

Indécidabilité, pavages et polyominos

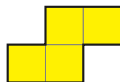
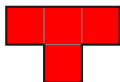
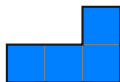
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Polyominoes

Definition A **polyomino** is a simply connected tile obtained by gluing together rookwise connected unit squares.



Definition A **tiling** of a region by a set of polyominoes is a partition of the region into images of the tiles by isometries.



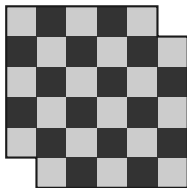
Definition A **tiling by translation** is a tiling where isometries are restricted to translations.

Tiling finite regions

Remark The combinatorics of tilings of finite regions is challenging, polyominoes make great puzzles.

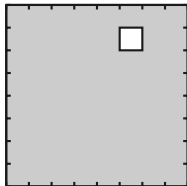
Can you tile with dominoes a $2m \times 2n$ rectangle with two opposite corners cut?

[Golomb 1965]



Can you tile with L-tiles a $2^n \times 2^n$ square with one cut unit square?

[Golomb 1965]



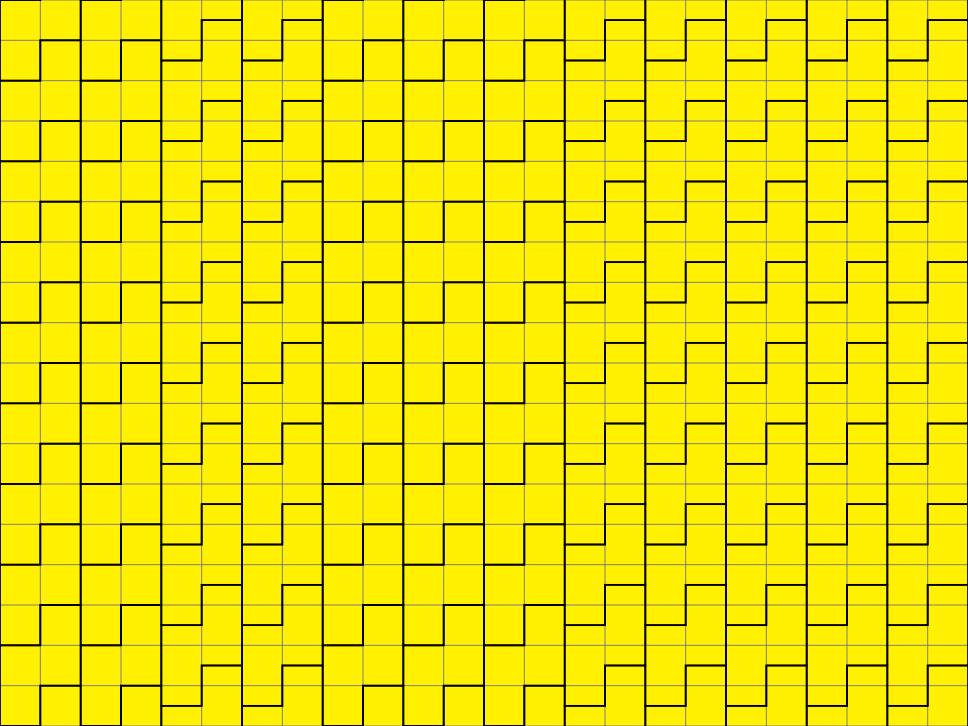
Tiling the plane

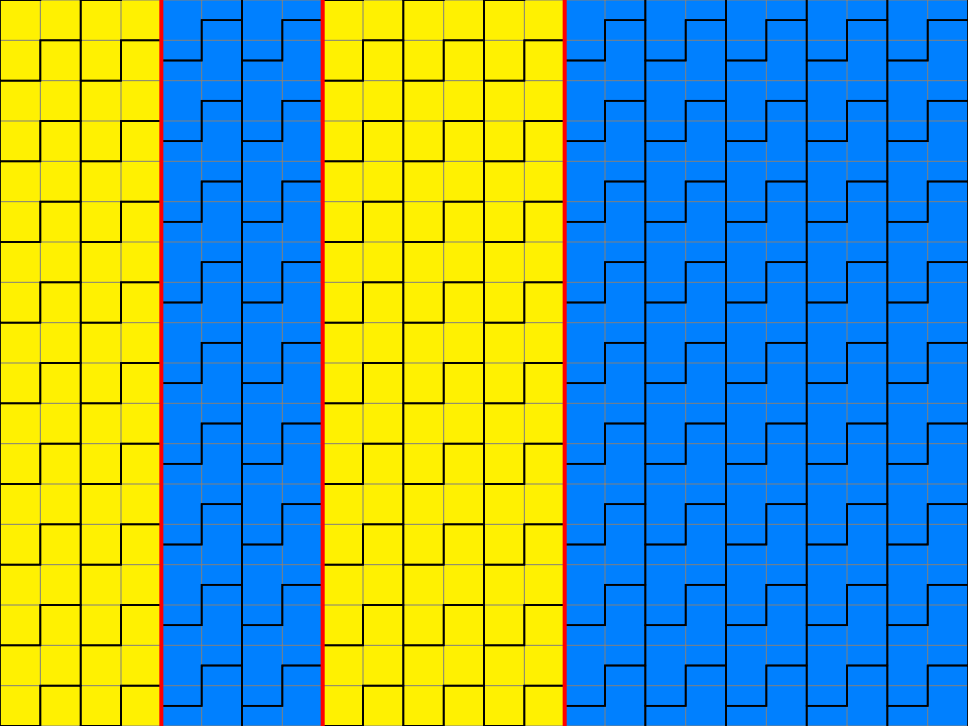
We consider tilings of **the whole Euclidian plane** by finite sets of polyominoes.

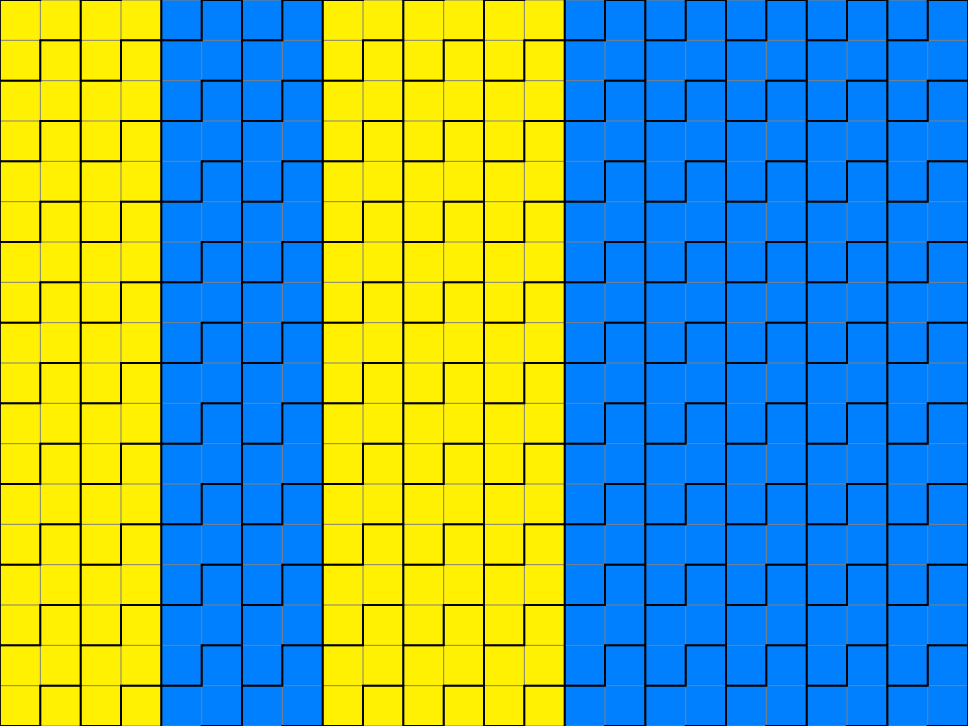
Definition A tiling is **discrete** if all the unit squares composing images of the polyominoes are aligned on the grid \mathbb{Z}^2 .

Lemma A tile set admits a tiling iff it admits a discrete tiling.

Sketch of the proof Non-discrete tilings have countably many infinite parallel fracture lines. By shifting along fracture lines, one constructs a discrete tiling from any non-discrete tiling.







The k -Polyomino Problem

We consider the **complexity** of the following two problems:

Polyomino Problem Given a finite set of polyominoes, decide if it can tile the plane.

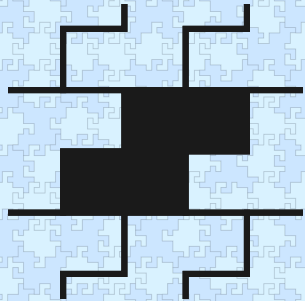
k -Polyomino Problem Given a set of k polyominoes, decide if it can tile the plane.

1. Complexity of tiling

2. The Polyomino Problem

3. The k -Polyomino Problem

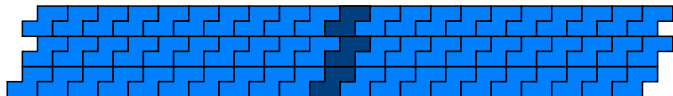
4. Conclusion



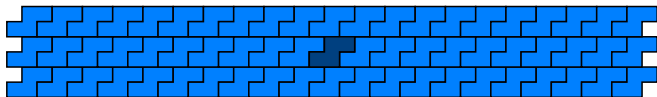
Periodic Tilings



Definition A tiling is **periodic** with period p if it is invariant by a **translation** of vector p .



Lemma If a finite set of polyominoes admits a **periodic** tiling then it admits a **biperiodic** tiling.

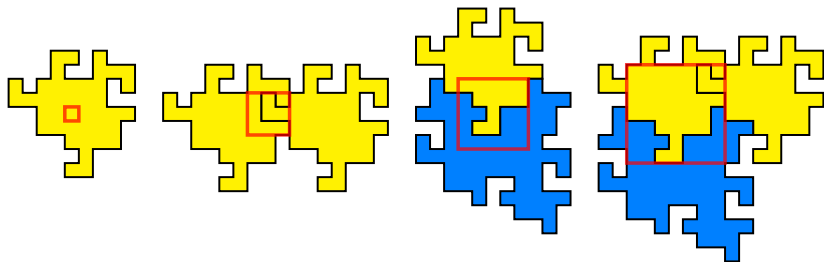


Lemma Finite sets of polyominoes tiling the plane biperiodically are **re (recursively enumerable)**.



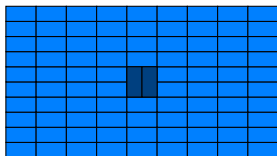
Lemma Finite sets of polyominoes tiling the plane are **co-re**.

Sketch of the proof Consider tilings of finite regions covering larger and larger squares. If the set does not tile the plane, by compactness, there exists a size of square it cannot cover with tiles.





Definition A tiling is **aperiodic** if it admits no non-trivial period.



Definition A set of polyominoes is **aperiodic** if it admits a tiling and all its tilings are aperiodic.

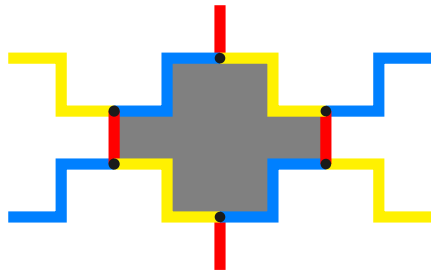
Remark When there is **no aperiodic** set of tiles, the $(k-)$ Polyomino Problem is **decidable**.

One polyomino by translation



Theorem[Wijshoff and van Leeuwen 1984] A **single polyomino** that tiles the plane **by translation** tiles it biperiodically. The problem is decidable.

Theorem[Beauquier and Nivat 1991] A single polyomino tiles the plane by translation iff it is a **pseudo-hexagon**.



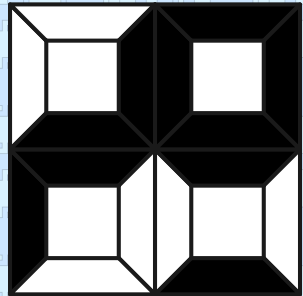
[Gambini and Vuillon 2007] This can be tested in $O(n^2)$.

1. Complexity of tiling

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4. Conclusion

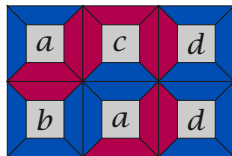
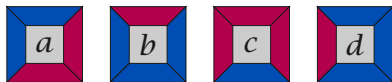


The Domino Problem (DP)

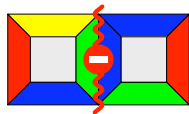
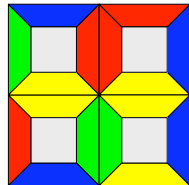
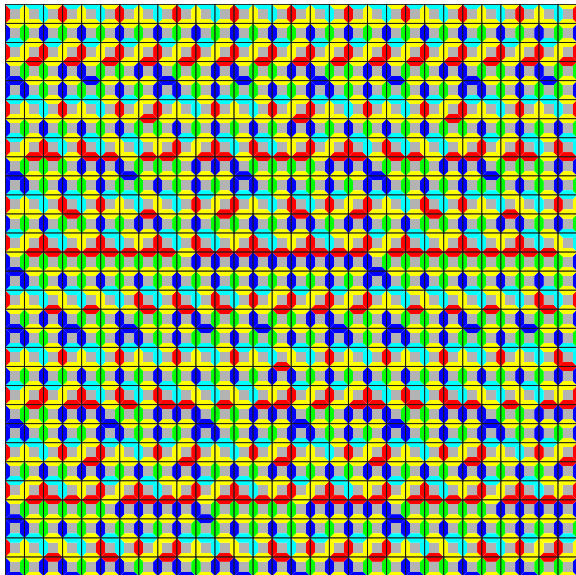


“Assume we are *given a finite set of square plates of the same size with edges colored, each in a different manner. Suppose further there are infinitely many copies of each plate (plate type). We are **not permitted to rotate or reflect a plate.** The question is to find an effective procedure by which we can **decide**, for each given finite set of plates, **whether we can cover up the whole plane** (or, equivalently, an infinite quadrant thereof) **with copies of the plates subject to the restriction that adjoining edges must have the same color.**”*

(Wang, 1961)



Wang tiles



Undecidability of DP



Theorem[Berger 1964] **DP** is **recursively undecidable**.

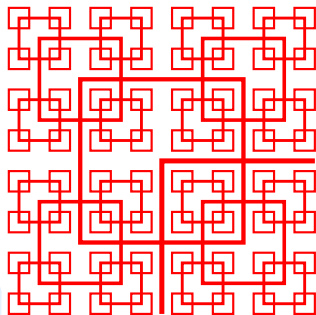
Remark To prove it one needs **aperiodic** tile sets.

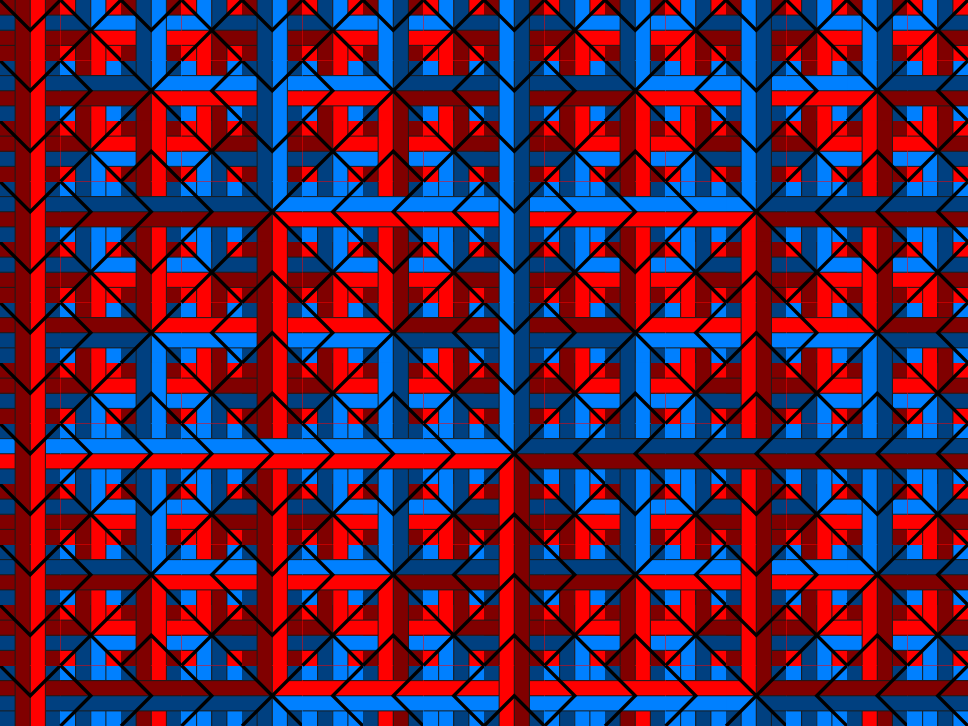
Idea of the proof

Enforce an (aperiodic) **self-similar structure** using local rules.

Insert a **Turing machine computation everywhere** using the structure.

Remark Plenty of different proofs!





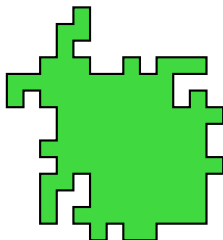
The Polyomino Problem is undecidable



Wang tiles are oriented unit squares with colors.

Colors can be encoded by **bumps and dents**.

A Wang tile can be **encoded** as a big pseudo-square polyomino with bumps and dents in place of colors.



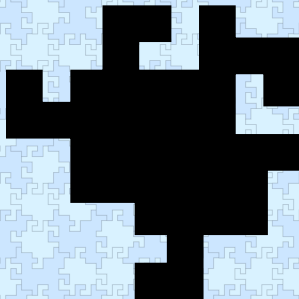
Theorem[Golomb 1970] The Polyomino Problem is **recursively undecidable**.

1. Complexity of tiling

2. The Polyomino Problem

3. The k -Polyomino Problem

4. Conclusion





Remark The reduction of Golomb encodes N Wang tiles into N polyominoes.

What about the **k -Polyomino Problem**?

(1) either it is decidable for all k and the family of algorithms is not itself recursive (*eg. set of Wang tiles with k colors*);

(2) either there exists a frontier between decidable and undecidable cases (*eg. Post Correspondence Problem*).

We will show that (2) holds.

Dented polyominoes



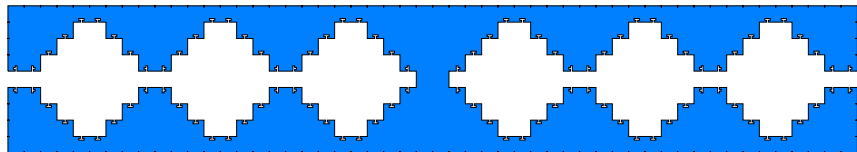
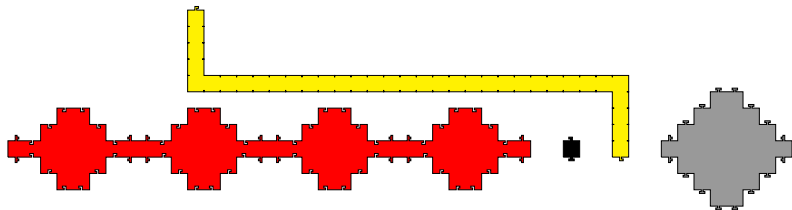
Computing with polyominoes relies on several levels of encoding. To lever the complexity of the tiles, we use dented polyominoes.

Definition A **dented polyomino** is a polyomino with edges labeled by a **dent shape** and an **orientation**. When considering tilings, dents and bumps have to match.

Lemma Every set of k dented polyominoes can be encoded as a set of k polyominoes, preserving the set of tilings.

Sketch of the proof Scale each polyomino by a factor far larger than bumps, then add bumps and dents along edges.

5 tiles

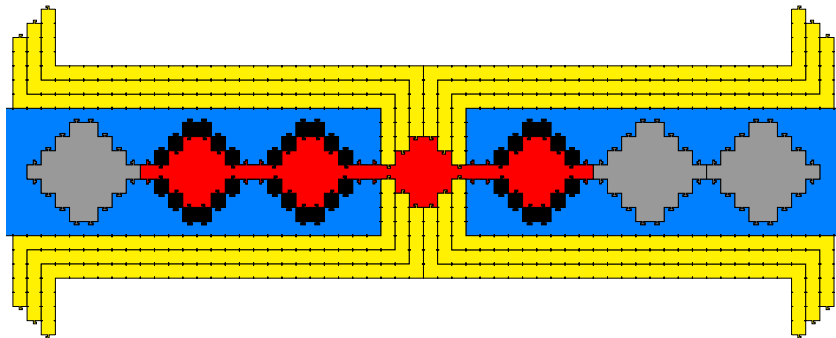


	<i>blank</i>	<i>bit</i>	<i>marker</i>	<i>inside</i>
shape				
bump				
dent				
		wire, tooth	meat, filler	tooth, filler
		meat	jaw	jaw

Encoding Wang tiles



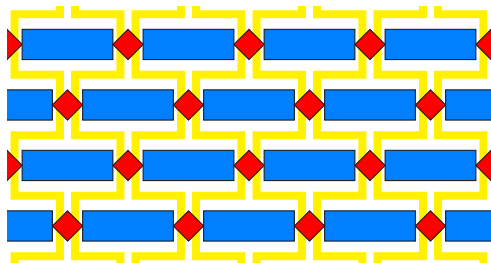
A **meat** is placed in between two **jaws** to select a tile. The gaps inside the **jaws** are filled by **fillers** and **teeth**. **Wires** connect Wang tiles.



Encoding a tiling by Wang tiles



Wang tiles are encoded and placed on a **regular grid**.
Tiles of a same diagonal are placed on a horizontal line
sharing jaws.





It remains to show to **difficult part of the proof**.

Why does every tiling codes a tiling by Wang tiles?

- (1) The polyominoes locally enforce Wang tiles coding;
- (2) Details on the encoding of colors enforce a same orientation for all Wang tiles in the plane.

Theorem[O 2009] The 5-Polyomino Problem is **undecidable**.

Tiling by translation



Previous encoding uses:

1 meat, 1 jaw, 1 filler, 4 wires, 4 teeth.

Theorem[O 2009] The 11-Polyomino Translation Problem is **undecidable**.

The problem is decidable for a single polyomino and undecidable for 11 polyominoes. What about $2 \leq k \leq 10$?

Even for $k = 2$, it seems that it is not trivial...

1. Complexity of tiling
2. The Polyomino Problem
3. The k -Polyomino Problem
- 4. Conclusion**



Patch Extension Problem Given a finite patch of tiles decide if it can be extended to a tiling.

Theorem[Myers 1974] There exists sets of Wang tiles for which the Patch Extension Problem is **undecidable**.

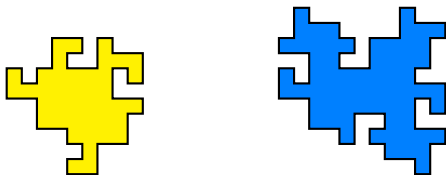
Remark The proof is **constructive** and gives an explicit aperiodic set of tiles.

Corollary There exists sets of 5 polyominoes for which the Patch Extension Problem is **undecidable**.

Aperiodic sets of polyominoes

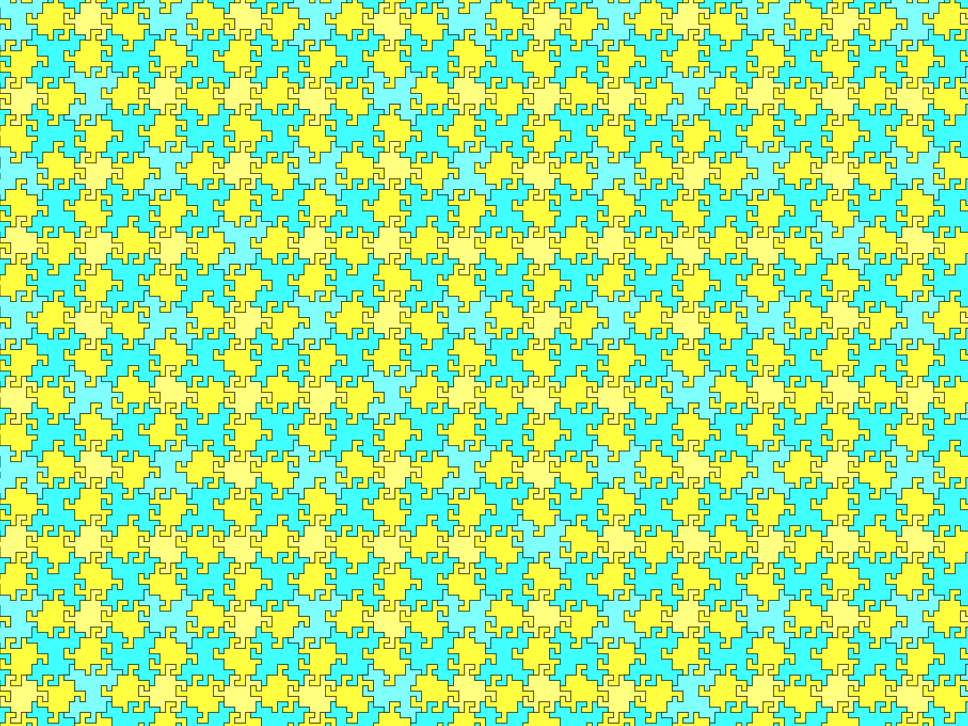
???

Remark If all tiling sets of polyominoes admit a biperiodic tiling for a given k , the k -Polyomino Problem is decidable.



Theorem[Goodman-Strauss 1999] There exists an **aperiodic set of 2 polyominoes**. (as interpreted by Jolivet)

Theorem[Ammann et al 1992] There exists an aperiodic set of **8 polyominoes for tiling by translation**.



		<i>Tiling</i>			
		1	2 – 4	5+	
aperiodic set	No	?	Yes	Yes	
k -PP	$O(n^2)$?	?	undecidable	
		1	2 – 7	8 – 10	11+

Tiling by translation

The following (old) problem is still open...

Open Problem Does there exist an aperiodic polyomino?