

# 104 tiles

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N. Ollinger (Escape, LIF)  
in Frac d'automne 2007

# 104 tiles

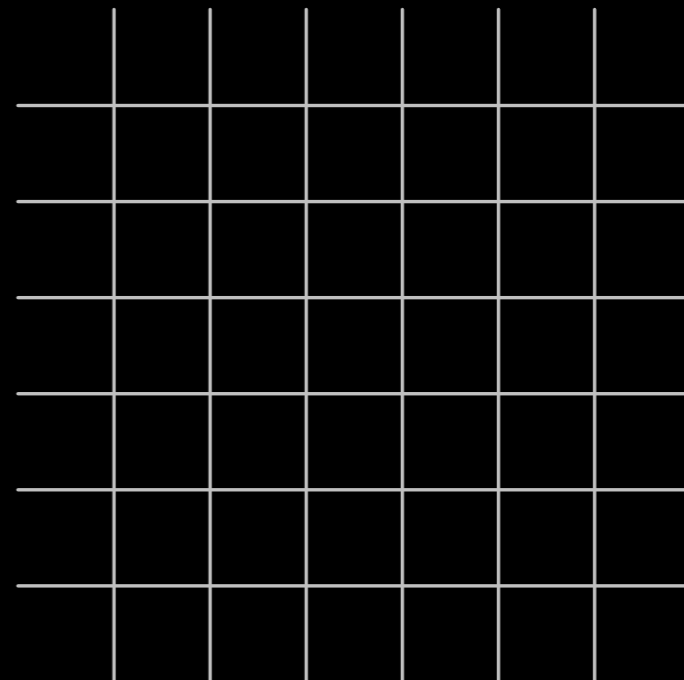
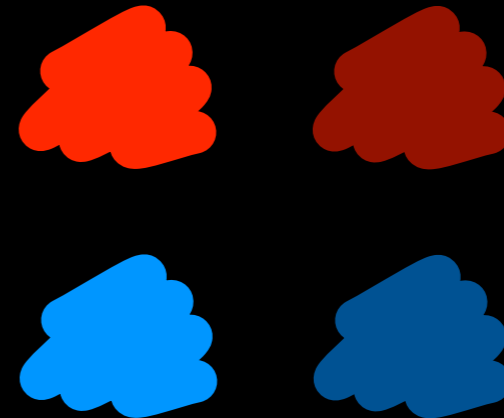
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# Foreword

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- Yet another presentation on the undecidability of the Domino Problem [**Berger 64**]
- Formal stuff (a few) stand on white background.
- Pictures (a lot of) stand on black background.
- This talk is very geometrical: we draw pictures on the Euclidian plane.

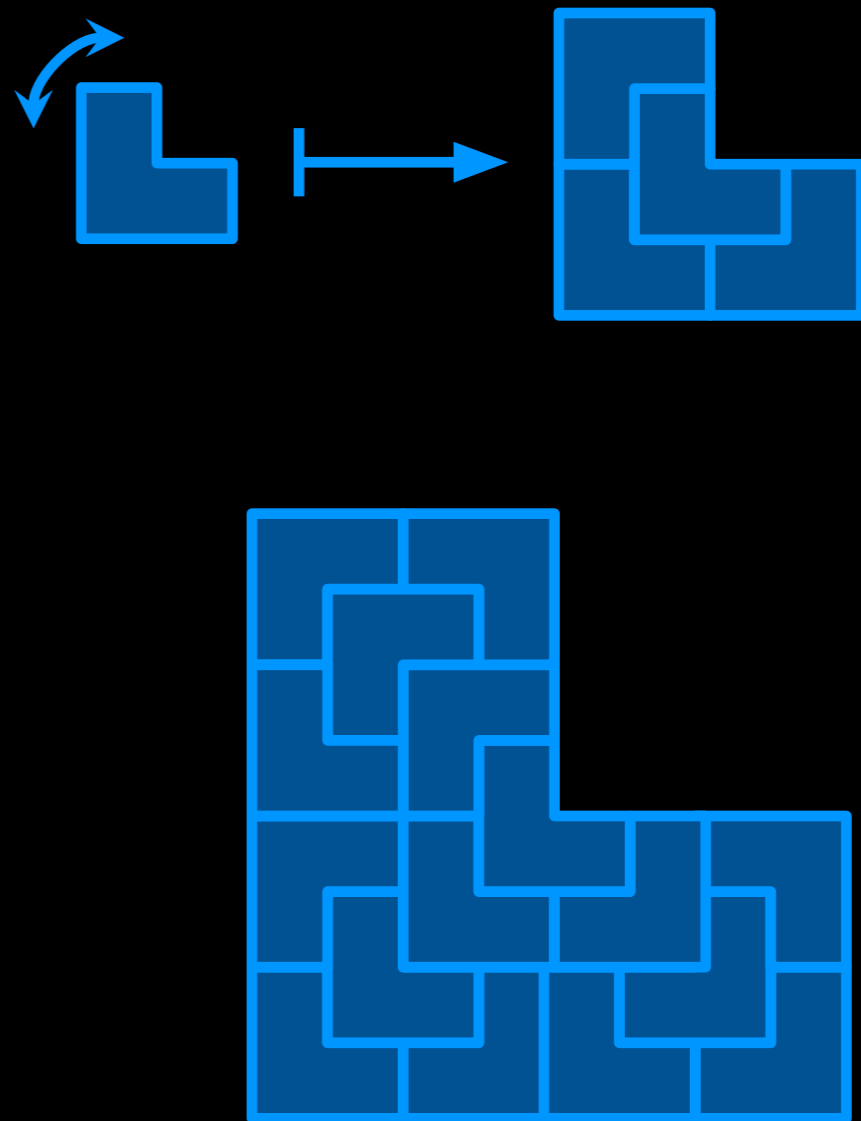


# 2x2 Substitutions

# Substitutions

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- Geometric substitutions provide a convenient recursive way to define aperiodic colorings of the plane.
- Subtle geometrical arguments are required: discuss dissection, inflation, scaling factor, etc.
- The classical L (or chair) substitution.

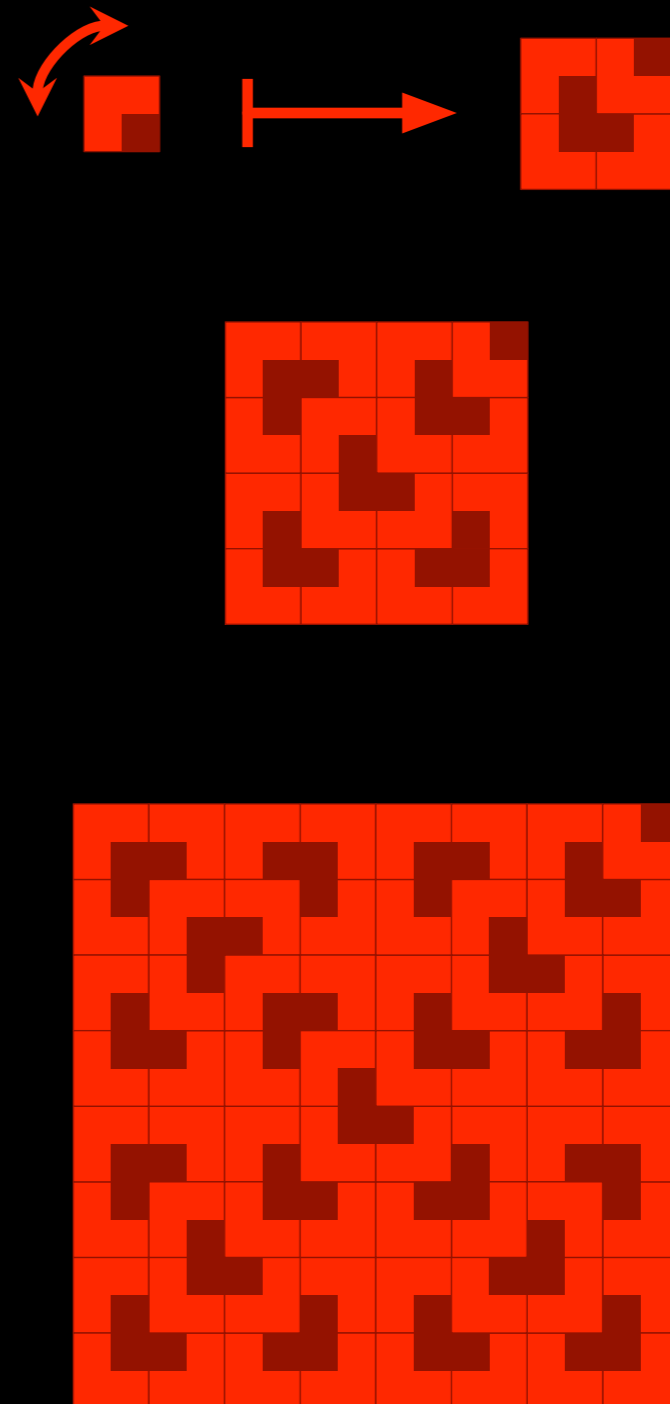


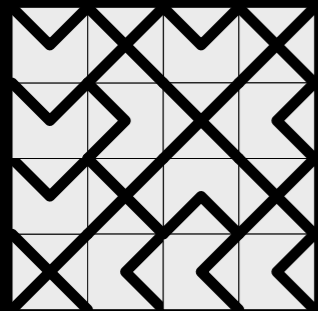
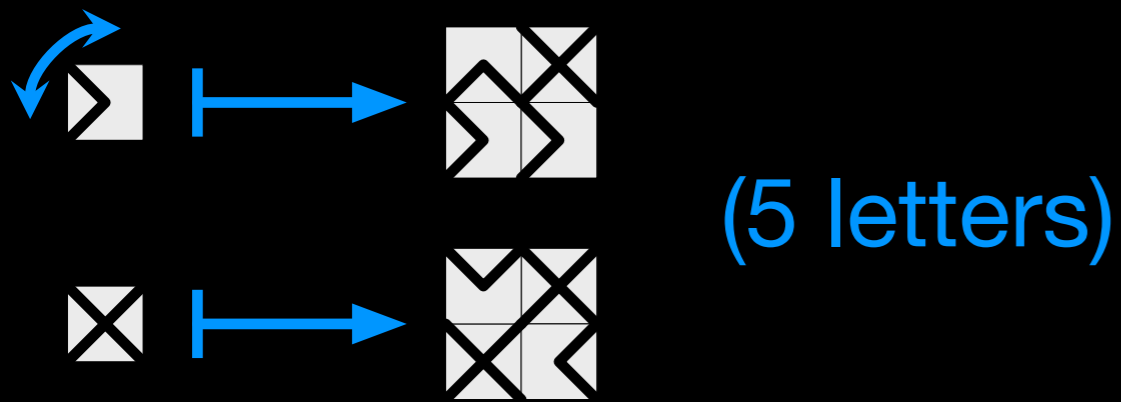
# 2x2 Substitutions

- A 2x2 substitution maps a finite alphabet to 2x2 squares of letters on that alphabet.
- The substitution is iterated to generate bigger squares.

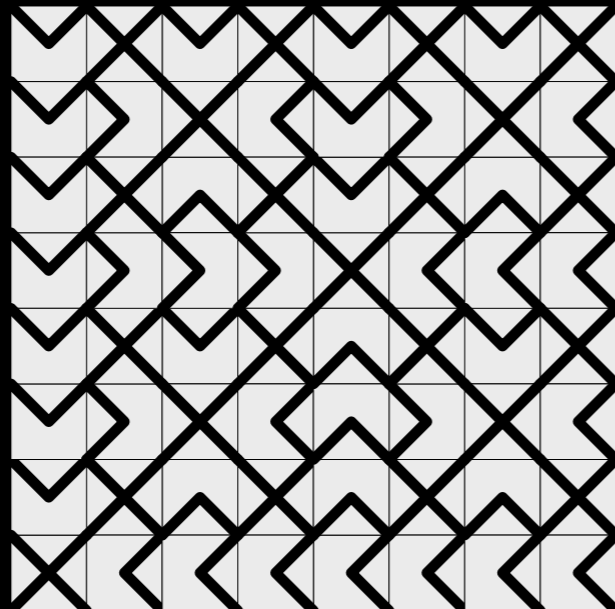
$$s : \Sigma \rightarrow \Sigma^{\boxplus}$$

$$S \circ \sigma_u = \sigma_{2u} \circ S$$

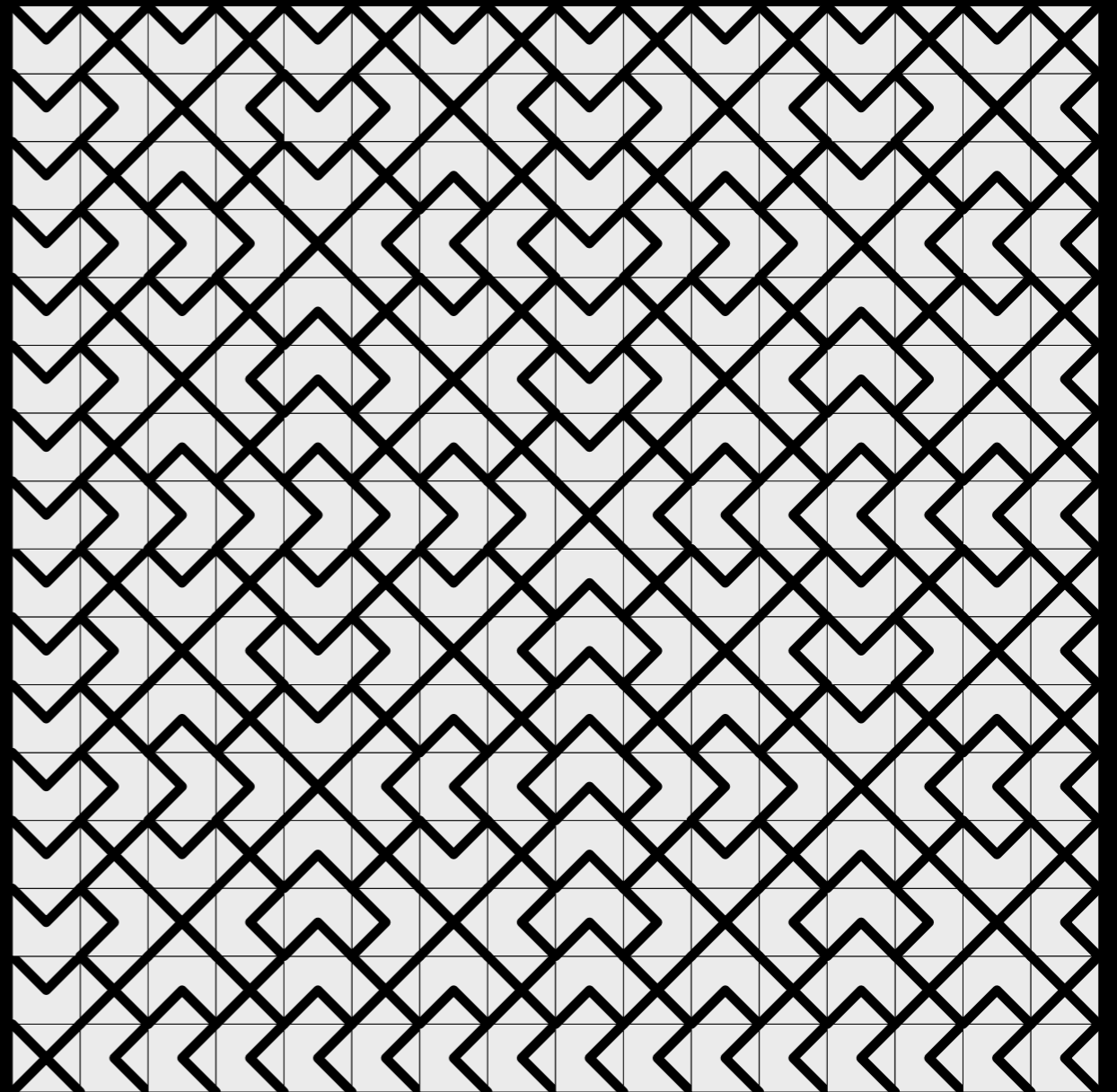




$s^2$  (X)



$s^3$  (X)



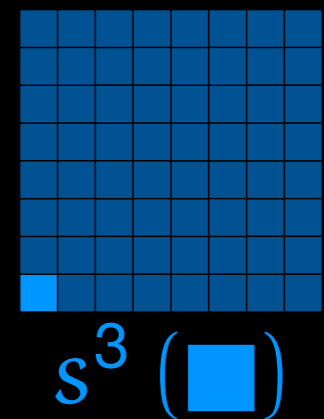
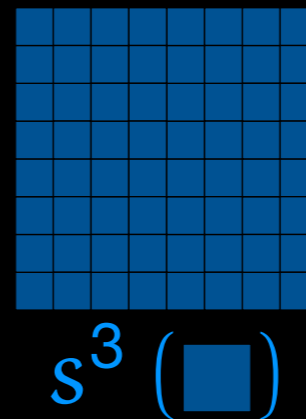
$s^4$  (X)

Another L-style substitution

# Pattern closure?

- What is a coloring of the plane generated by a substitution?
- One possibility is to consider a *closure* of generated patterns up to translations.
- Difficult to check...

$$C \in \Omega_s \text{ if } \forall P \prec_{\mathcal{F}} C, \exists a, k, P \prec s^k(a)$$



$$\Omega_s = \left\{ \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \square & \square & \square \\ \hline \square & \square & \square \\ \hline \end{array} \right\}$$



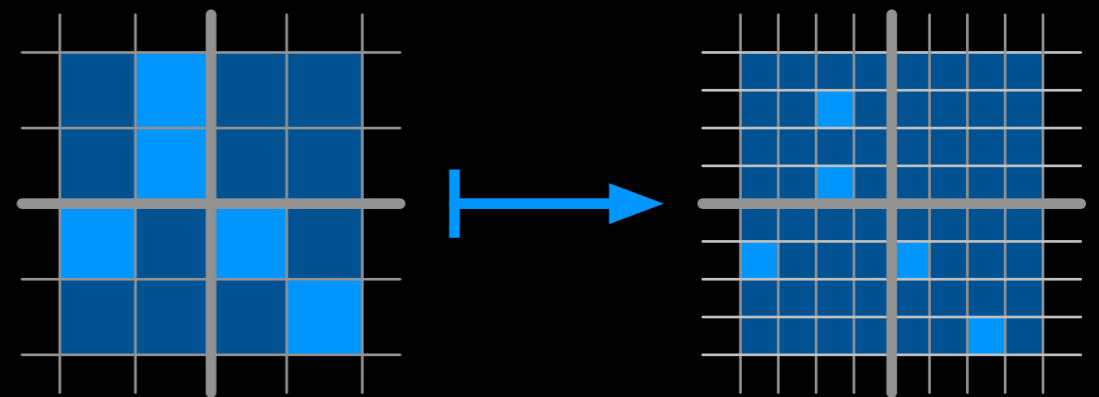
# Limit set!

- Extend the substitution to a global map inflating around 0.
- Take the limit set of iterations of this map closed up to translations.
- More colorings!

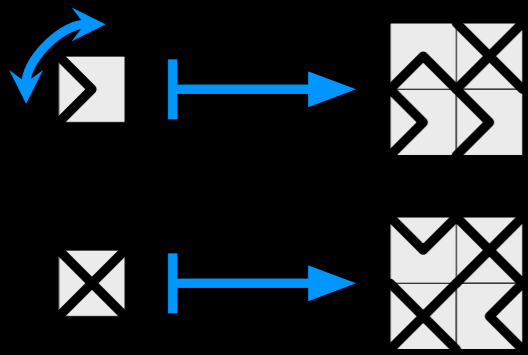
$$\Lambda_s = \bigcap_n \Lambda_s^{(n)} \text{ where}$$

$$\Lambda_s^{(0)} = \Sigma^{\mathbb{Z}^2}$$

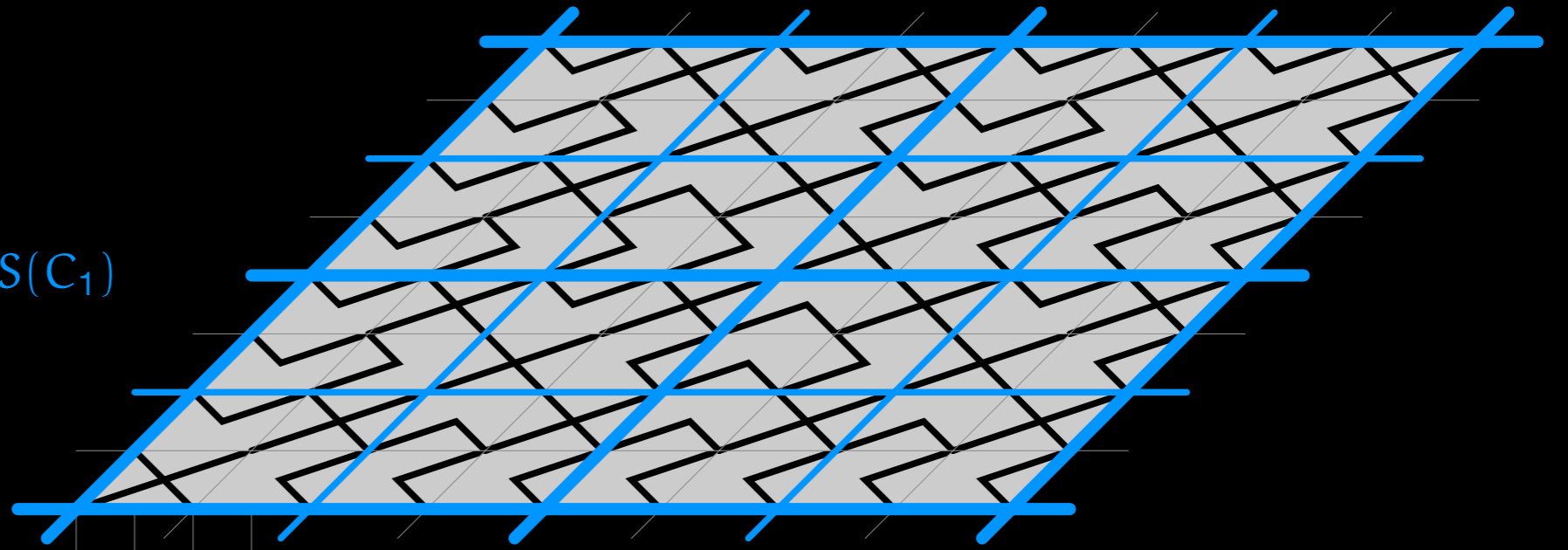
$$\Lambda_s^{(n+1)} = \left\langle S(\Lambda_s^{(n)}) \right\rangle_{\sigma}$$



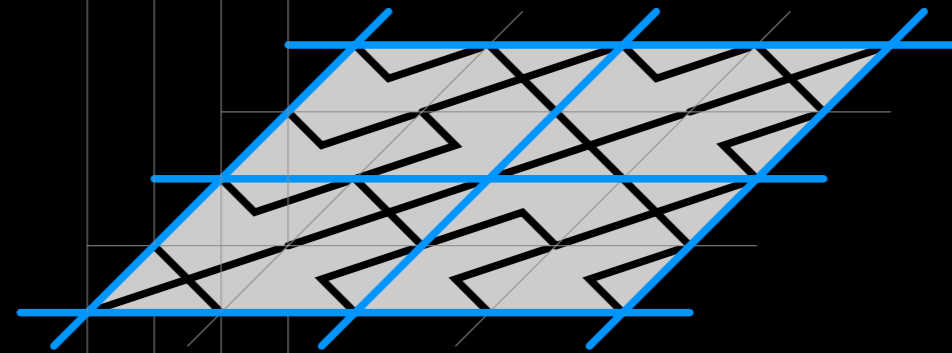
$$\Lambda_s = \Omega_s \cup \left\{ \begin{array}{c} y \\ \text{[2x2 grid with center square highlighted]} \\ x \end{array} \right\}_{x,y \in \mathbb{Z}^2}$$



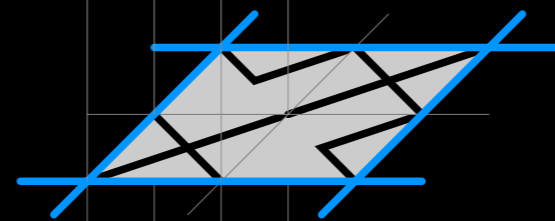
$$C_0 = \sigma_0 \circ S(C_1)$$



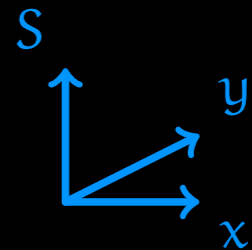
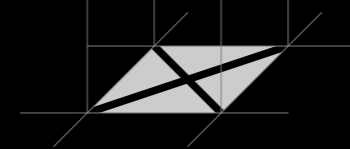
$$C_1 = \sigma_1 \circ S(C_2)$$



$$C_2 = \sigma_2 \circ S(C_3)$$



$$C_3 = \sigma_3 \circ S(C_4)$$



History of a coloring in the limit set

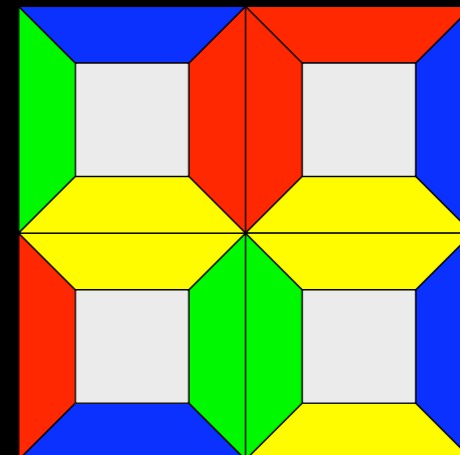
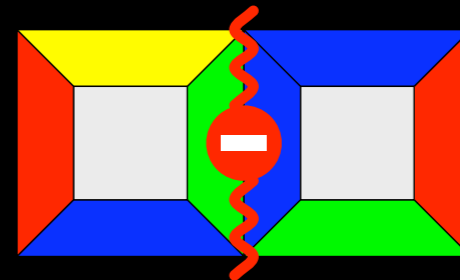
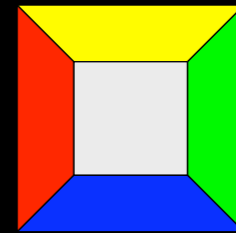
Tiling with Wang tiles

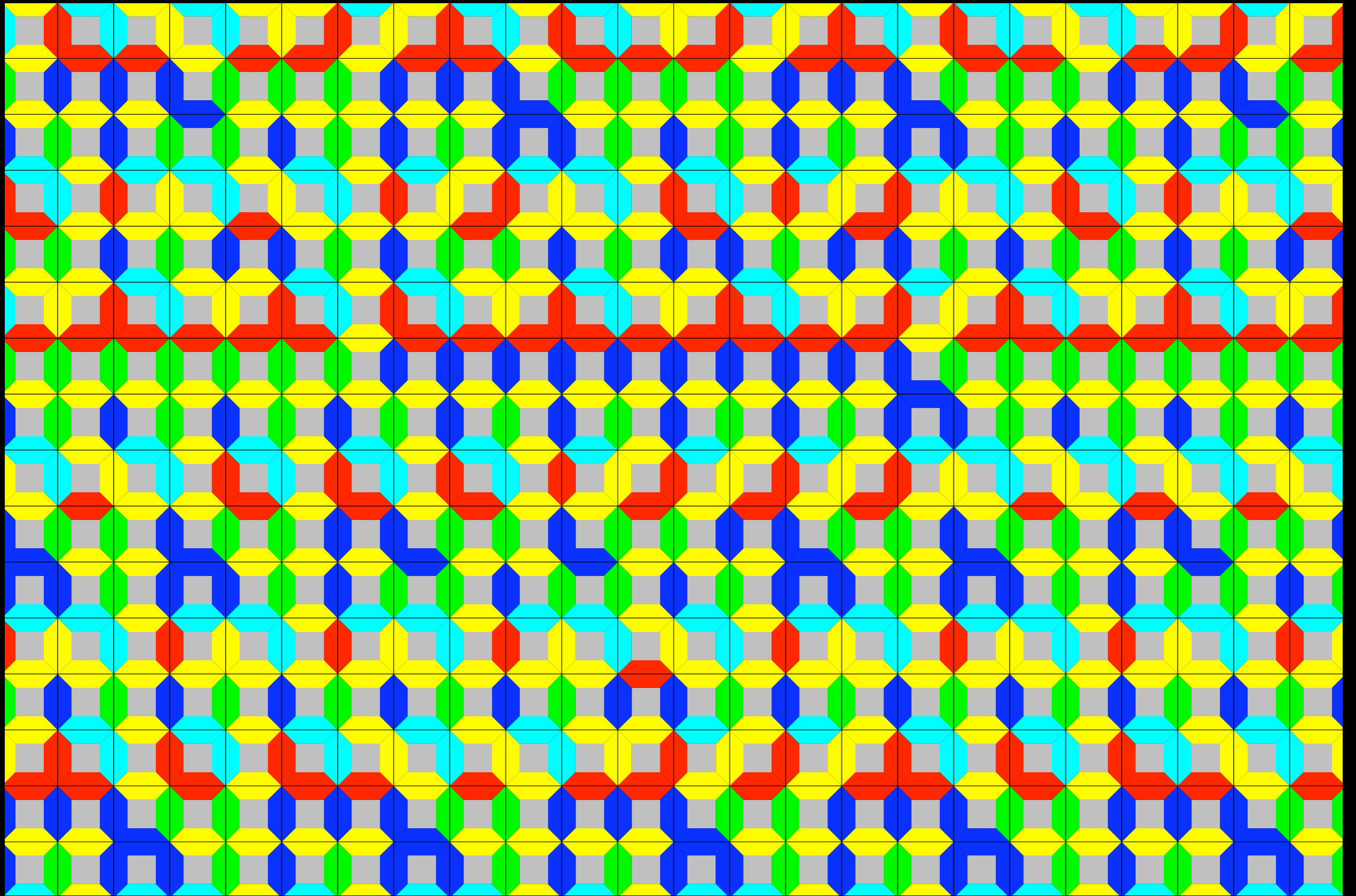
# Wang tiles

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- A Wang tile is a unit square with colored edges.
- A tiling is a finite set of Wang tiles.
- Tiling is done with translations only (no rotation) by matching colors along edges.

$$\tau \subseteq \mathbb{C}^4$$



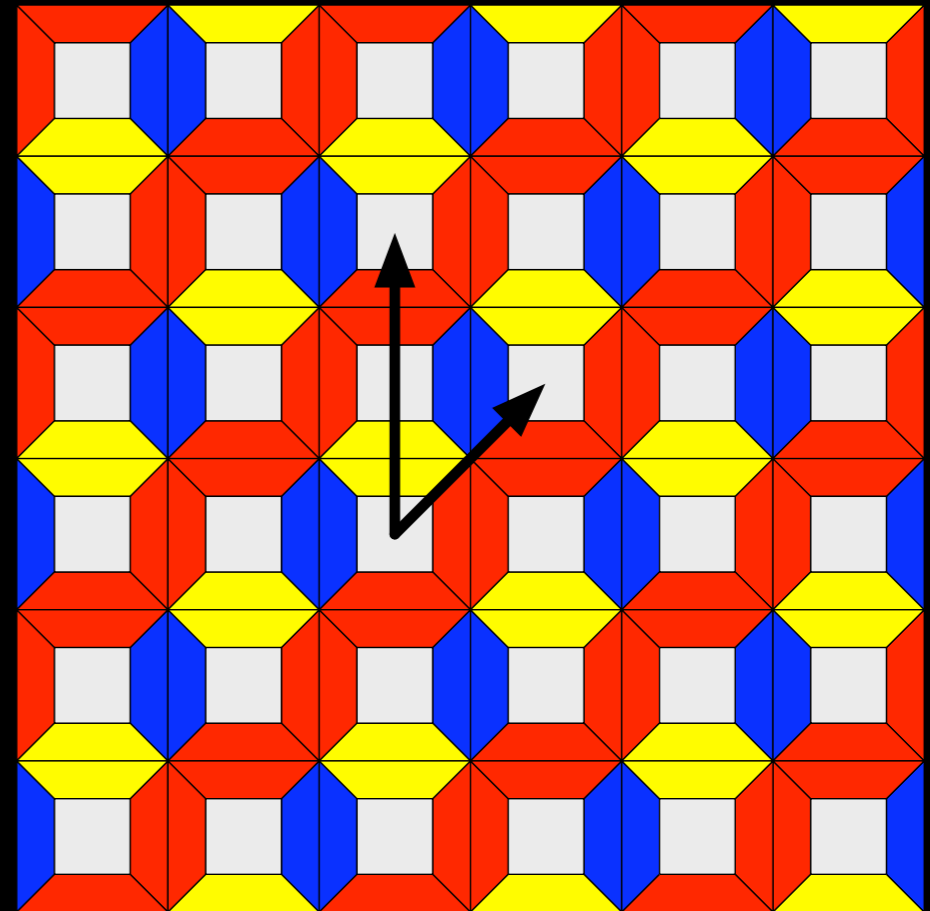


Tiling the whole plane with Wang tiles

# Periodicity

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- A tiling is periodic if it admits a periodicity vector.
- A tiling is biperiodic if it admits two non colinear periodicity vectors.
- If a tileset admits a periodic tiling then it admits a biperiodic tiling.
- A tileset is aperiodic if it admits no periodic tiling.



# The Domino Problem

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- **The Domino Problem:** given a tileset, decide if it admits a tiling.
- A long time ago, Wang conjectured there exists no aperiodic tileset and thus the Domino Problem is recursive.
- In 1964, his PhD student Berger proved that the problem is not recursive. To achieve this goal he encoded Turing machine computations inside aperiodic tilings.
- A whole talk could be devoted to the history of aperiodicity, the quest for small aperiodic tileset, the different proofs and techniques to prove the non recursivity of the domino problem... **Not today!**

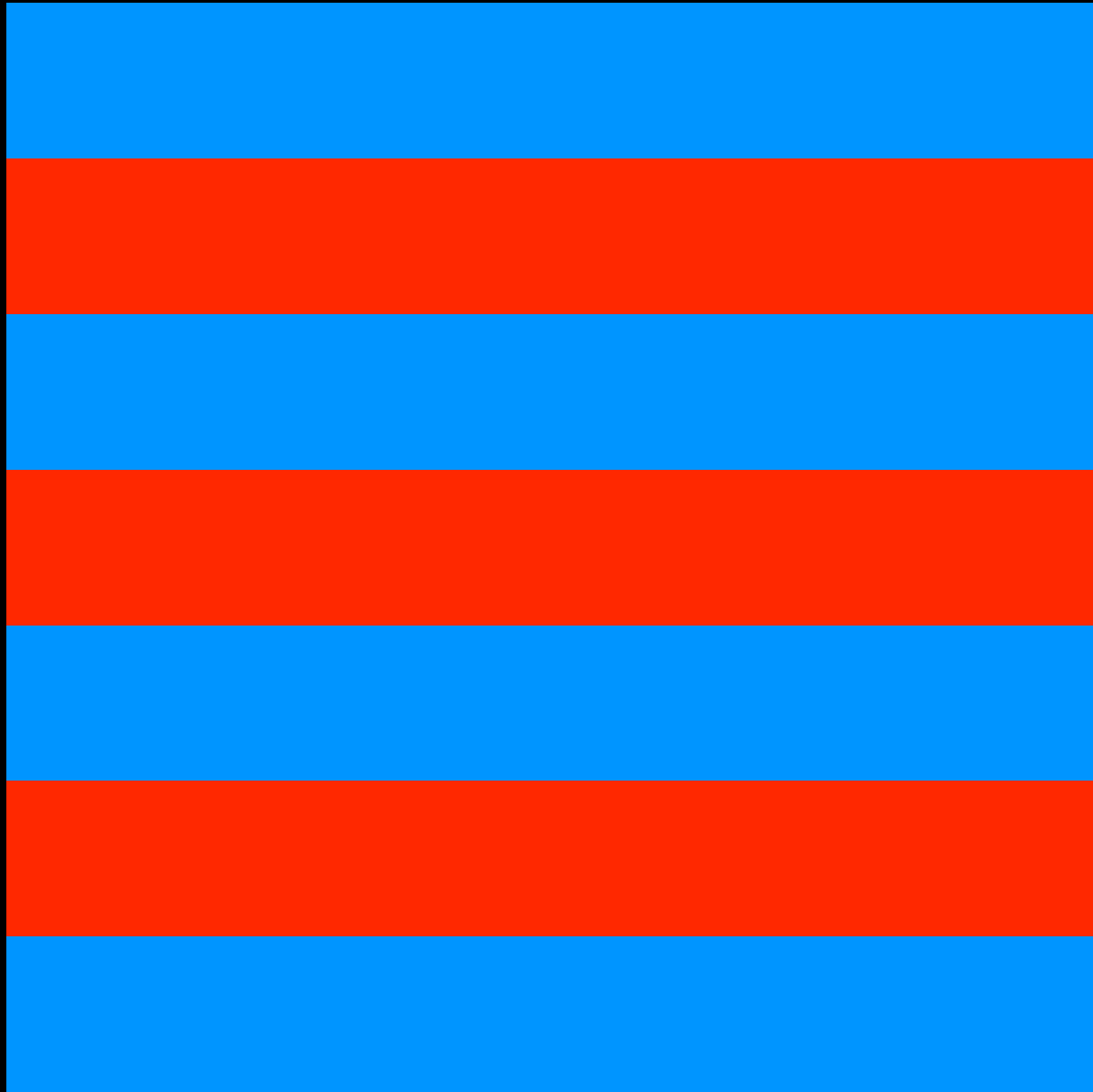
# Au menu

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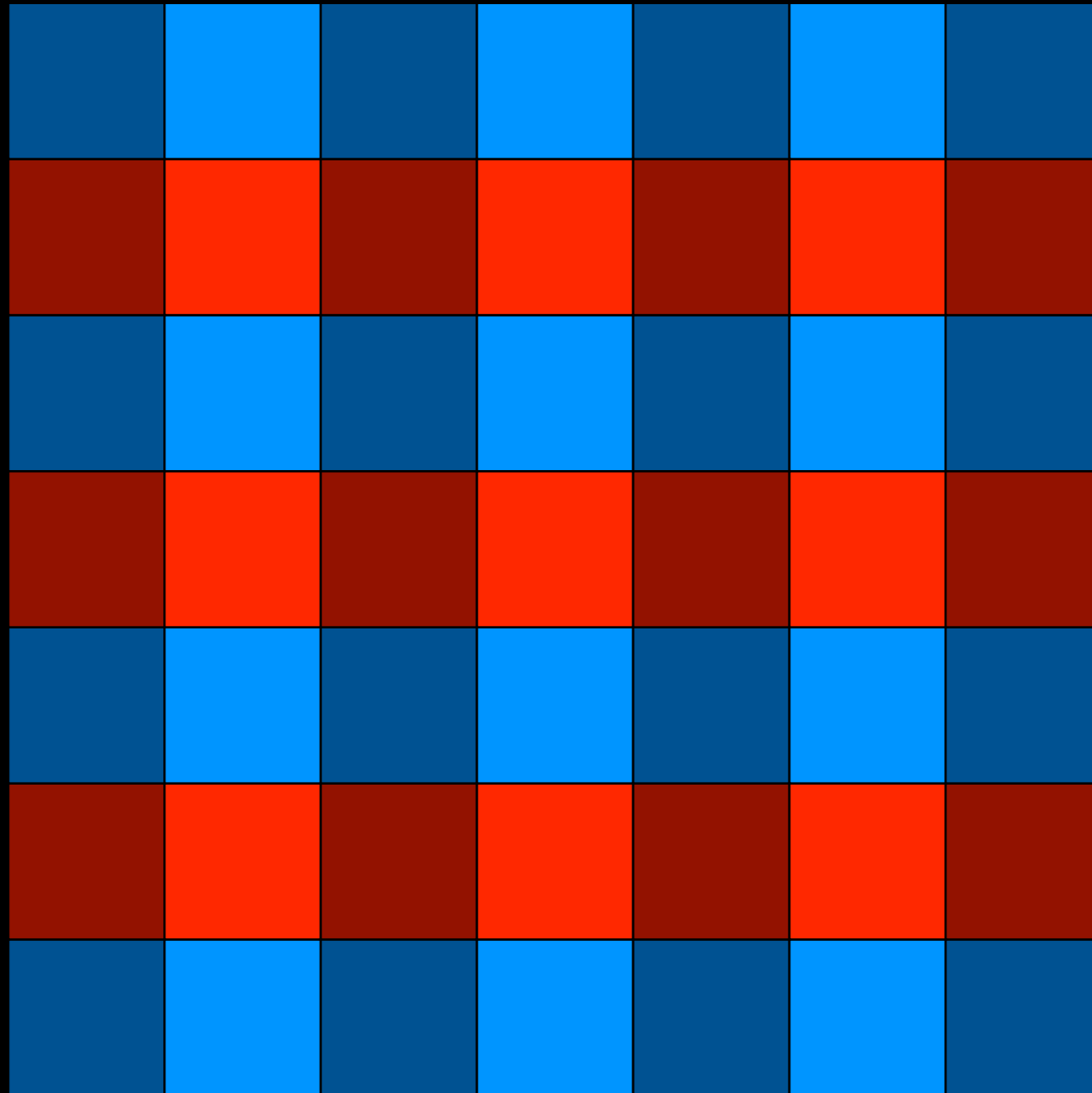
- Define a particular set of 104 tiles;
- Prove in details its aperiodicity;
- Briefly explain how to complete the Domino Problem construction.
- This is a mix of different tools and ideas from:
  - [Berger 64]** *The Undecidability of the Domino Problem*
  - [Robinson 71]** *Undecidability and nonperiodicity for tilings of the plane*
  - [Grünbaum Shephard 89]** *Tilings and Patterns, an introduction*
  - [Durand Levin Shen 05]** *Local rules and global order, or aperiodic tilings*



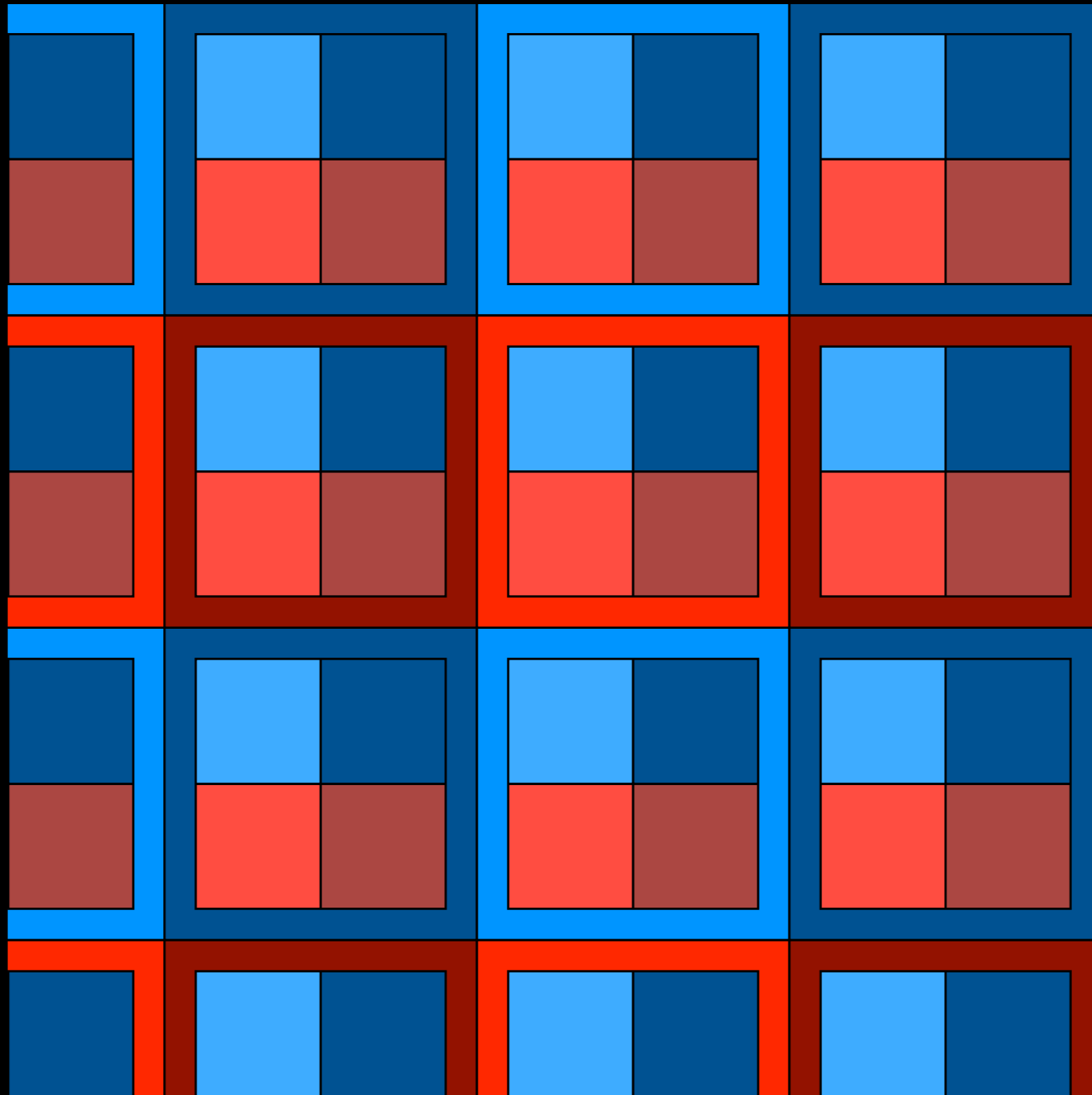
A set of 104 tiles



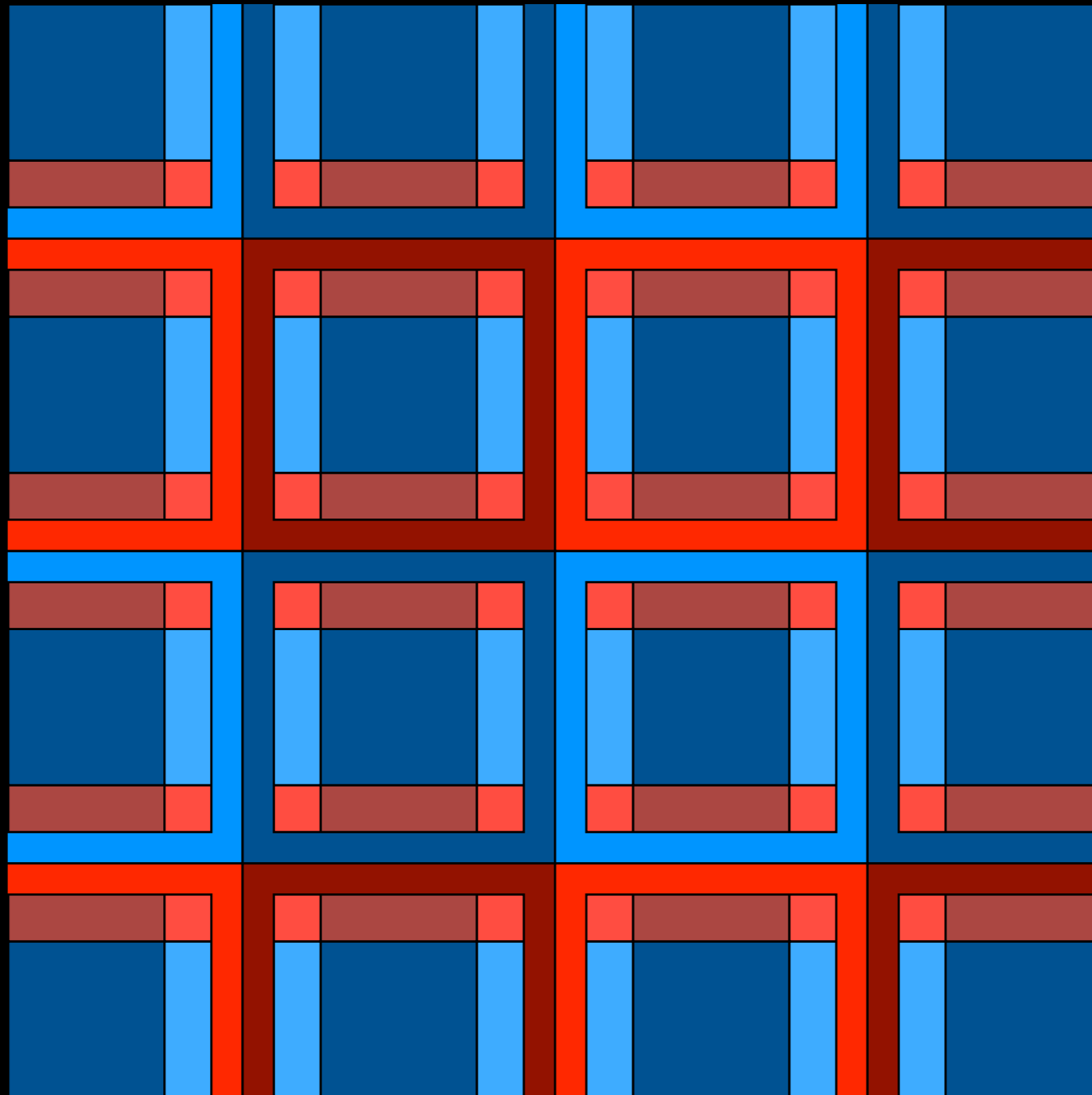
Encoding an infinite number of grids



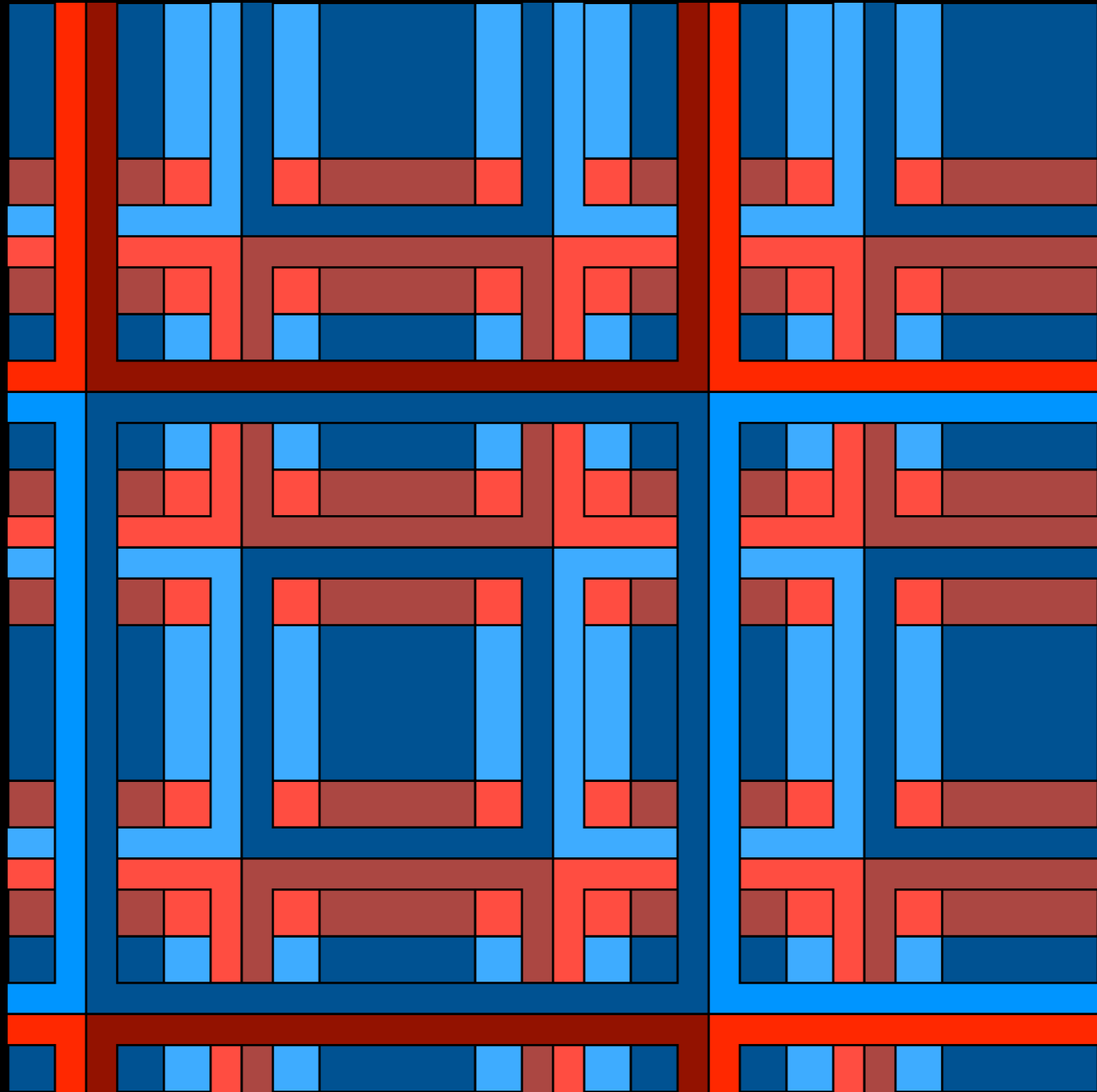
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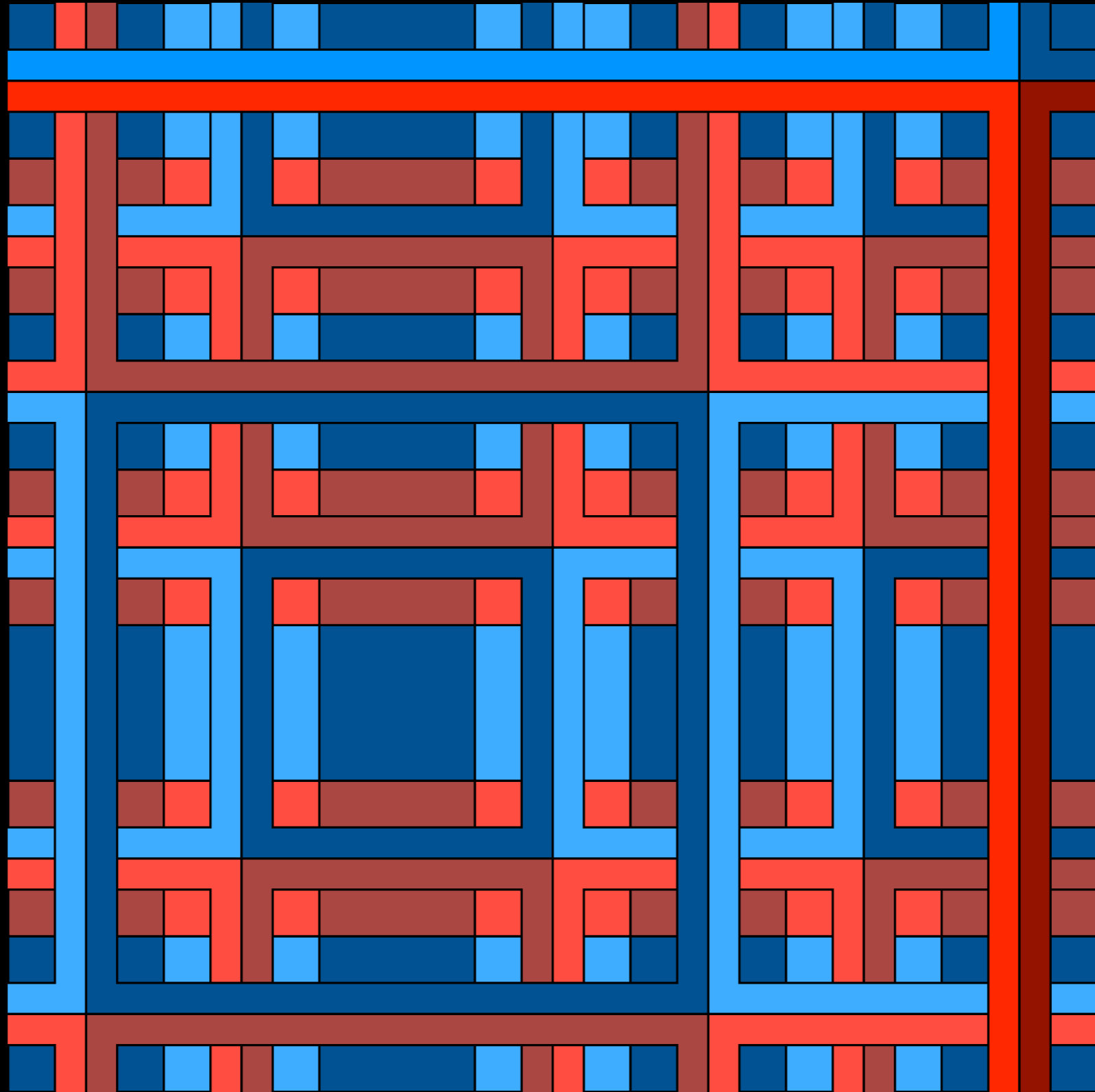
Encoding an infinite number of grids



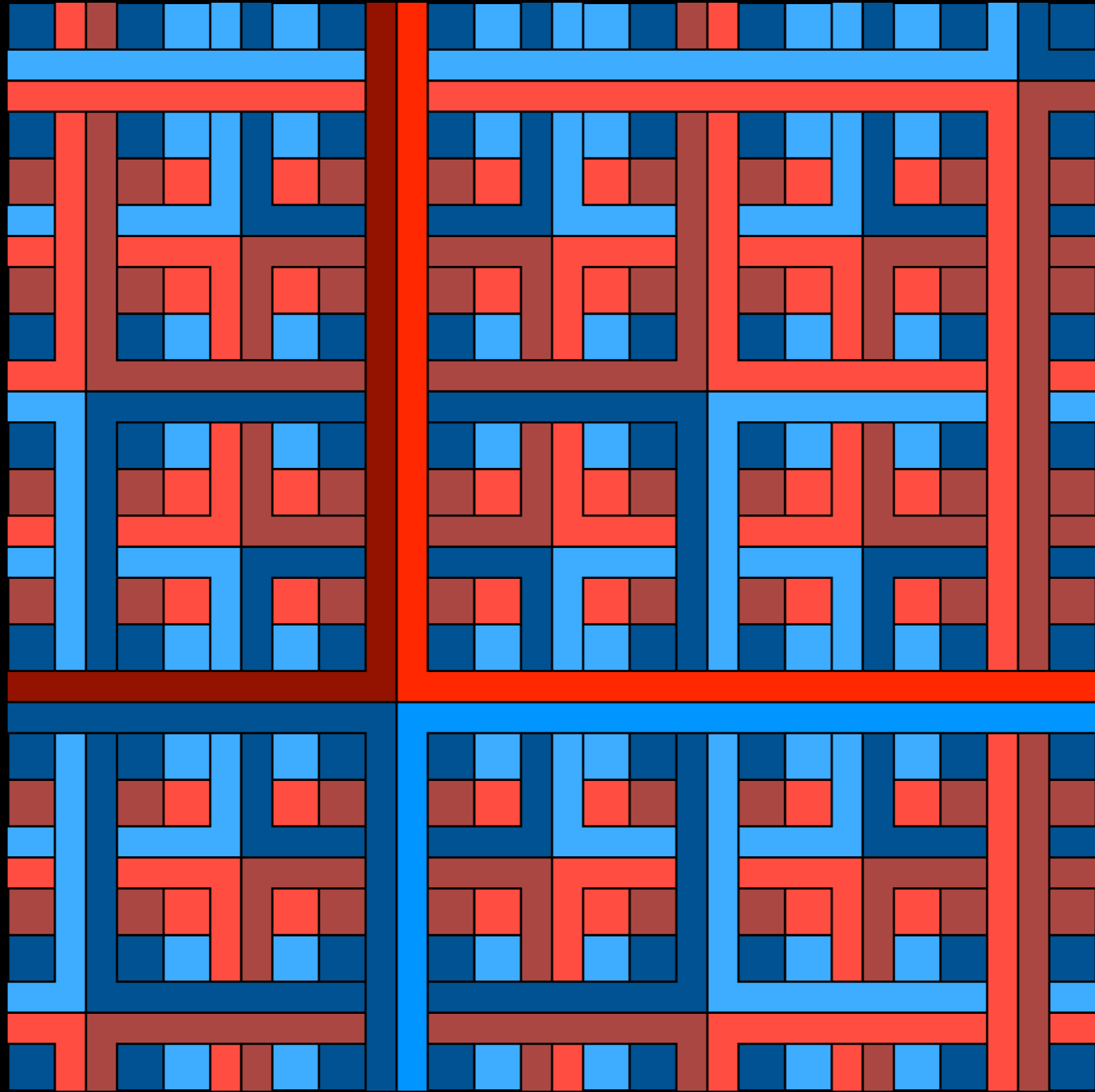
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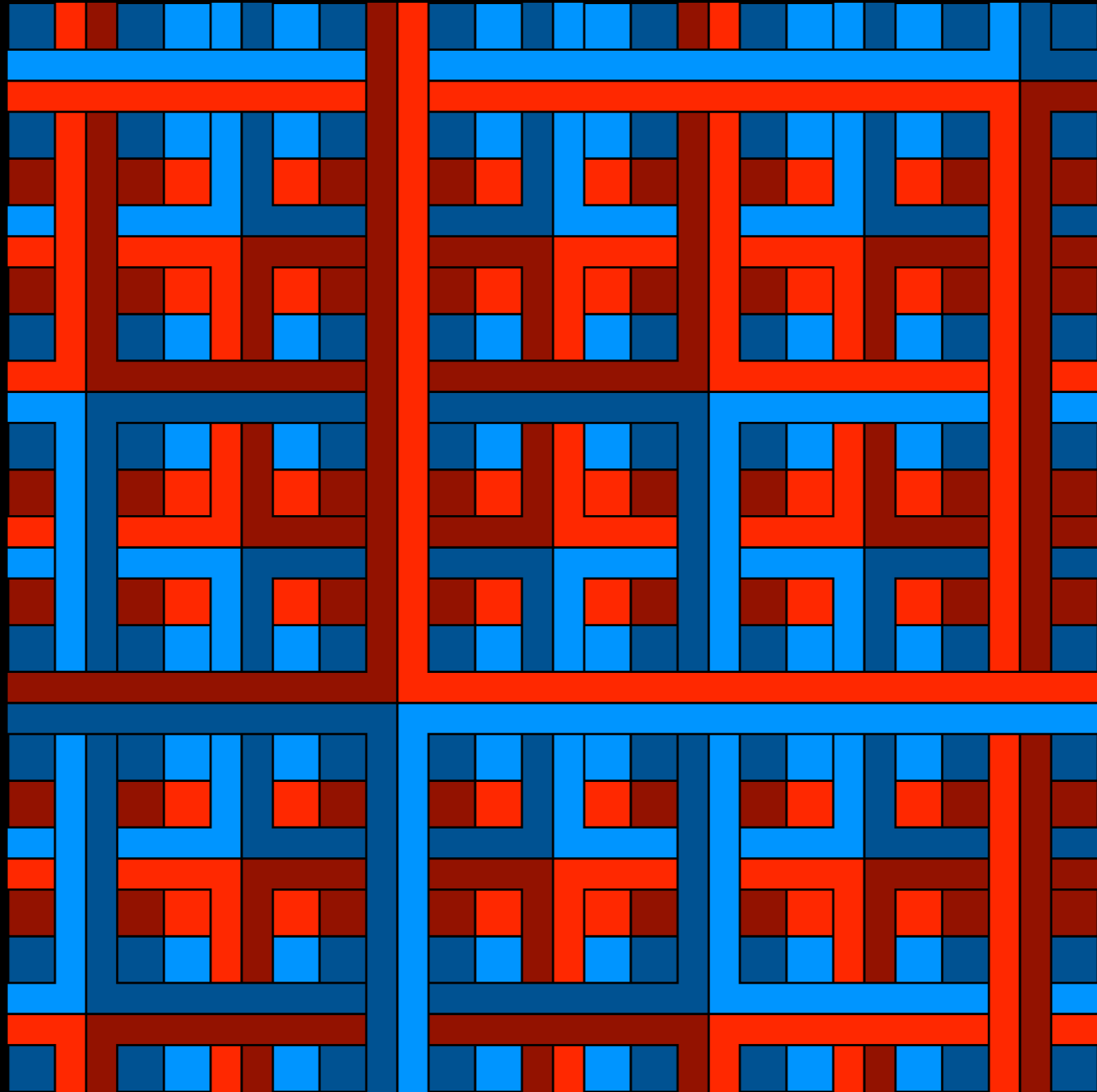


Encoding an infinite number of grids



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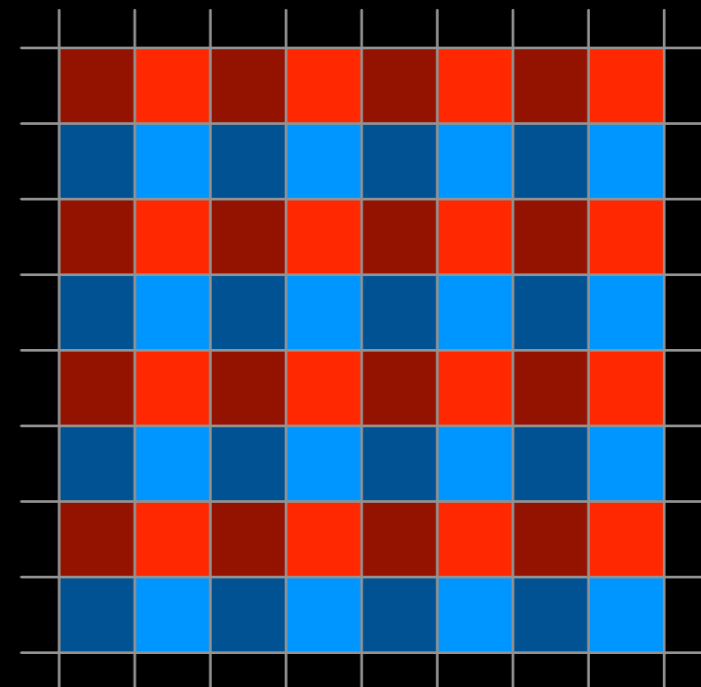
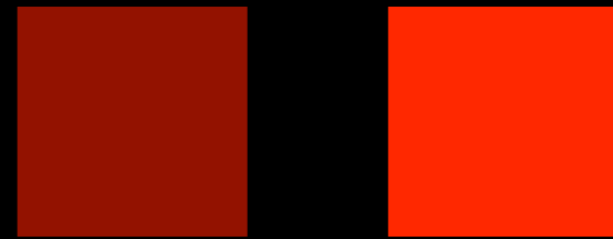


Encoding an infinite number of grids

# Layer 1. initialization

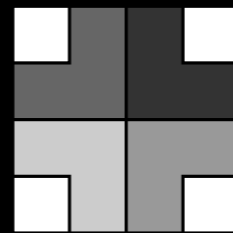
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- 4 tiles to group tiles 2×2.
- Alternate red/blue vertically.
- Alternate light/dark horizontally.
- Simple and very constrained matching rule.



# Layer 2. grids

- Tiles carry wires, two on each edge.
- Three kinds of tiles: X, H, V.
- Pairs of wires should be compatible.
- Matching rule: wire colors should match.
- Restrict layer 1 by kind.



4



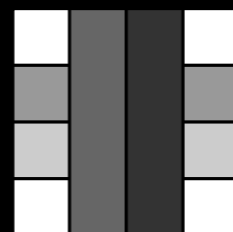
8



16



32



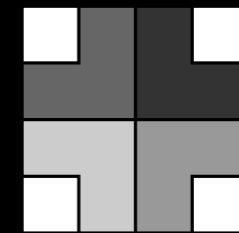
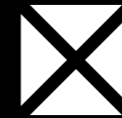
16



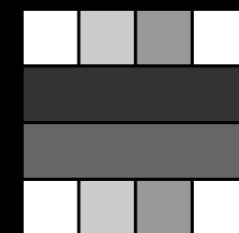
32

# Layer 3. square-ifier

- One more bit of information per edge to force squares.
- 5 different tiles with arrows.
- Matching rule: arrows should match along edges.
- Restrict arrows by kind.
- **Important:** red wire is a one-way lane.

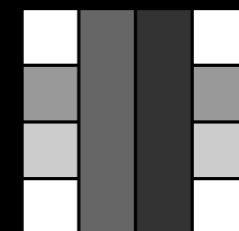


8



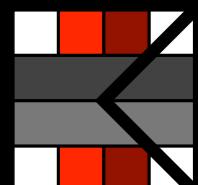
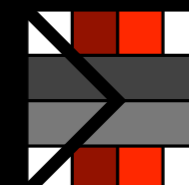
~~64~~

48



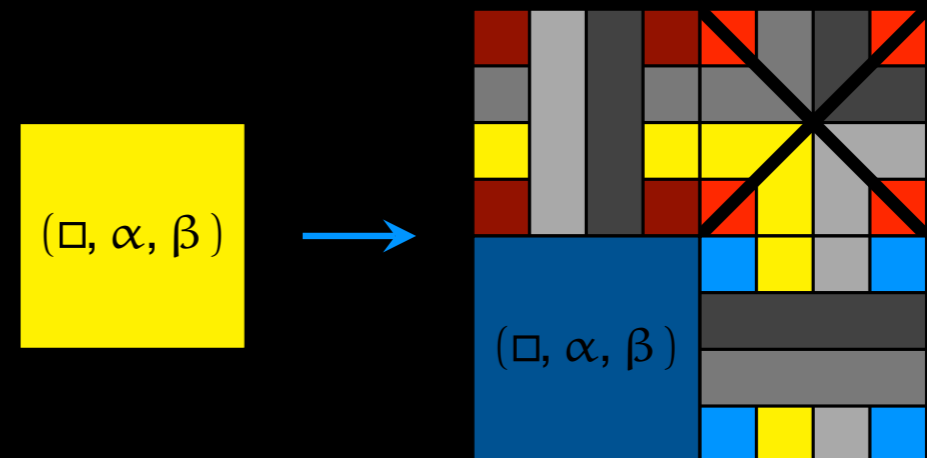
~~64~~

48

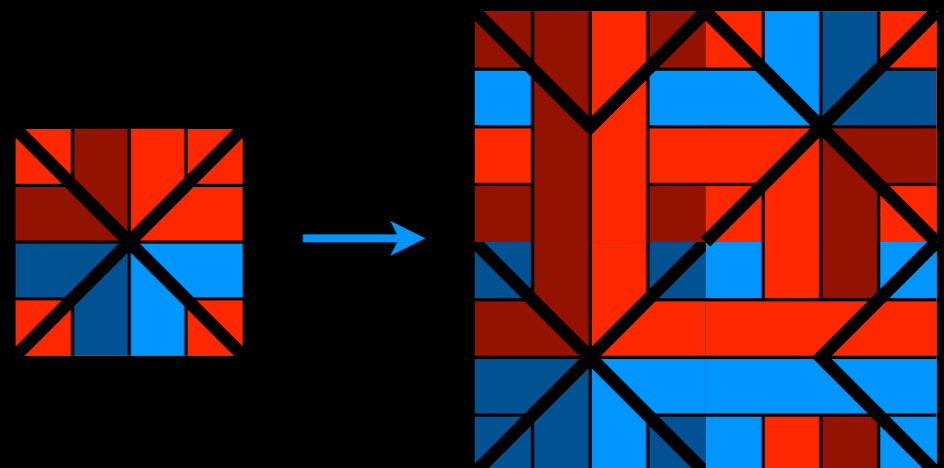


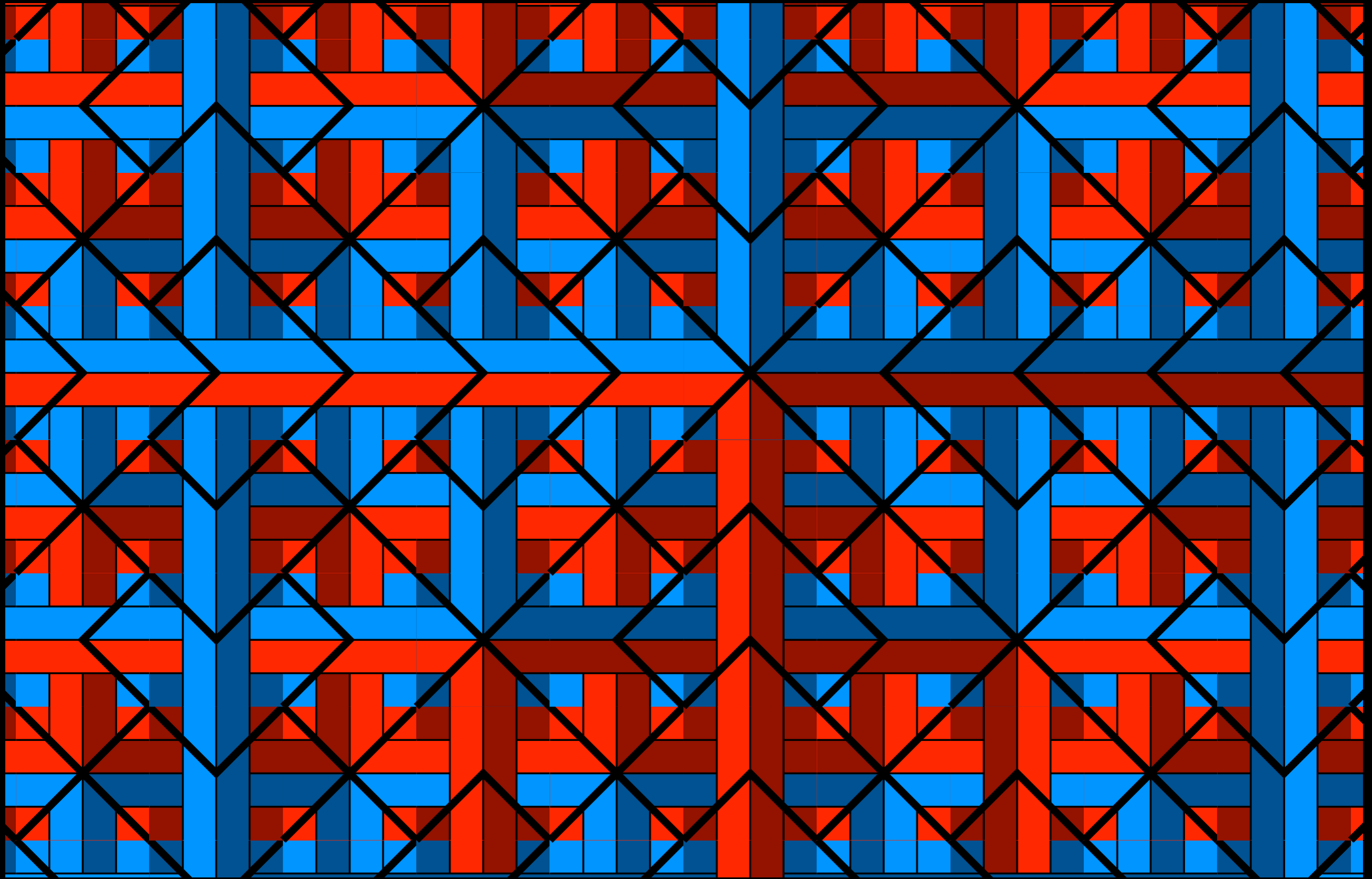
# Canonical substitution

- Copy the tile in the SW corner but for layer 1.
- Put the only possible X in NE that carry layer 1 of the original tile on SW wire.
- Propagate wires colors.
- Let H et V tile propagate layer 3 arrows.
- The substitution is injective.

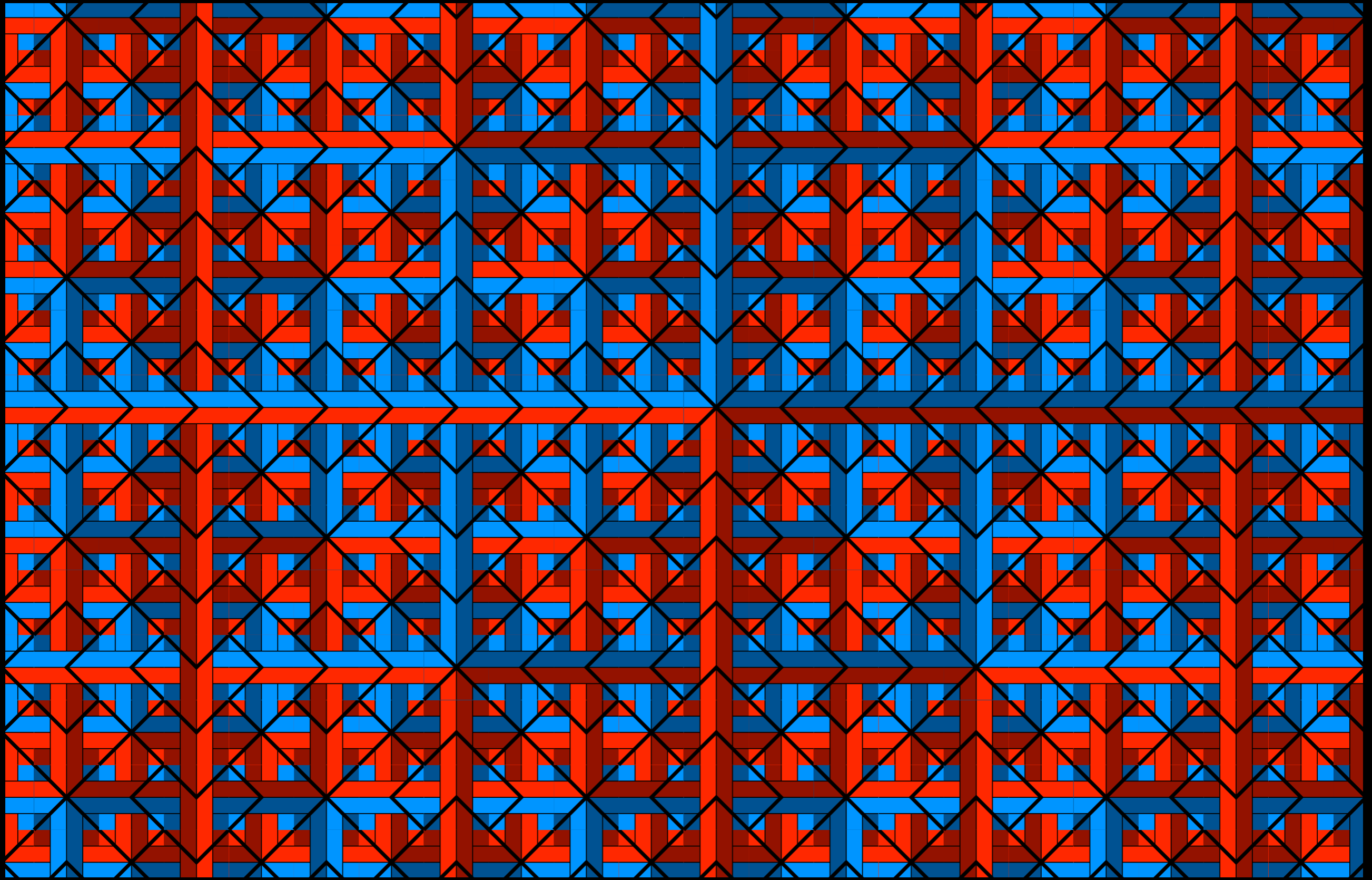


example:





The tiling can tile the whole plane...

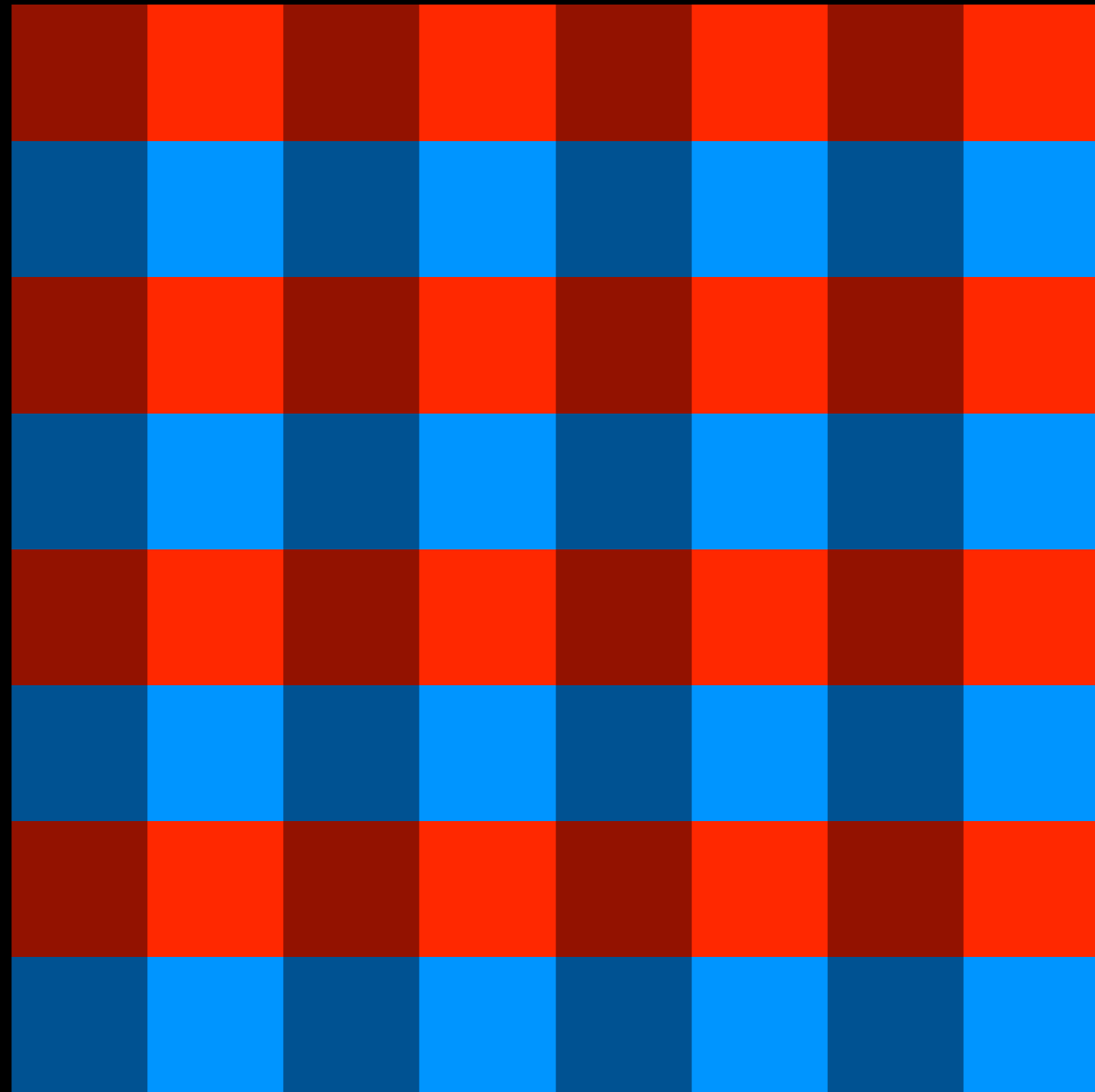


...extend *ad lib!*

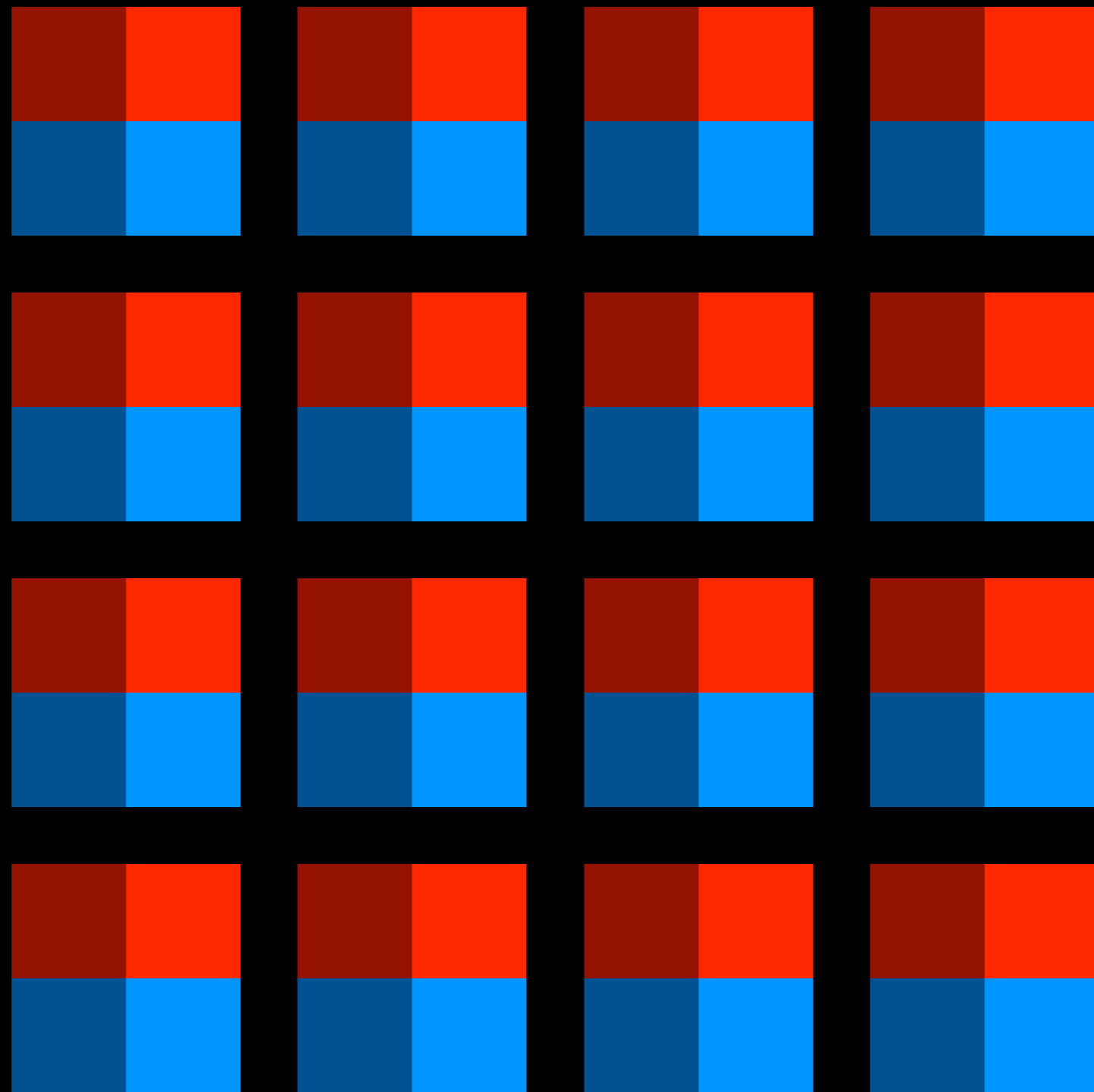
This set of tiles is aperiodic



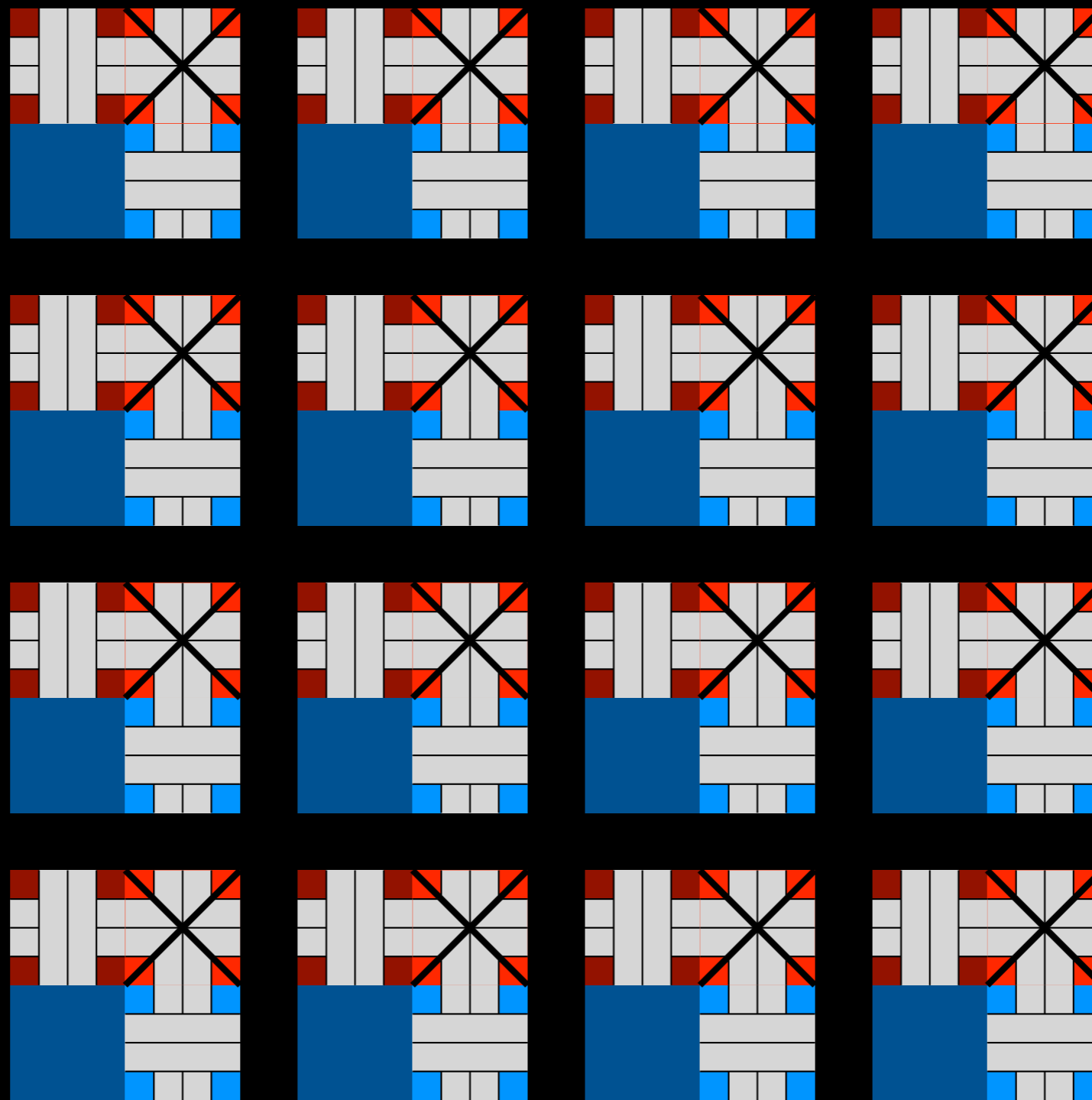




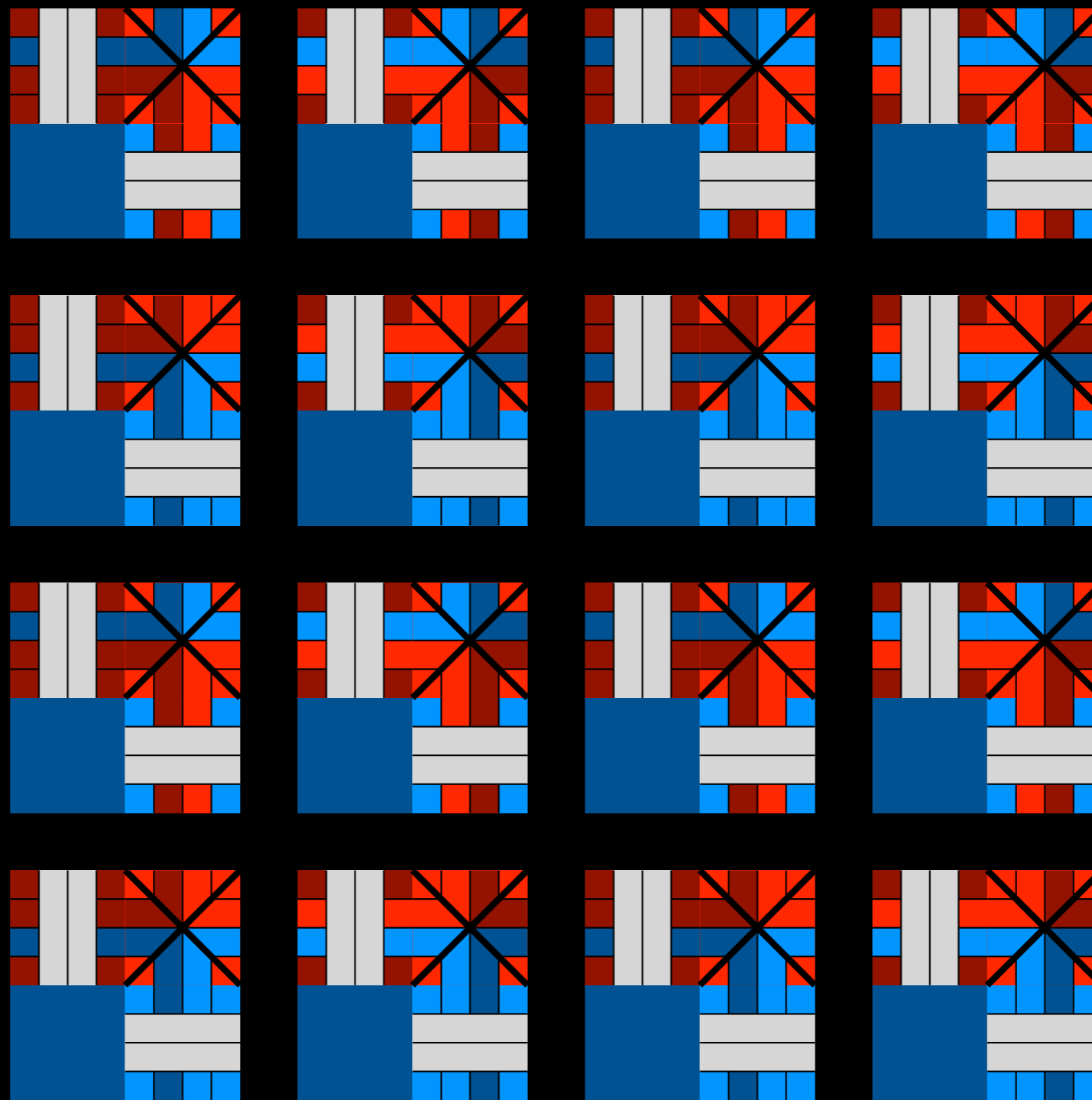
Lemma 1. each tiling from this set is an image of the canonical substitution.



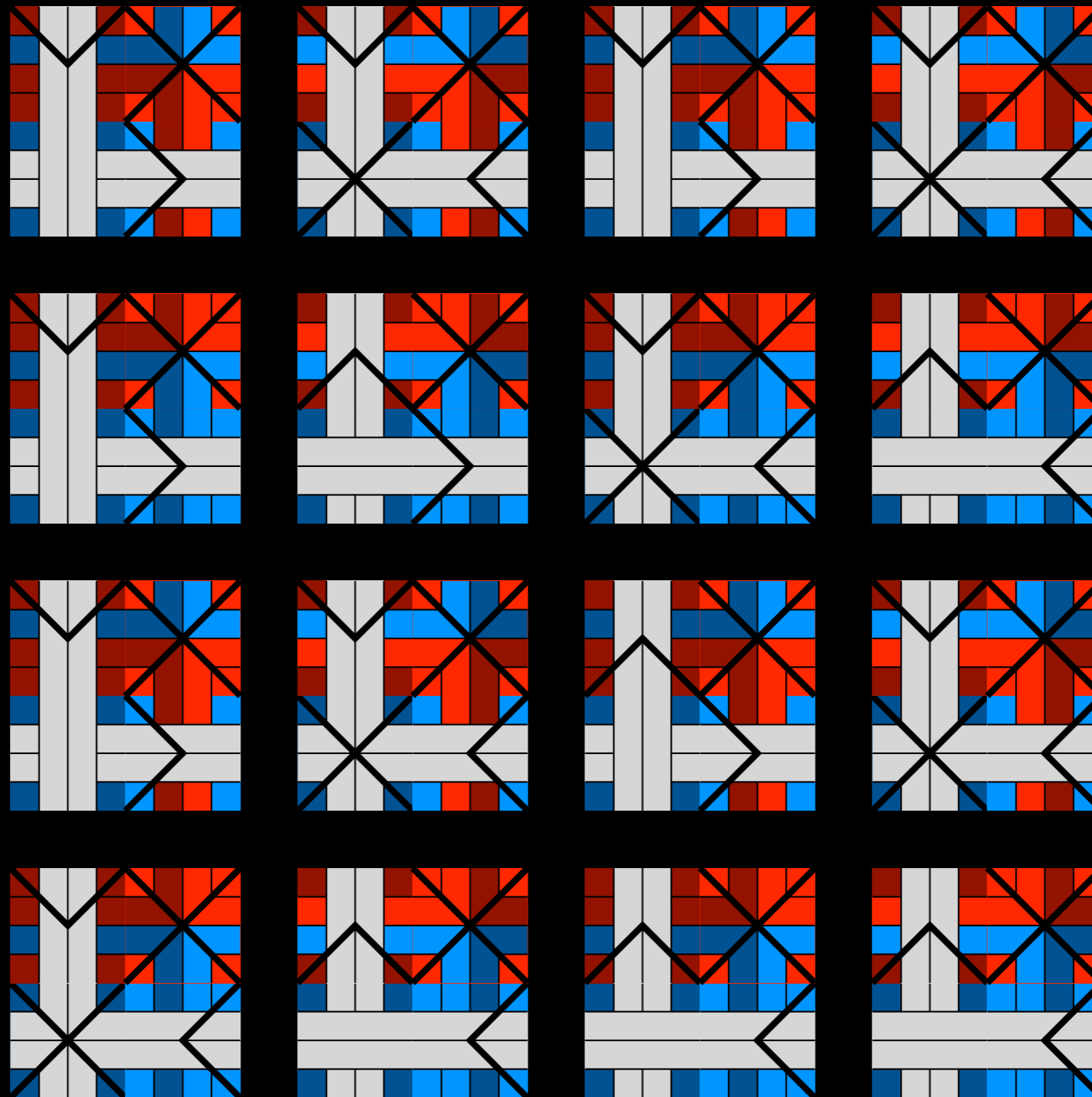
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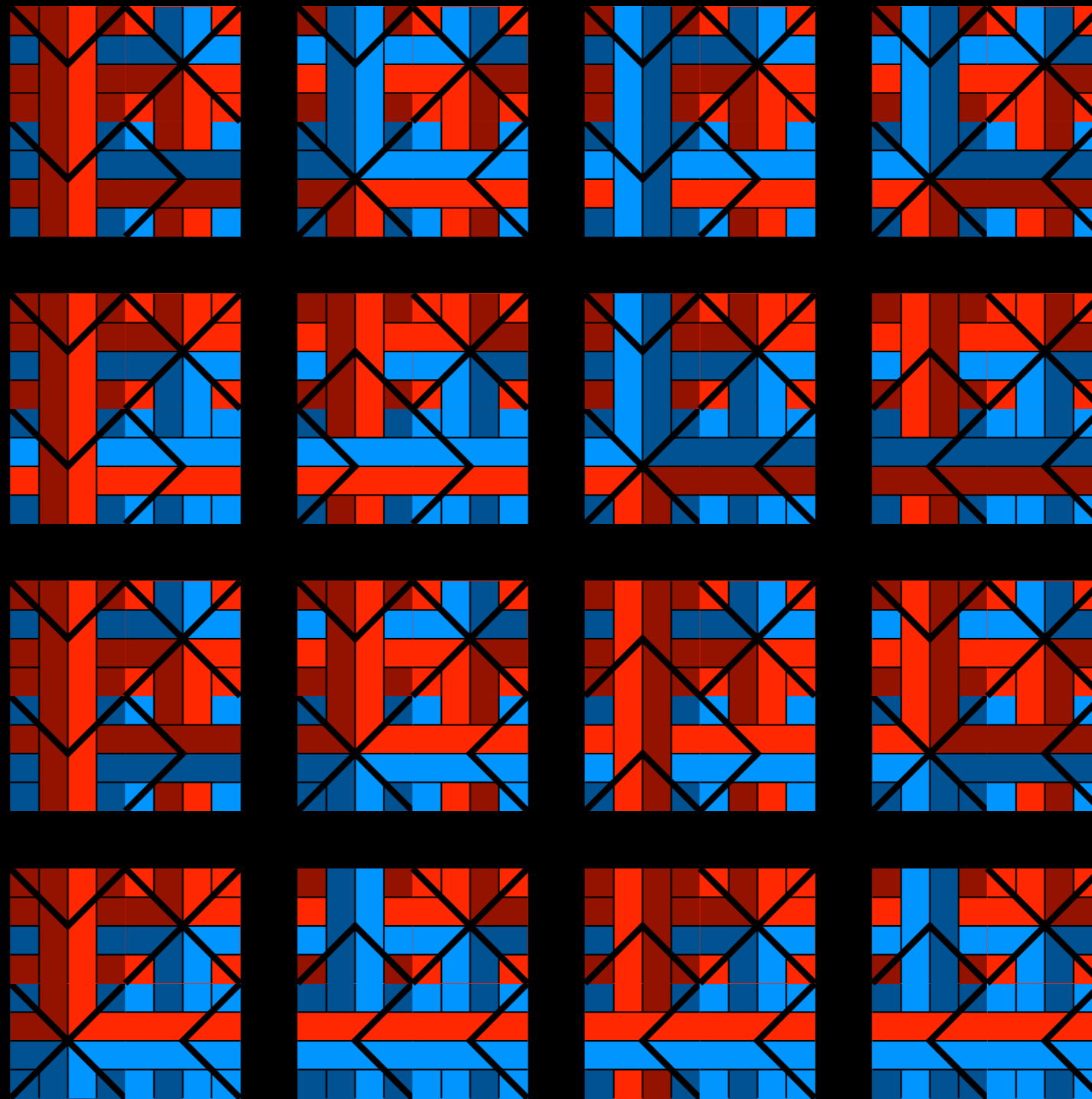
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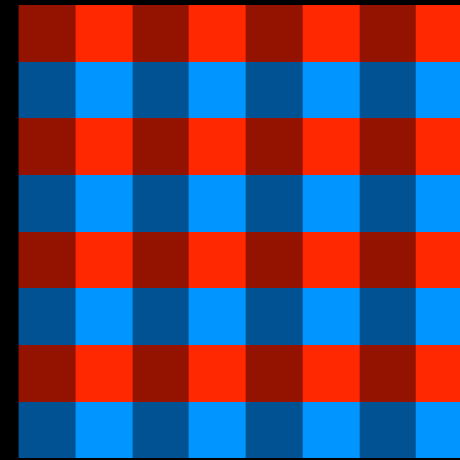
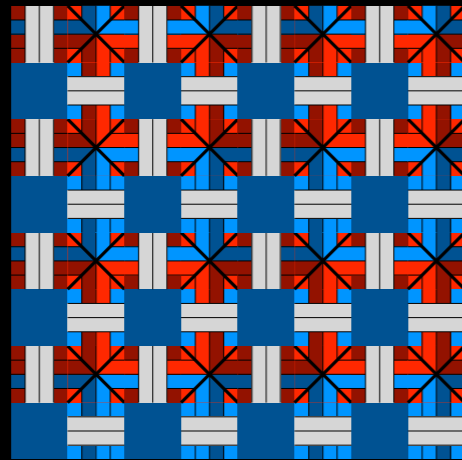
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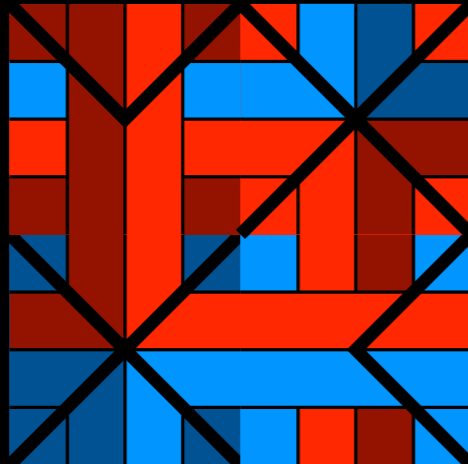
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Layer 1

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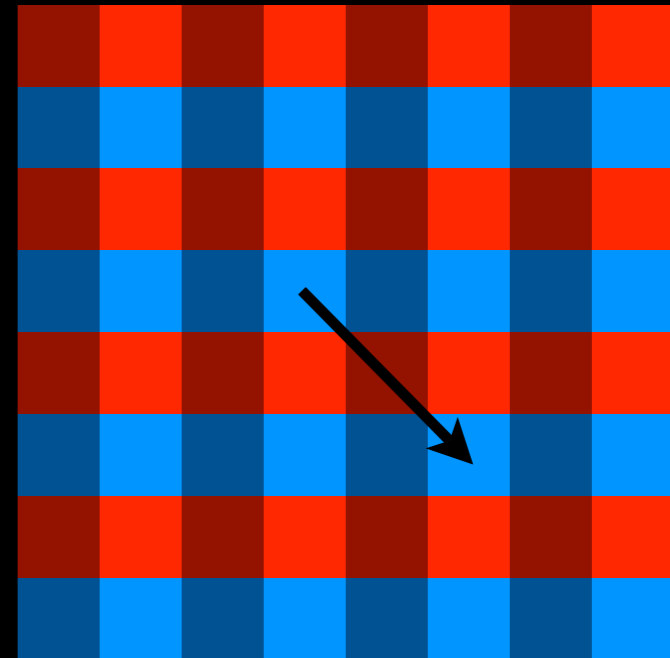
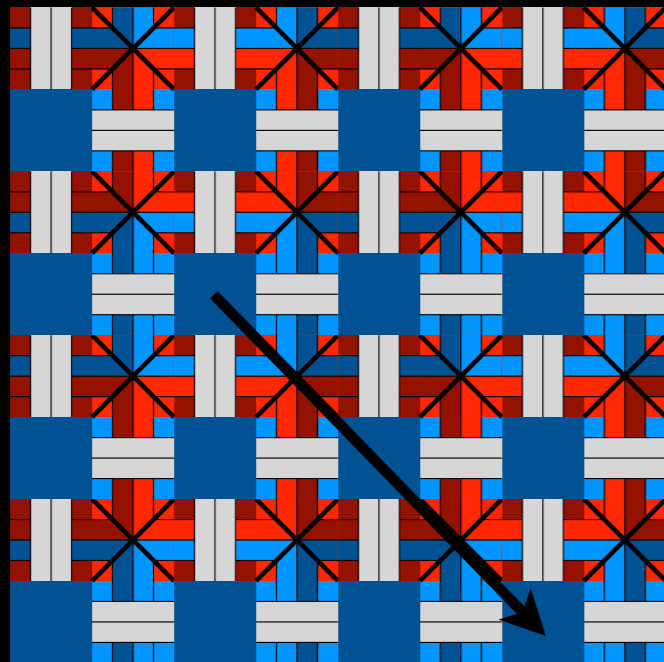
we just remove wires and squarifiers that propagate



Layers 2 & 3

Lemma 2. the preimage of each tiling is a (valid) tiling.



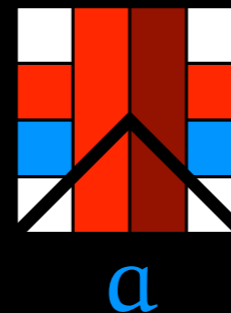
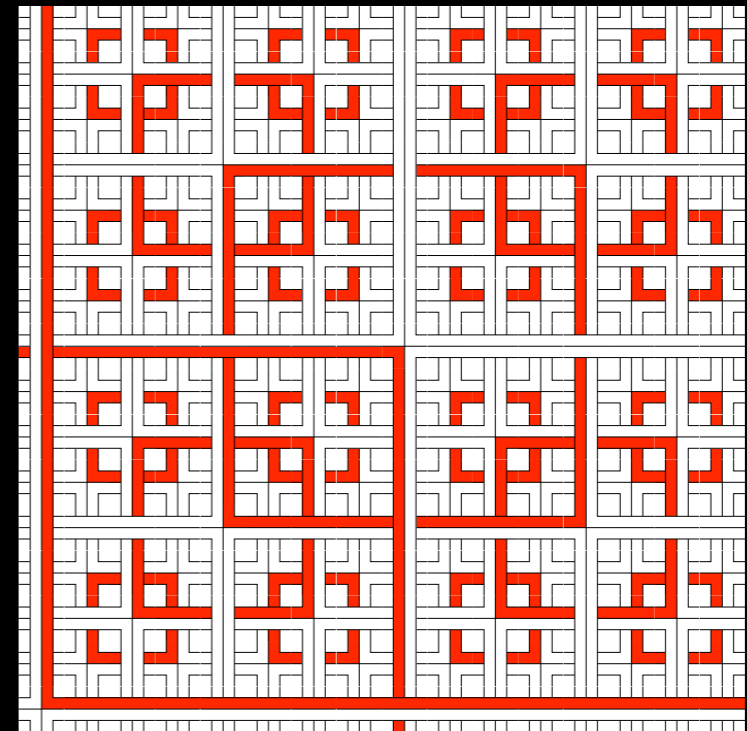


Lemma 3. a periodic tiling of period  $z$  has a preimage of period  $z/2$ .

Back to the Domino Problem

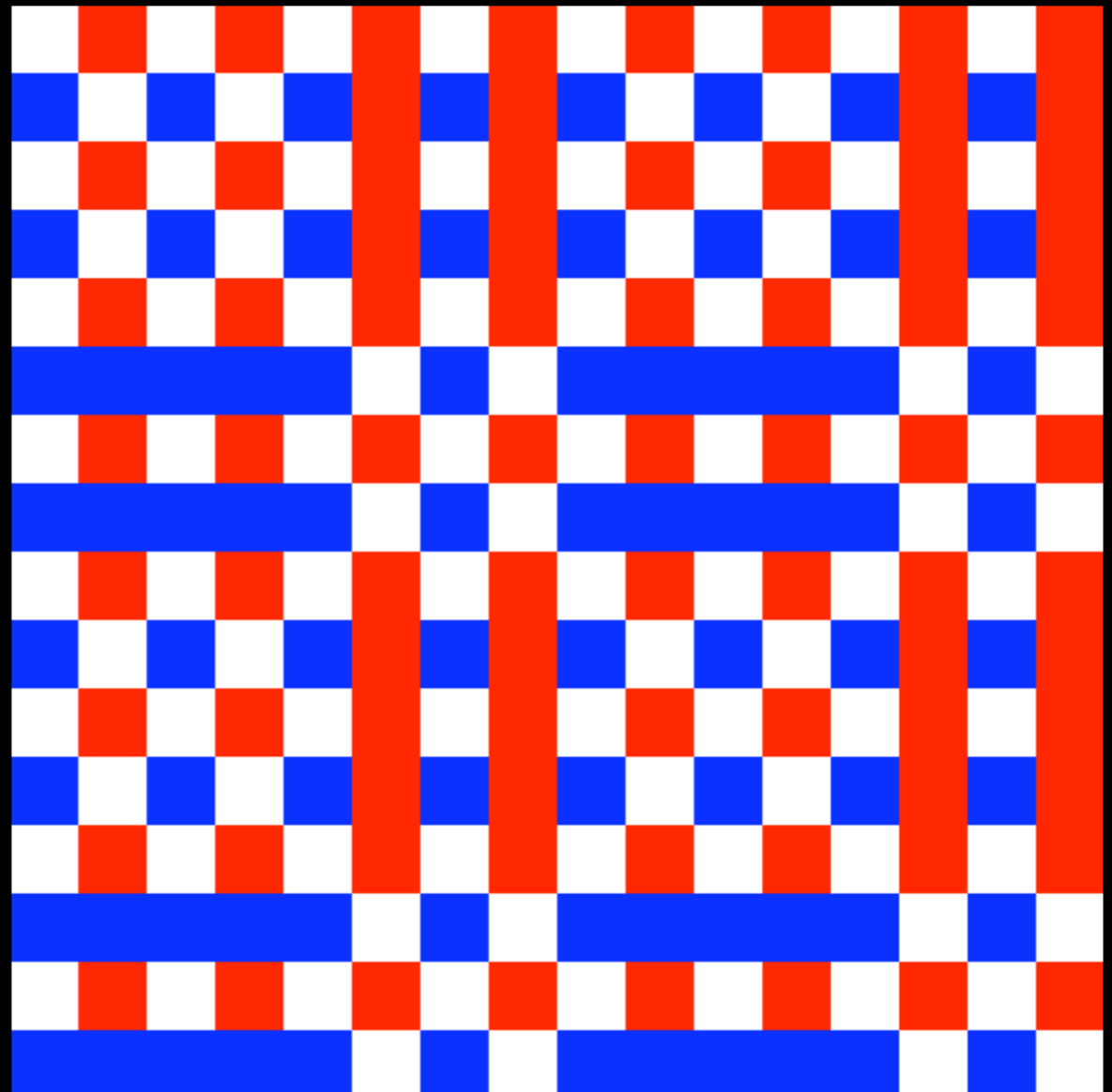
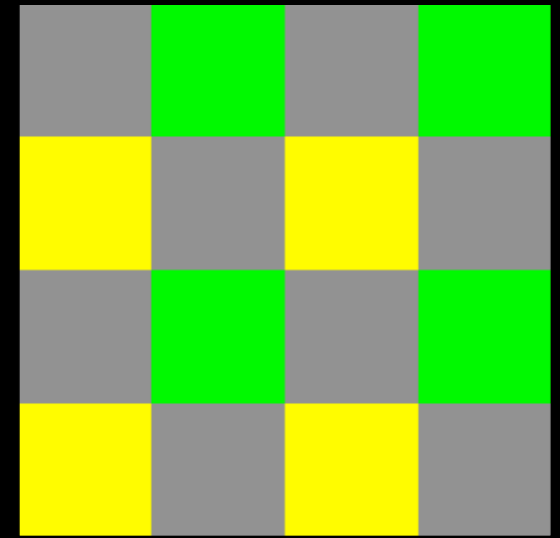
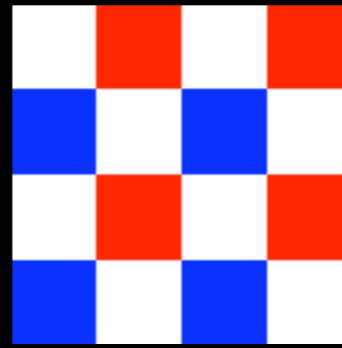
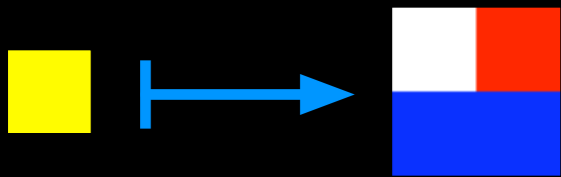
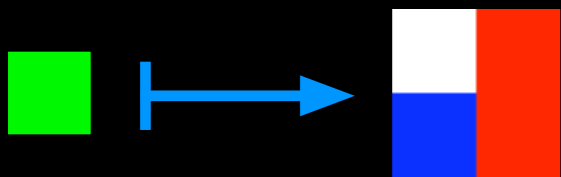
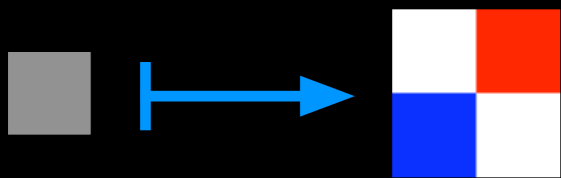
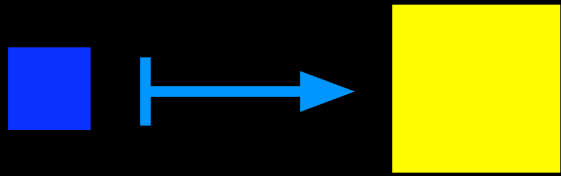
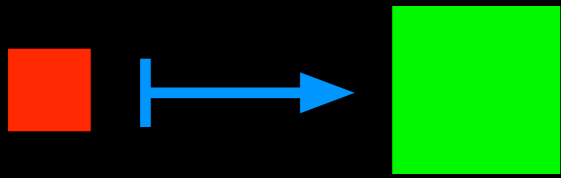
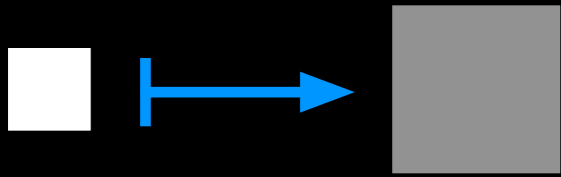
# Encode substitutions

- Select one color of wires (say red): the sequence of squares encode a tree.
- Put a letter on each square.
- Apply the substitution rule to go from one square to the next one.
- Up to projection the substitution limit set is encoded in the set of tilings.

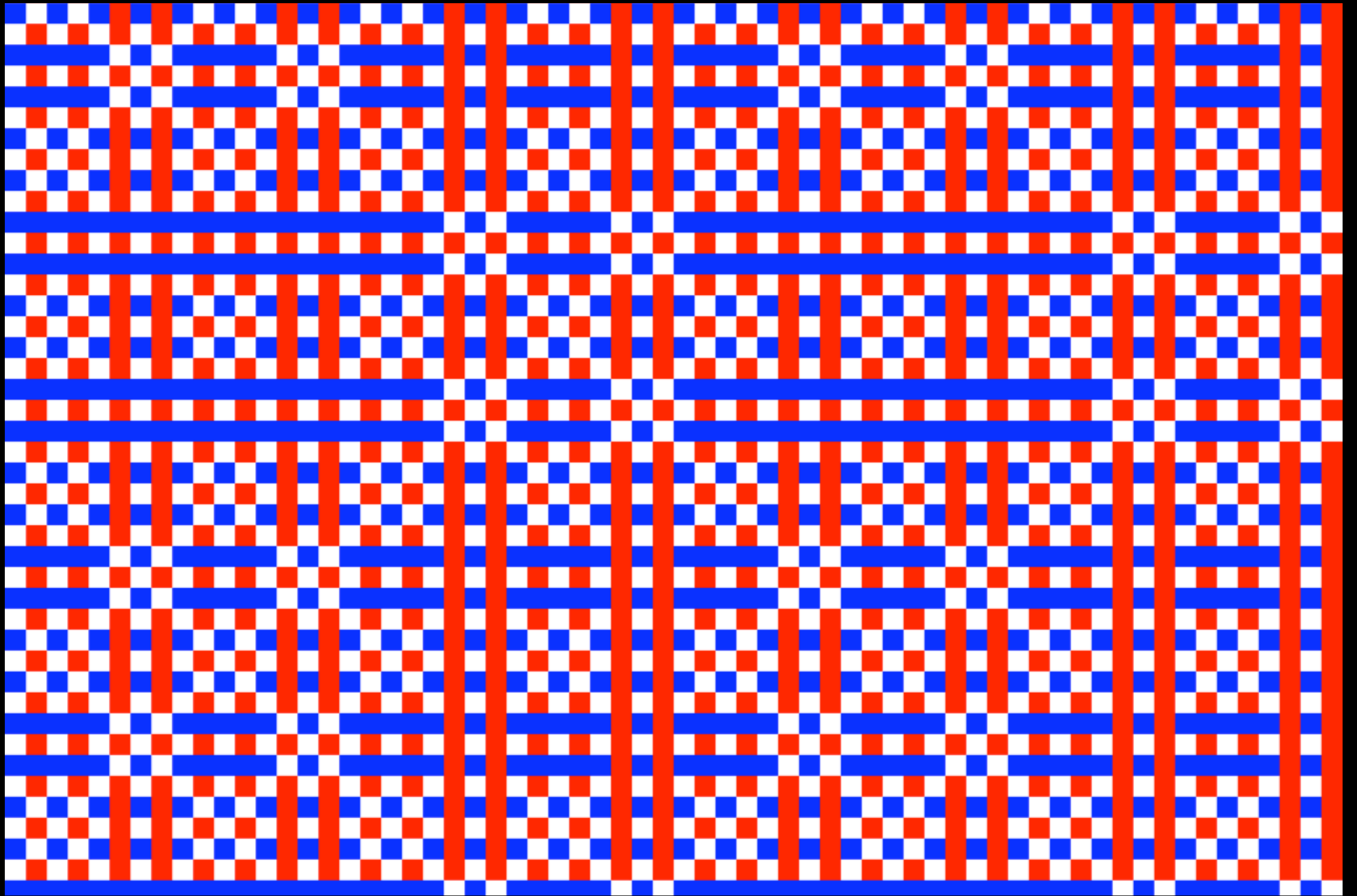


$$b = s(a) \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

$a$



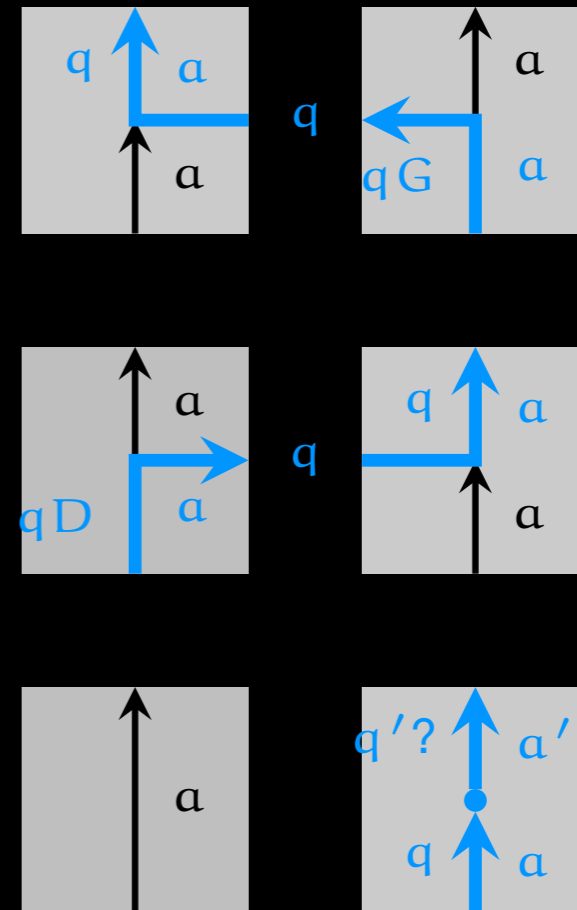
The O2 substitution



O2 iteration

# Turing computation

- Turing machines can be simulated with tilesets.
- Put an initial computation at the SW corner of each  $O_2$  computation square.
- Remove the halting state.
- The tileset tiles the plane iff the Turing machine does not halt.



That's all folks!