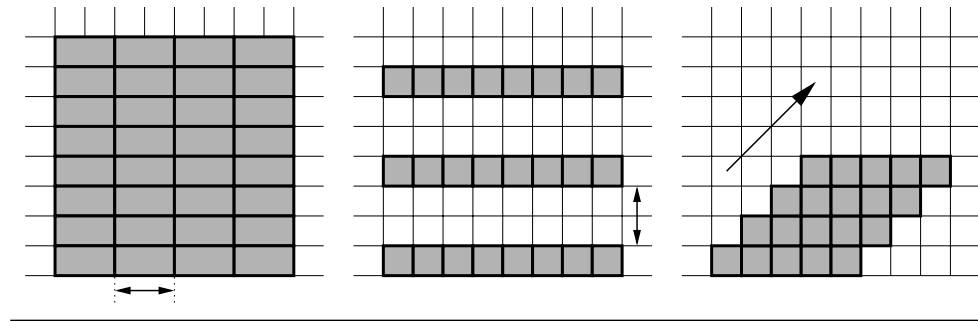
Theorem. “If there exists a strongly universal object, then any universal object is strongly universal” is a bulking property.

$$\varphi = (\exists x, \psi(x)) \Rightarrow \forall x, ((\forall y, \phi(y, x)) \Rightarrow \psi(x))$$

Proof



P⁻CS



$$O^m \circ \mathcal{A}^n \circ \sigma^k \circ O^{-m}$$

P⁻CS and Bulking

Definition. The $\langle m, n, k \rangle$ regular P⁻CS transformation of a CA \mathcal{A} is the CA $\mathcal{A}^{\langle m, n, k \rangle}$ where

$$\mathcal{A}^{\langle m, n, k \rangle} = o^m \circ \mathcal{A}^{mn} \circ o^{-m} \circ \sigma^k.$$

Lemma. \mathcal{A} simulates \mathcal{B} if and only if \mathcal{A} simulates \mathcal{B} regularly.

Theorem. Regular P⁻CS induces a Bulking.

Proof

- ▶ Obj is the set of CA and $\text{Pow} = \mathbb{N}^* \times \mathbb{N}^* \times \mathbb{Z}^*$.
- ▶ $f(\mathcal{A}, \langle m, n, k \rangle) = \mathcal{A}^{\langle m, n, k \rangle}$
- ▶ $R(\mathcal{A}, \mathcal{B})$ if there is an injection φ from $S_{\mathcal{A}}$ into $S_{\mathcal{B}}$ such that

$$\bar{\varphi} \circ \mathcal{A} = \mathcal{B} \circ \bar{\varphi}$$

- ▶ $\langle m, n, k \rangle \cdot \langle m', n', k' \rangle = \langle mm', nn', k' + n'k \rangle$
- ▶ Then we need to check the axioms...

Intrinsic Universality

- ▶ There exists a strongly universal CA on P^-C .
- ▶ Universal CA coincide with universal CA on P^-CS .

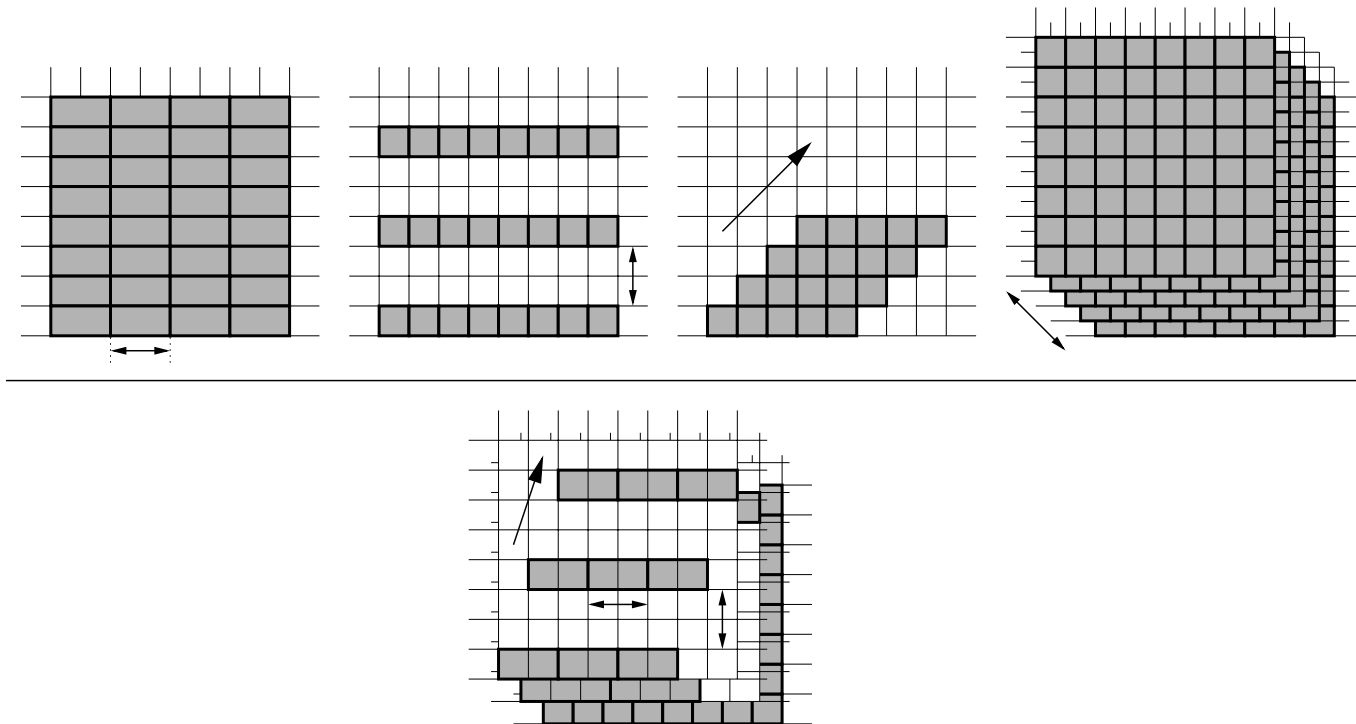
Robust Definition. An intrinsically universal CA is a strongly universal CA on P^-CS .

Theorem[Rapaport 98]. There is no real-time universal CA.

Some results on P^-CS

- ▶ Undecidability anywhere (thanks to Nilpotent CA),
- ▶ Infinitely many equivalence classes,
- ▶ Classical properties are compatible with P^-CS :
trivial CA, two-state CA, first-neighbors CA,
Number-conserving CA, Totalistic CA, Reversible CA.
- ▶ **But** the preorder induces no semi-lattice.

PCST



$$\prod_{i=1}^j \sigma^{m_i} \circ \mathcal{A}^{n_i} \circ \sigma^{k_i} \circ \sigma^{-m_i}$$

P⁻CST and Bulking

Definition. The $\prod \langle m_i, n_i, k_i \rangle$ regular P⁻CST transformation of a CA \mathcal{A} is the CA $\mathcal{A}^{\prod \langle m_i, n_i, k_i \rangle}$ where

$$\mathcal{A}^{\prod \langle m_i, n_i, k_i \rangle} = \prod_i o^{m_i} \circ \mathcal{A}^{m_i n_i} \circ o^{-m_i} \circ \sigma^{k_i}.$$

Lemma. \mathcal{A} simulates \mathcal{B} if and only if \mathcal{A} simulates \mathcal{B} regularly.

Theorem. Regular P⁻CST induces a Bulking.

Proof

- ▶ Obj is the set of CA and $\text{Pow} = \cup_j (\mathbb{N}^* \times \mathbb{N}^* \times \mathbb{Z}^*)^j$.
- ▶ $f(\mathcal{A}, \prod \langle m_i, n_i, k_i \rangle) = \mathcal{A} \prod \langle m_i, n_i, k_i \rangle$
- ▶ $R(\mathcal{A}, \mathcal{B})$ if there is an injection φ from $S_{\mathcal{A}}$ into $S_{\mathcal{B}}$ such that

$$\bar{\varphi} \circ \mathcal{A} = \mathcal{B} \circ \bar{\varphi}$$

- ▶ $\prod \langle m_i, n_i, k_i \rangle \cdot \prod \langle m'_j, n'_j, k'_j \rangle =$
 $\prod \langle m_i m'_j, n_i n'_j, k'_j + n'_j k_i \rangle$
- ▶ Then we need to check the axioms...

The semi-lattice

Theorem. P^-CST induces a sup semi-lattice with the natural operation $\mathcal{A} \times \mathcal{B}$ as a sup operation.

- ▶ Ideals play an interesting role:
 - Reversible CA build a principal ideal,
 - Non-chaoticity build an ideal.

Going further...

- ▶ Continue studying P^- CST.
- ▶ Continue the exploration of classical properties.
- ▶ Is the ideal of non-universal CAs principal ?