Substitutions combinatoires et pavages

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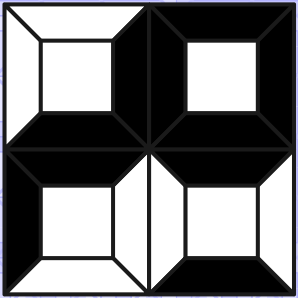
1. Sofic Tilings

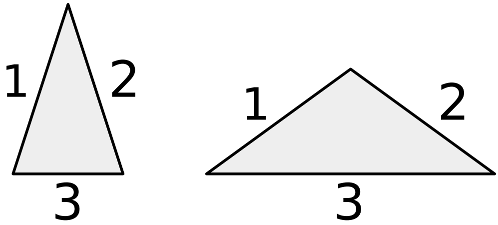
2. Combinatorial Substitutions

3. Main result

4. Conclusion

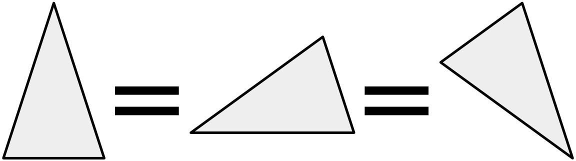


Tiles and tilings



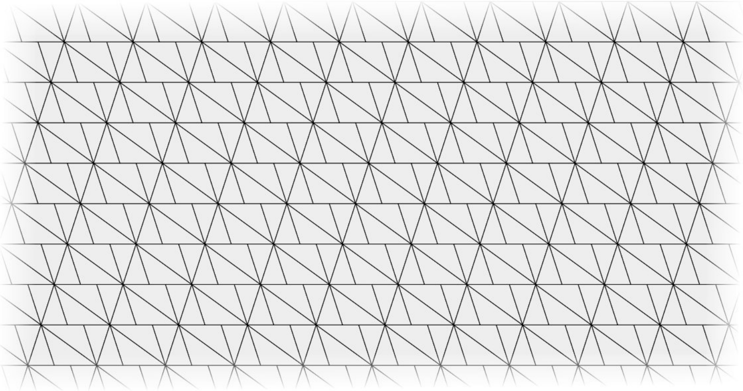
Tile polytope of $\mathbb{R}^d$ with finitely many (numbered) facets.

Tiles and tilings



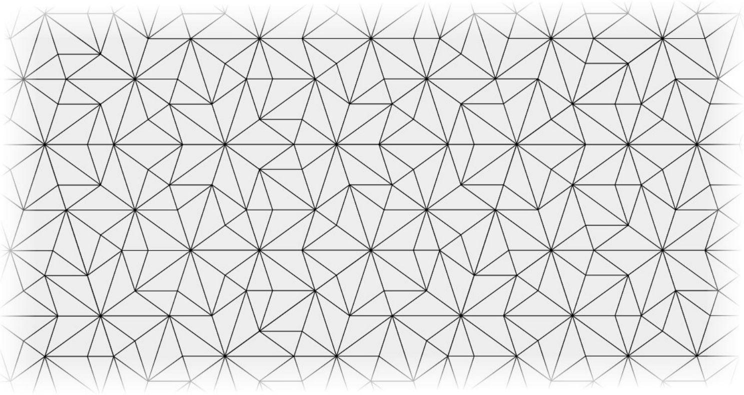
Tiles are here considered up to translations and rotations.

Tiles and tilings



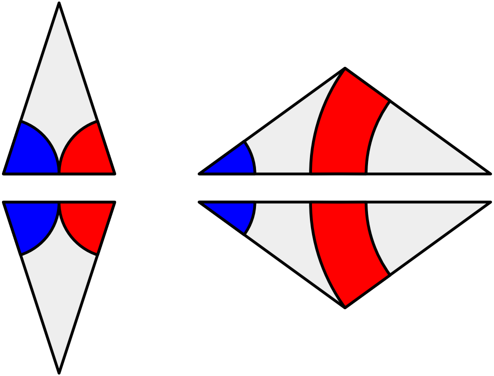
Tiling covering of $\mathbb{R}^d$ by **facet-to-facet** tiles.

Tiles and tilings



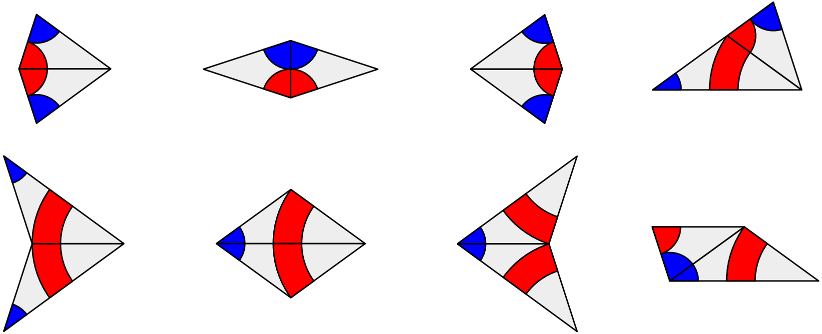
Tiling covering of $\mathbb{R}^d$ by **facet-to-facet** tiles.

Decorations



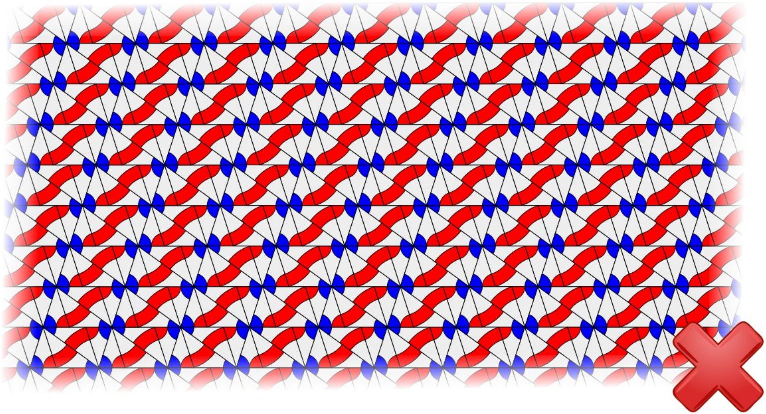
Decoration maps each point of tile boundaries to a color.

Decorations



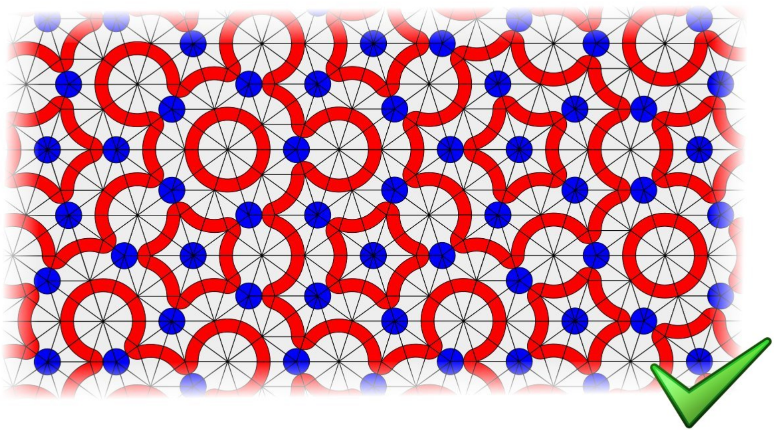
matching if decorations are equal over common facets.

Decorations



Decorated tiling tiling by matching decorated tiles.

Decorations



Decorated tiling tiling by matching decorated tiles.

Sofic tilings

Decorated tile set $\tau \rightsquigarrow$ set Λ_τ of decorated tilings.

Let π be the map which removes tile decorations.

Definition A set of tilings is **sofic** if it can be written as $\pi(\Lambda_\tau)$, where τ is a **finite** decorated tile set.

Sofic tilings

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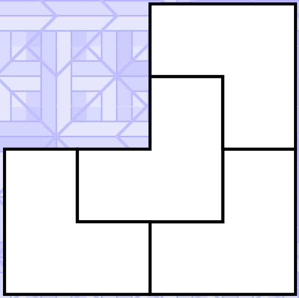
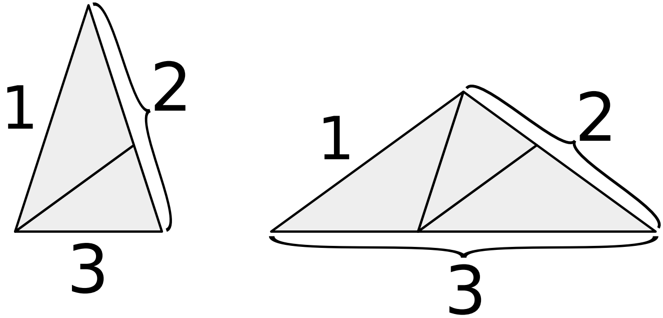
What (interesting) properties on tilings can (or cannot) be enforced by soficity?

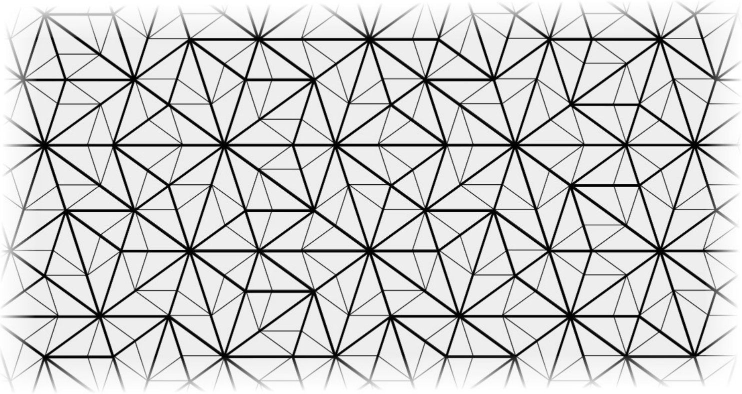


Macro-tiles and macro-tilings



Macro-tile finite partial tiling with (numbered) **macro-facets**.

Macro-tiles and macro-tilings



Macro-tiling macro-facet-to-facet tiling by macro-tiles.

Combinatorial substitution

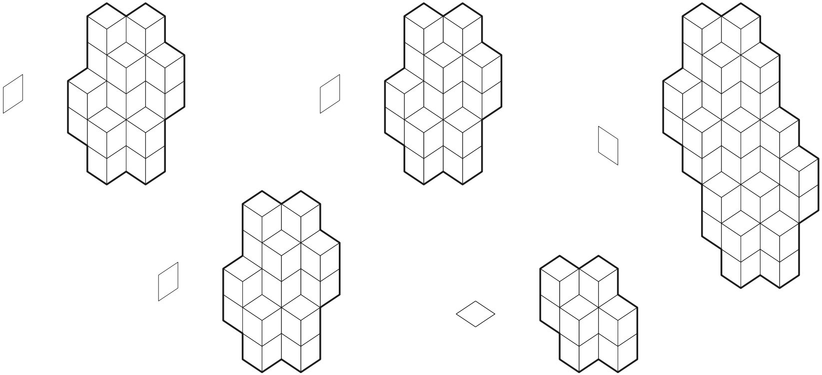
Definition A **combinatorial substitution** is a finite set of pairs (tile, macro-tile).

Let $\sigma = \{(P_i, Q_i)\}_i$ be a combinatorial substitution.

Preimage under σ of a tiling by the P_i's: macro-tiling by the Q_i's with the same **combinatorial structure**.

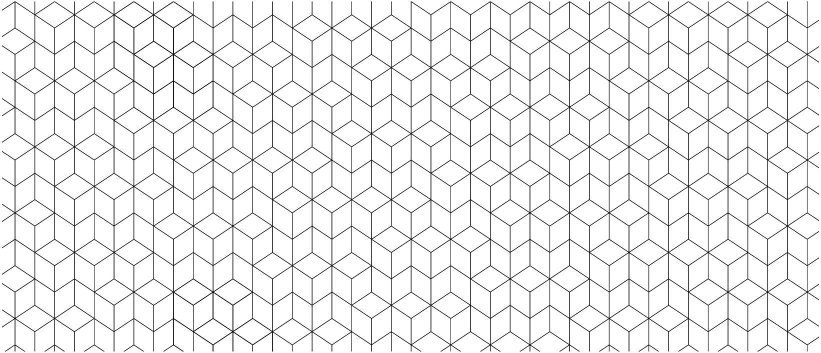
Definition The **limit set** of a combinatorial substitution σ is the set of tilings which admit an infinite sequence of preimages under σ.

Example



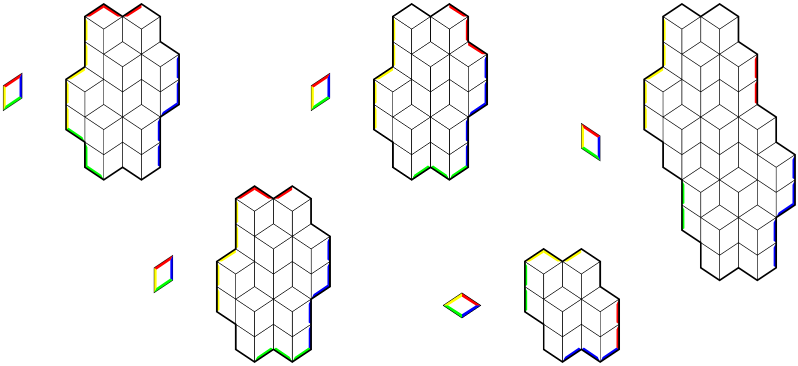
From Rauzy generalized substitution...

Example



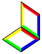
From Rauzy generalized substitution. . .

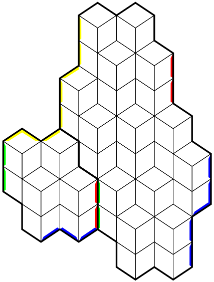
Example



. . . to Rauzy combinatorial substitution.

Example

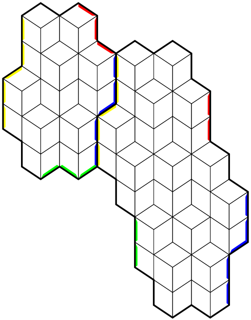




Tiles match in a tiling as macro-tiles in its image...

...and conversely.

Example



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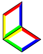

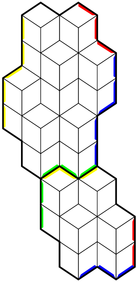

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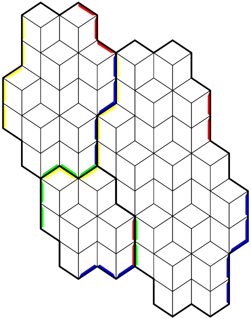

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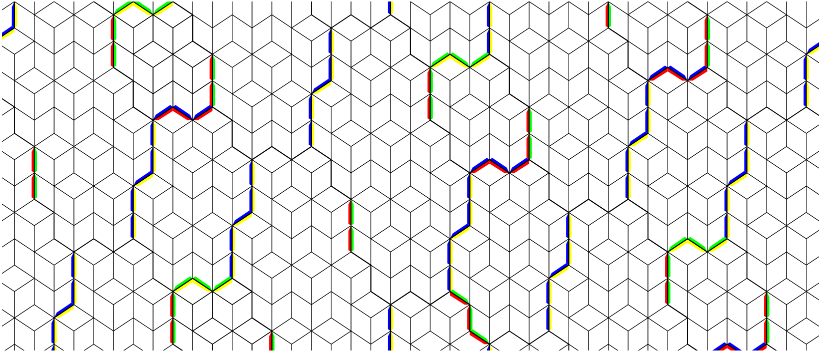
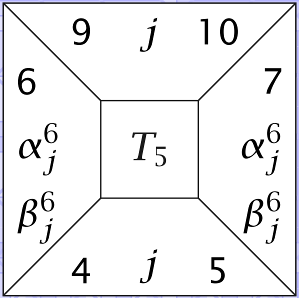
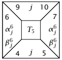


Any tiling decomposes into macro-tiles.


9
j
10
6
7
α_j^6
T_5
α_j^6
β_j^6
β_j^6
4
j
5

Main result

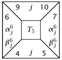


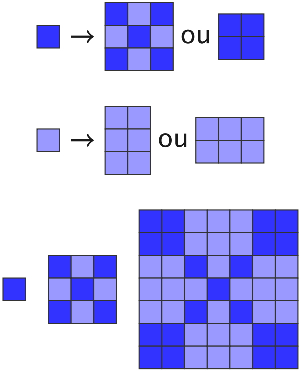
Theorem[FO 2010] The **limit set** of a **good** combinatorial substitution is **sofic**.

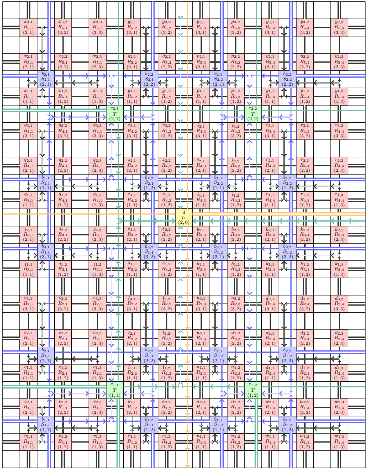
Remark The result is **constructive**: given a substitution we recursively construct a decorated set of tiles.

This extends (and simplifies?) previous similar results.

Mozes 1990

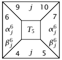


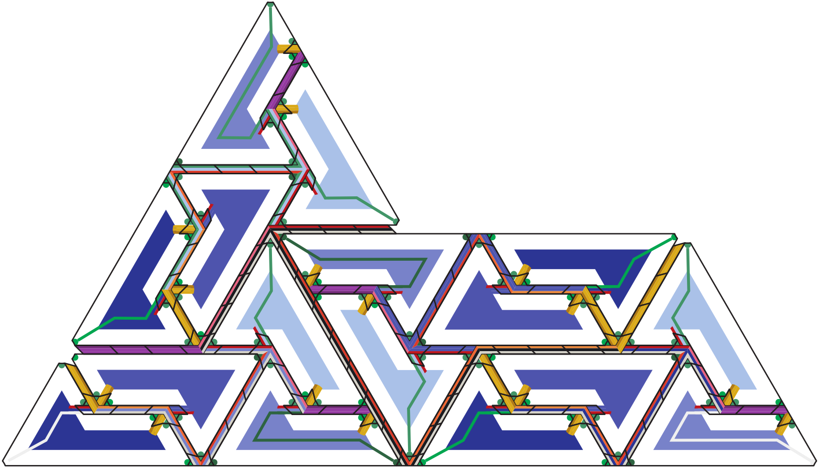




Theorem[Mozes 1990] The limit set of a **non-deterministic rectangular substitution** is sofic.

Goodman-Strauss 1998



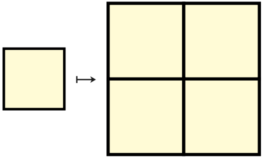


Theorem[Goodman-Strauss 1998] The limit set of **homothetic substitution** *(+ some hypothesis)* is sofic.

Recipe [O 2008]

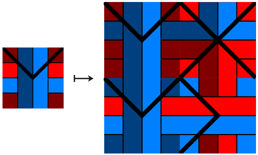
A **substitution** s generates a **limit set** $\Lambda_s = \bigcap_t \mathrm{Img}(s^t)$.

The limit set is the set of colorings admitting an **history** $(c_i)_{i\in\mathbb{N}}$ where $c_i = s(c_{i+1})$.

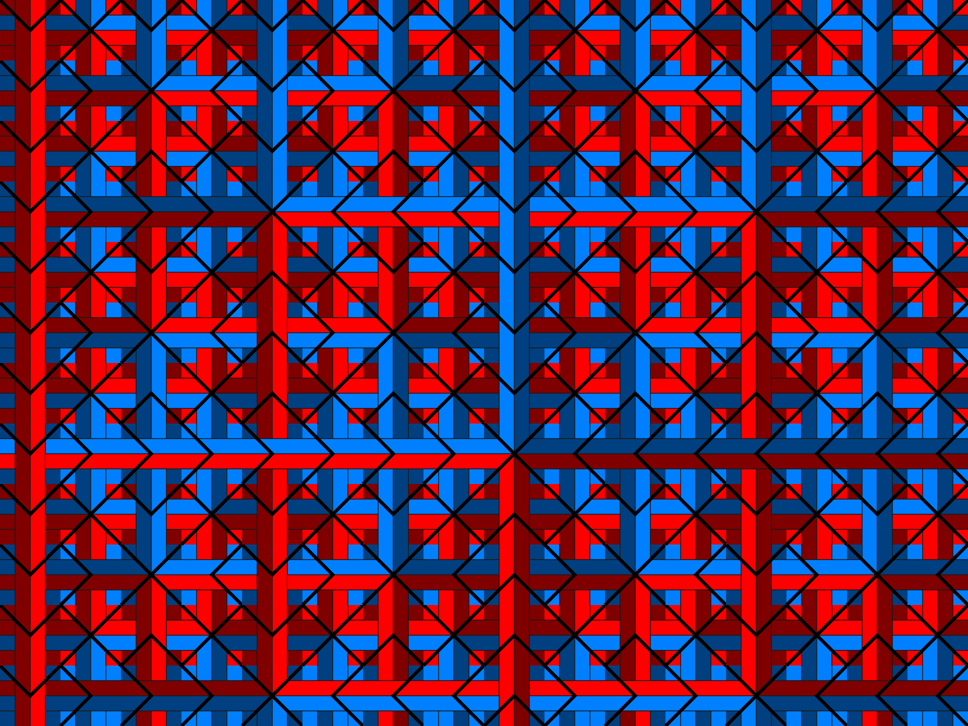


A tile set τ **simulates** a tile set τ' with an encoding $f : \tau' \to \tau^n$ if tilings by τ decompose via f in tilings by τ'.

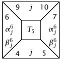
Tilings of a **self-simulating** tile set τ with encoding s are the **limit set** of s.



To encode Λ_s via **local matching rules** decorate s into a **locally checkable** $s^\bullet$ embedding a whole history.



Self-simulation

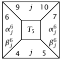


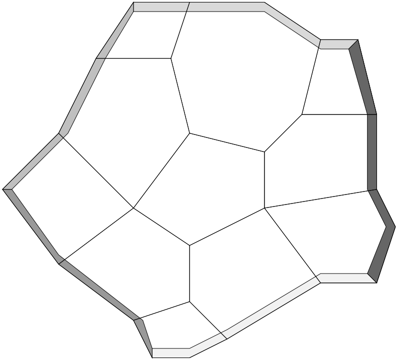
Definition A decorated tile set τ **self-simulates** if it admits tilings and there are τ-macro-tiles s.t.

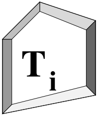
1. any τ-tiling is also a macro-tiling by these τ-macro-tiles;
2. each τ-macro-tile is **combinatorially equivalent** to a τ-tile.

Proposition If τ **self-simulates**, then $\pi(X_\tau)$ is a **subset of the limit set** of the combinatorial substitution with pairs τ-macro-tile/equivalent τ-tile.

A self-simulating decorated tile set τ

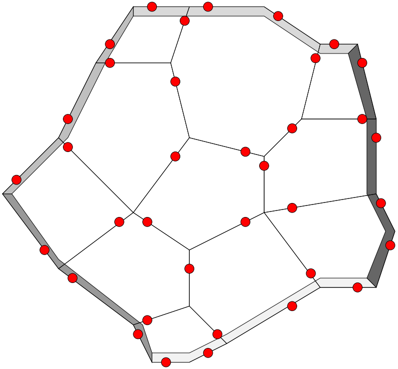


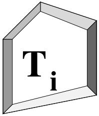




Fix a set of macro-tiles and let $T_1, \ldots, T_n$ be all their tiles.

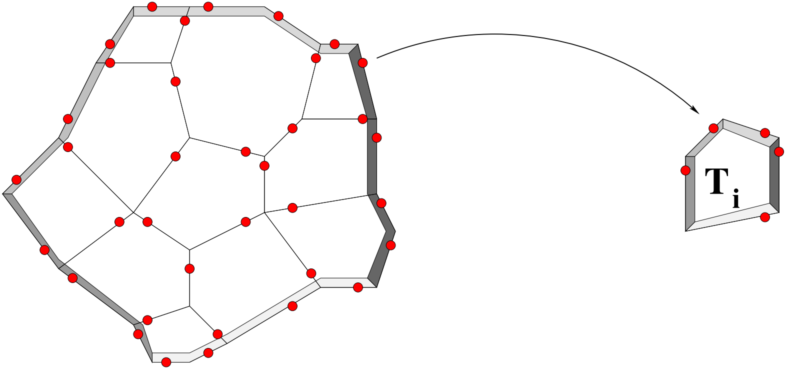
A self-simulating decorated tile set τ





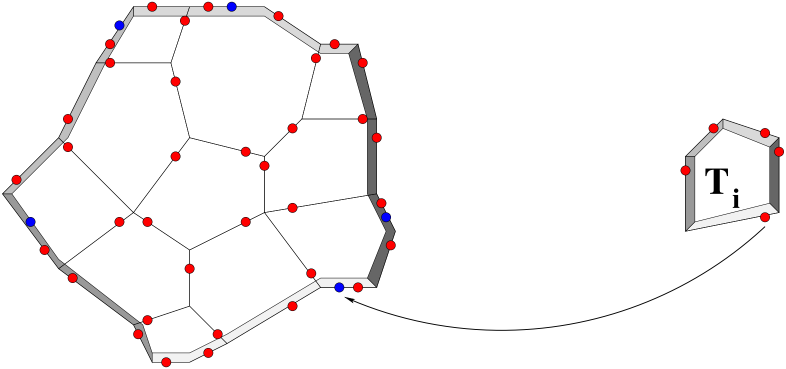
To enforce τ-tilings to be τ-macro-tilings: decorations specify tile neighbors within macro-tiles and mark macro-facets.

A self-simulating decorated tile set τ



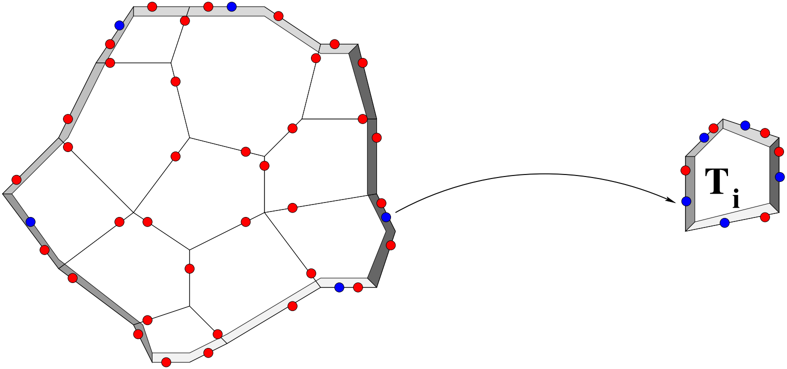
This yields so-called macro-indices on tile facets.

A self-simulating decorated tile set τ



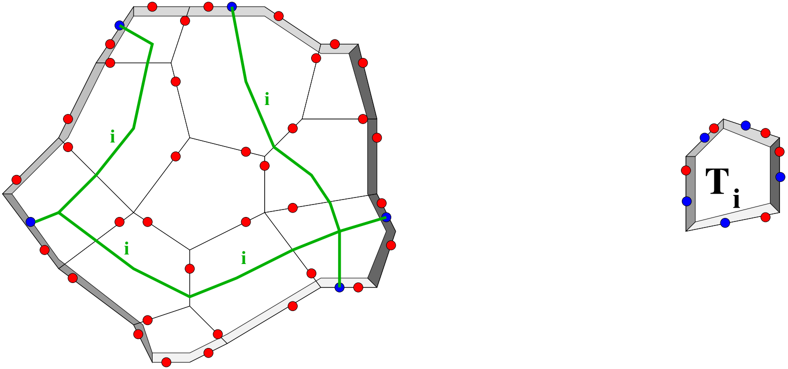
The macro-indices of facets of a τ-tile must then be encoded on the corresponding macro-facets of its simulating τ-macro-tile.

A self-simulating decorated tile set τ



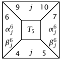
This yields so-called neighbor-indices on tile facets.

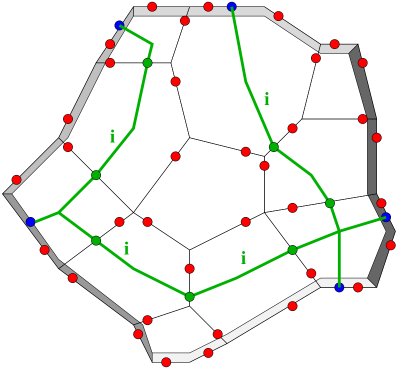
A self-simulating decorated tile set τ

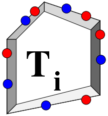


We force these neighbor-indices to come from the same tile T_i, called parent-tile, by carrying its index i between macro-facets, where it is converted into the suitable neighbor-index.

A self-simulating decorated tile set τ

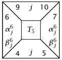


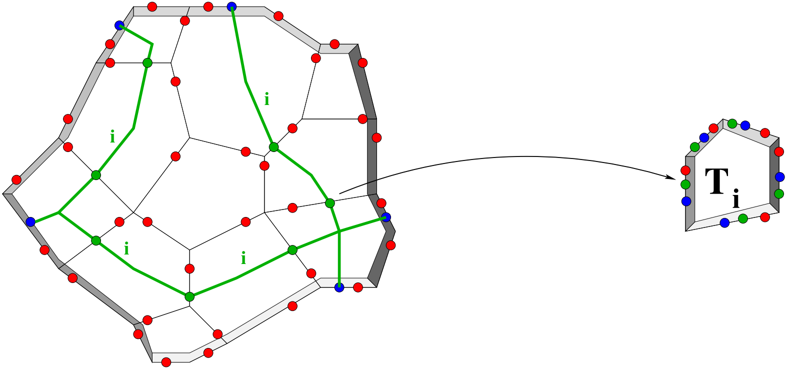




Such tile indices are encoded on facets by so-called parent-index.

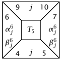
A self-simulating decorated tile set τ

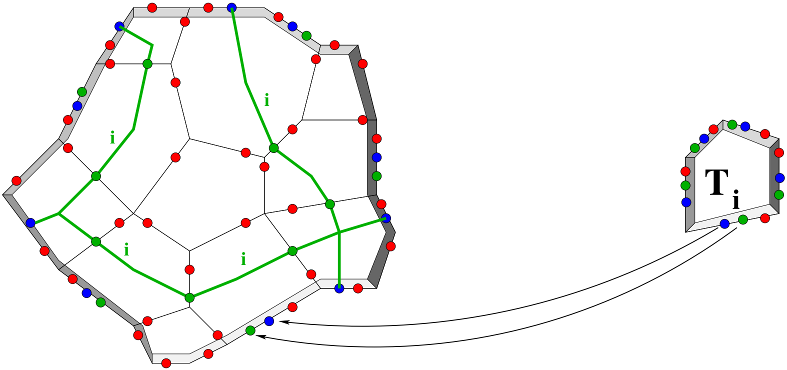




This yields, once again, a new index on each tile facets...

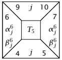
A self-simulating decorated tile set τ

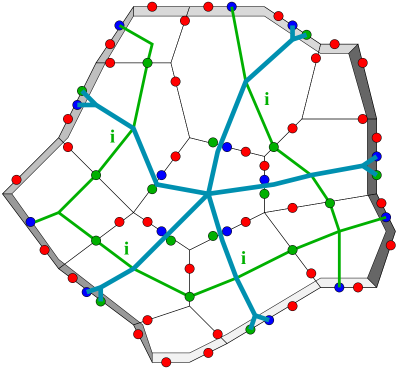


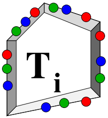


But the trick is that the neighbor-indices and parent-indices of facets of a τ-tile can be encoded on the corresponding <u>big enough</u> macro-facets of the equivalent τ-macro-tile <u>without any new index</u>!

A self-simulating decorated tile set τ

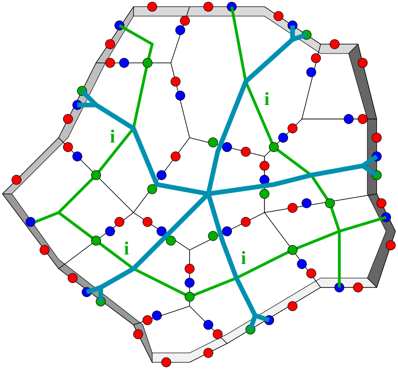


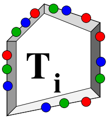




In big enough macro-tiles, we can then carry these pairs of neighbor/parent indices up to a central tile along a star-like network.

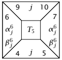
A self-simulating decorated tile set τ

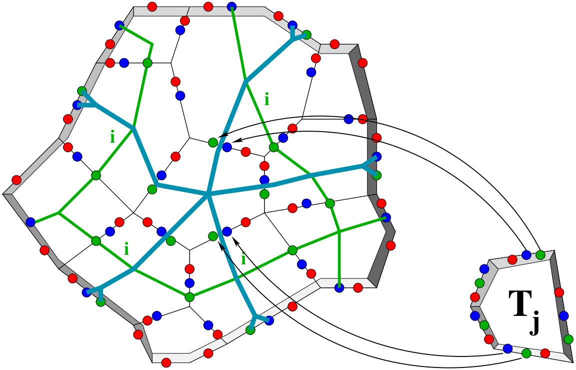




On internal facets not crossed by this network, we copy the macro-index on the neighbor-index (this redundancy is later used).

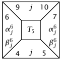
A self-simulating decorated tile set τ

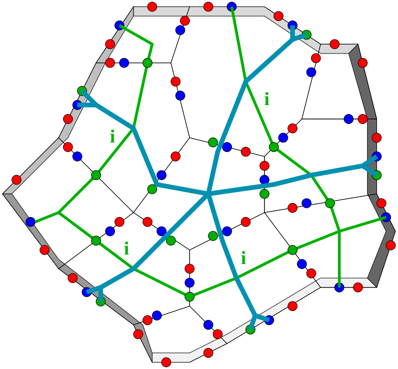




The pairs on a central τ-tile can be those of any non-central τ-tile (from which the central τ-tile is said to derive).

A self-simulating decorated tile set τ





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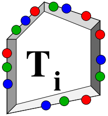
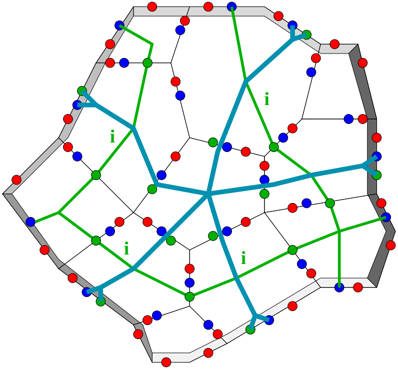
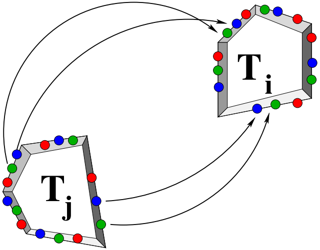


The τ-macro-tile with parent-index i is combinatorially equivalent to T_i endowed with the pairs of the central τ-tile.
<u>But is it a τ-tile?</u>

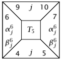
A self-simulating decorated tile set τ

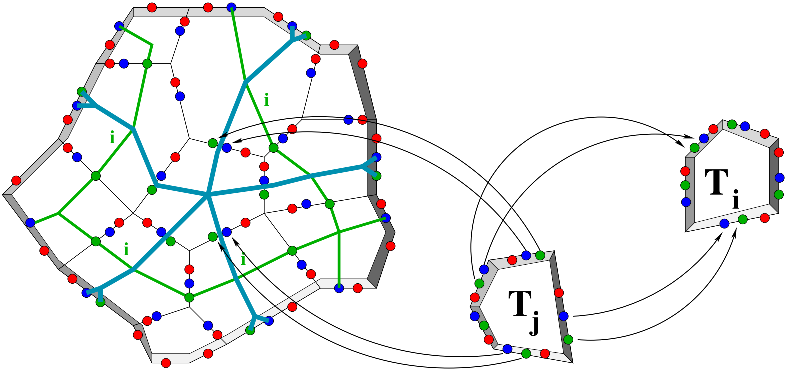




If T_i is a central tile, then its pairs can be derived from any non-central τ-tile (as for any central tile)...

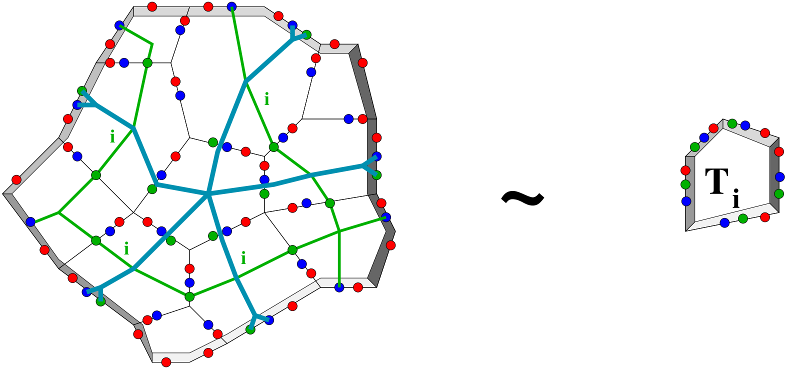
A self-simulating decorated tile set τ





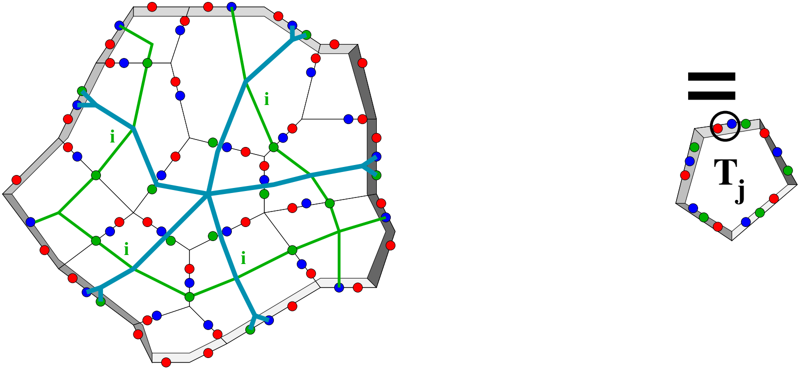
...in particular from the non-central τ-tile from which are also derived the pairs of the central τ-tile of our τ-macro-tile.

A self-simulating decorated tile set τ



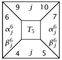
In this case, the equivalent decorated T_i is a derived central τ-tile.

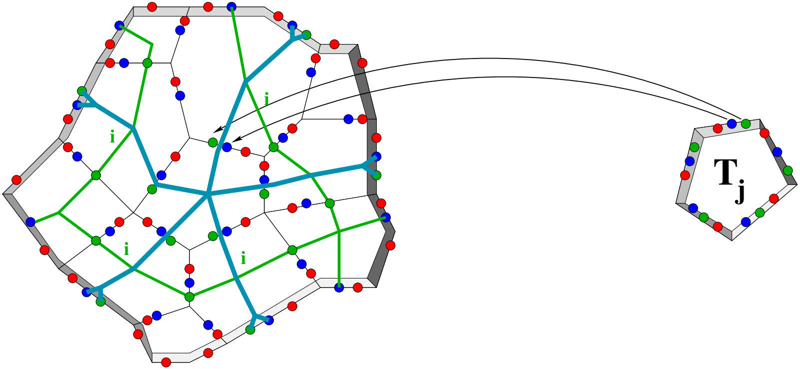
A self-simulating decorated tile set τ



Otherwise, consider the non-central τ-tile from which derives our central τ-tile; at least one facet is internal and not crossed by a network: its neighbor and macro indices are equal (by redundancy).

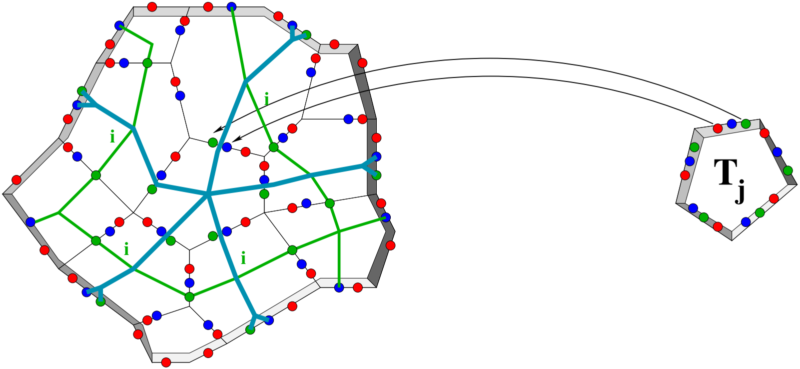
A self-simulating decorated tile set τ





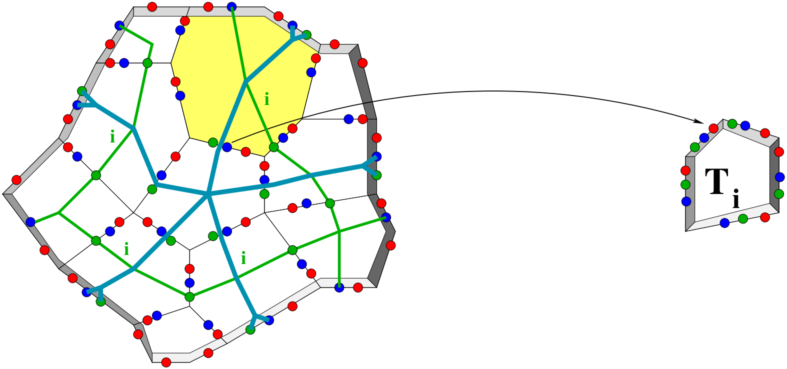
Thus, by copying the neighbor and parent indices (derivation)...

A self-simulating decorated tile set τ



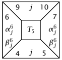
. . . one copies a macro-index on our central τ-tile, and thus on the whole corresponding network branch.

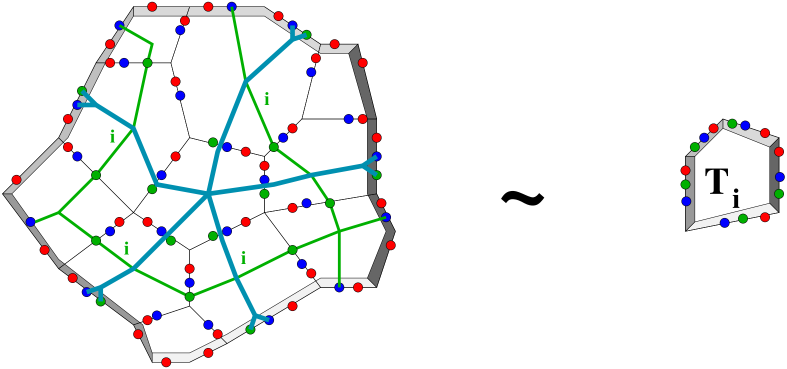
A self-simulating decorated tile set τ



A tile on this k-th branch which also knows the parent-index i can then force this macro-index to be the one on the k-th facet of a decorated T_i (recall that all the decorated T_i have the same one).

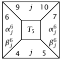
A self-simulating decorated tile set τ





In this case, the equivalent decorated T_i is the non-central τ-tile from which derives the central τ-tile of our τ-macro-tile.

Main result



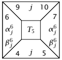
Theorem[FO 2010] The **limit set** of a **good** combinatorial substitution is **sofic**.

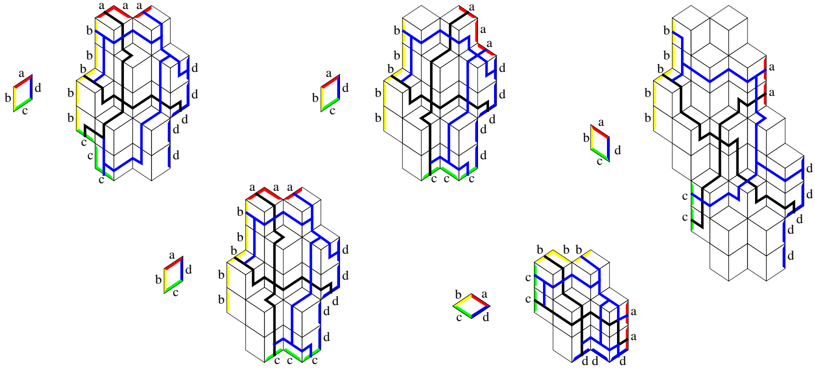
Remark No need to care about **geometry**.

But, what is a **good** combinatorial substitution?

Definition A **good** combinatorial substitution is both **connecting** and **consistent**.

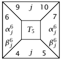
Connecting





Intuitively A substitution is **connecting** if there is enough space inside macro-tiles to wire the networks.

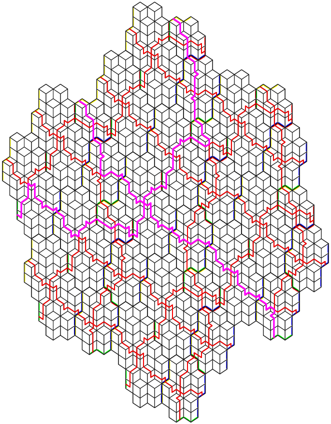
Consistent



Definition A combinatorial substitution is **consistent** if any tiling by macro-tiles admits a preimage under the substitution.

Remark This is where the **geometrical** consistency hides.

Open Pb **Characterize** consistent combinatorial substitutions.



1. Sofic Tilings

2. Combinatorial Substitutions

3. Main result

4. Conclusion

Conclusion and open problem

???

The **global** hierarchical structure associated to a substitution system can be enforced by **local** matching rules.

Open Pb Is it possible to describe the **geometry** of tiles by finite local **combinatorial** constraints?