A
Small Minimal Aperiodic Reversible Turing machine

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joint work with J. Cassaigne, R. Torres and A. Gajardo

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Motivation

Solve a conjecture that we had with J. Kari a few years ago:

Periodicity and Immortality in Reversible Computing

Jarkko Kari (Dpt. of Mathematics, University of Turku, Finland)
Nicolas Ollinger (LIF, Aix-Marseille Université, CNRS, France)

Toruń, Poland — August 27, 2008

J. Kari and N. Ollinger. Periodicity and Immortality in Reversible Computing.
E. Ochmański and J. Tyszkiewicz (Eds.): **MFCS 2008**, LNCS 5162, pp. 419–430, 2008.

. . .

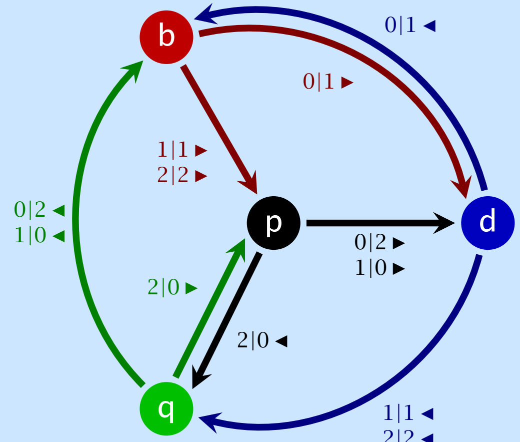
Open Problems with conjectures

Conjecture 1 It is undecidable whether a given complete 2-**RCM** admits a periodic configuration. *(proven if you remove complete or replace 2 by 3)*

Conjecture 2 There exists a complete **RTM** without a periodic configuration. *(known for DTM [BCN02])*

Conjecture 3 It is undecidable whether a given complete **RTM** admits a periodic configuration. *(known for DTM [BCN02])*

Theorem To find if a given **complete reversible Turing machine** admits a **periodic orbit** is Σ_1^0-complete.



1. Dynamics of Turing machines

Turing machines

Definition A **Turing machine** is a triple (Q, Σ, δ) where Q is the finite set of states, Σ is the finite set of tape symbols and $\delta : Q \times \Sigma \to Q \times \Sigma \times \{◀, ▶\}$ is the transition function.

Transition $\delta(s, a) = (t, b, d)$ means:
"in state s, when reading the symbol a on the tape,
replace it by b move the head in direction d and enter state t."

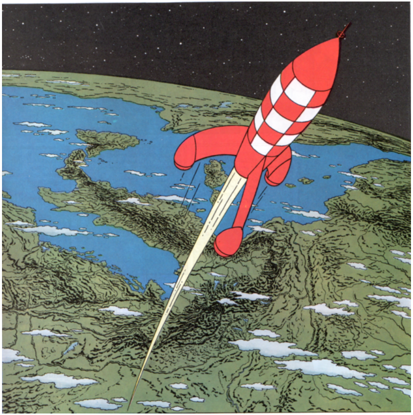
Remark We do not care about blank symbol or initial and final states, we see Turing machines as dynamical systems.

Moving head dynamics

$$X_H = Q \times \mathbb{Z} \times \Sigma^{\mathbb{Z}} \cup \Sigma^{\mathbb{Z}}$$
$$T_H : X_H \to X_H$$

$\cdots 000000\mathbf{b}000000000 \cdots$
$\cdots 0000001\mathbf{d}00000000 \cdots$
$\cdots 000000\mathbf{b}110000000 \cdots$
$\cdots 0000001\mathbf{p}10000000 \cdots$
$\cdots 00000010\mathbf{d}0000000 \cdots$
$\cdots 0000001\mathbf{b}01000000 \cdots$
$\cdots 00000011\mathbf{d}1000000 \cdots$
$\cdots 0000001\mathbf{q}11000000 \cdots$
$\cdots 000000\mathbf{b}101000000 \cdots$
$\cdots 0000001\mathbf{p}01000000 \cdots$
$\vdots$

Hergé. *On a marché sur la lune*. Casterman, 1954.



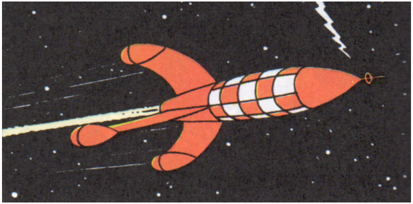
Long shot

Moving tape dynamics

$$X_T = Q \times \Sigma^{\mathbb{Z}}$$
$$T_T : X_T \to X_T$$

$\cdots 0000000\mathbf{b}00000000 \cdots$
$\cdots 0000001\mathbf{d}00000000 \cdots$
$\cdots 0000000\mathbf{b}11000000 \cdots$
$\cdots 0000001\mathbf{p}10000000 \cdots$
$\cdots 0000010\mathbf{d}00000000 \cdots$
$\cdots 0000001\mathbf{b}01000000 \cdots$
$\cdots 0000011\mathbf{d}10000000 \cdots$
$\cdots 0000001\mathbf{q}11000000 \cdots$
$\cdots 0000000\mathbf{b}10100000 \cdots$
$\cdots 0000001\mathbf{p}01000000 \cdots$
$\vdots$

Hergé. *On a marché sur la lune*. Casterman, 1954.



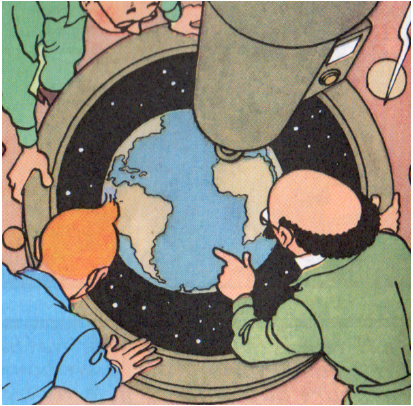
Tracking shot

Trace subshift

$$S_T \subseteq (Q \times \Sigma)^\omega$$

0 0 1 1 0 0 1 1 1 0
b d b p d b d q b p . . .

Hergé. *On a marché sur la lune*. Casterman, 1954.



Point of view shot

Simple dynamical properties

Definition A point $x \in X$ is **periodic** if it admits a **period** $p > 0$ such that $T^p(x) = x$.

Definition A machine is **periodic** if every point is periodic.

Remark Periodicity implies uniform periodicity: $T^p = \mathrm{Id}$.

Theorem[KO08] The **periodicity problem** is Σ_1^0-complete.

Definition A machine is **aperiodic** if it has no periodic point.

Partial vs complete machines

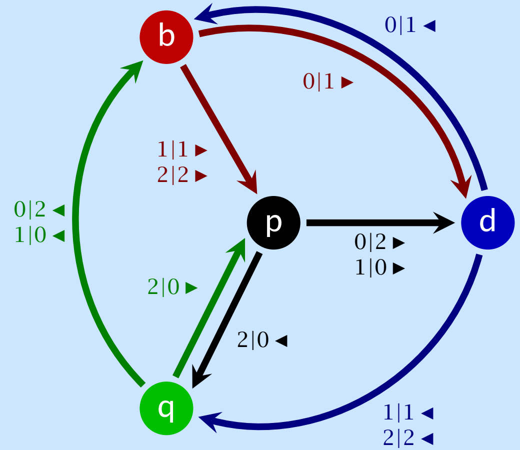
Definition A TM is **complete** if δ is completely defined, otherwise **undefined transitions** of a partial δ correspond to **halting configurations**.

Definition A point is **mortal** if it eventually **halts**.

Thm[Hooper66] The **immortality problem** is Π_1^0-complete.

Rk **Mortality** is different from **totality** which is Π_2^0-complete.

Thm[KO08] The result is the same for **reversible TM**.



2. Reversible Turing machines

Reversible Turing machines

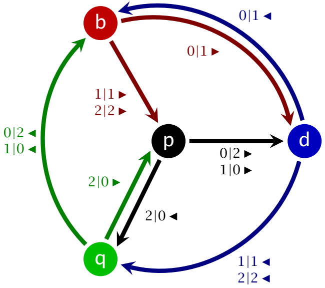
Intuitively, a TM is **reversible** if there exists another TM to compute backwards: "$T_2 = T_1^{-1}$". **Forget technical details**...

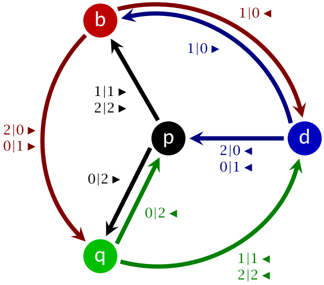
Definition A TM is **reversible** if δ can be decomposed as:

$$\begin{array}{rcl} \delta(s,a) & = & (t,b,\rho(t)) \qquad \text{where } (t,b) = \sigma(s,a) \\ \rho & : & Q \to \{◀,▶\} \\ \sigma & \in & \mathfrak{S}_{Q\times\Sigma} \end{array}$$

Remark $\delta^{-1}(t,b) = (s,a,♦(\rho(s)))$

A complete RTM



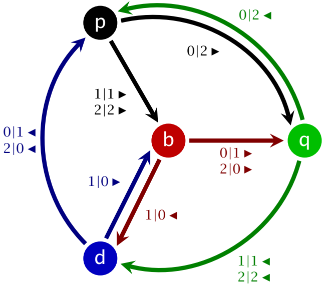


It is **time-symmetric**:

its own inverse up to

state/symbol permutation.

1	⇔	2
b	⇔	p
d	⇔	q



Searching for a reduction

We want to prove the following:

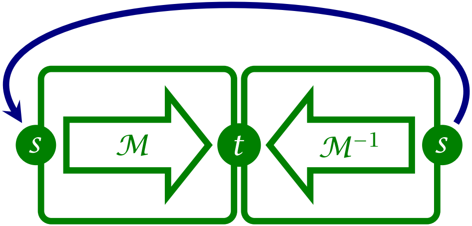
Theorem To find if a given **complete reversible Turing machine** admits a **periodic orbit** is Σ_1^0-complete.

In the partial case we use the following tool:

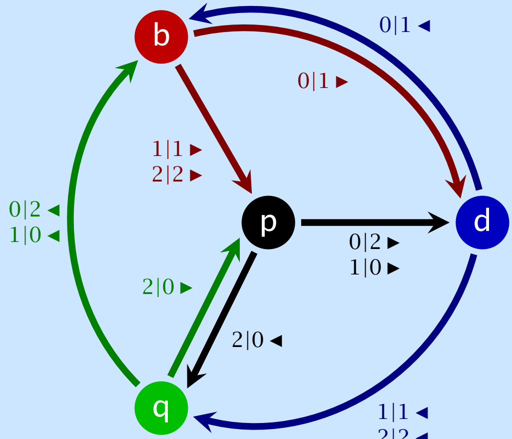
Prop[KO08] To find if a given **(aperiodic) RTM** can reach a given state t from a given state s is Σ_1^0-complete.

The partial case

Principle of the reduction Associate to an (aperiodic) RTM $\mathcal{M}$ with given s and t a new machine with a periodic orbit if and only if t is reachable from s.



We need to find a way to **complete** the constructed machine.



3. a SMART machine

Cassaigne machine

Conj[Kůrka97] Every **complete** TM has a **periodic** point.

Thm[BCN02] No, here is an **aperiodic** complete TM.

Rk It relies on the **bounded search** technique [Hooper66].

In 2008, I asked **J. Cassaigne** if he had a reversible version of the BCN construction...

...he answered with a small machine $\mathfrak{C}$ which is a reversible and (drastic) simplification of the BCN machine.

Aperiodicity

Proposition The machine $\mathfrak{C}$ is **aperiodic**.

Idea of the proof (R. Torres)

1. The behavior starting from a tape of 0 is aperiodic;
2. Every block of 0 eventually grows;
3. Thus $\mathfrak{C}$ is aperiodic.

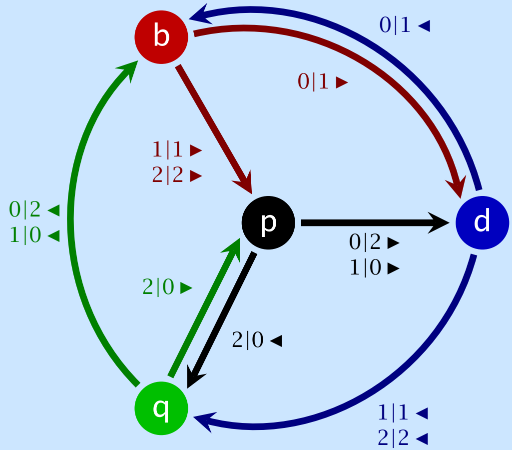
Minimality

The behavior of $\mathfrak{C}$ can be precisely described (R. Torres).

Proposition The behavior starting from a tape of 0 is **dense**.

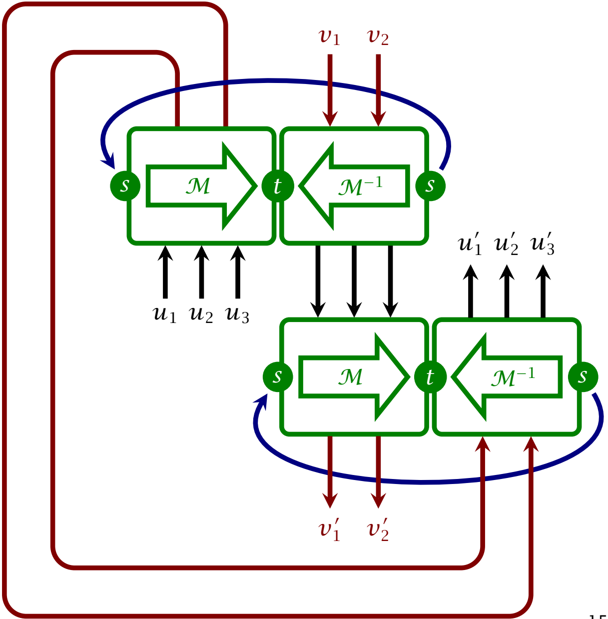
Proposition The trace subshift of $\mathfrak{C}$ is **minimal**.

Proposition The trace subshift of $\mathfrak{C}$ is **substitutive**.

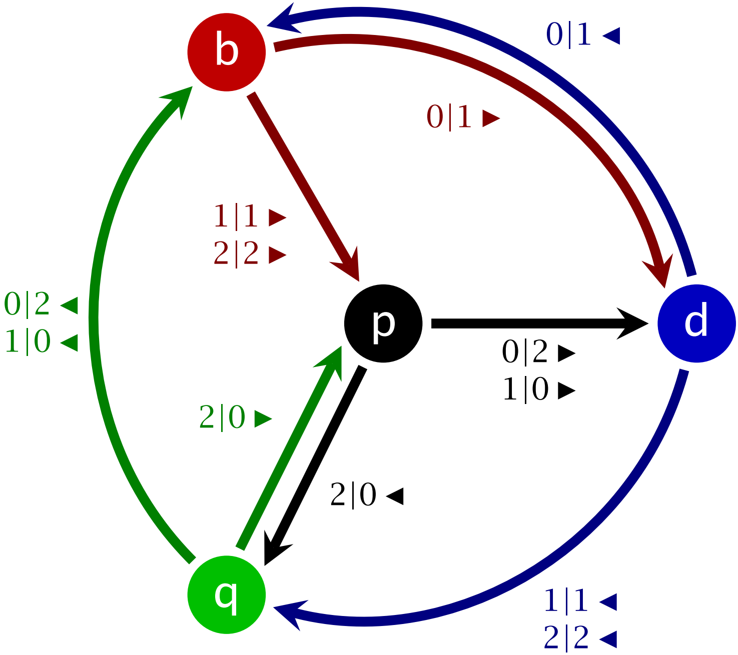


4. Embedding the machine

Embedding trick



Use the transitions of the **Cassaigne machine** to connect the u'_i to the u_i and the v'_i to the v_i.


b
0|1 ◄
0|1 ►
1|1 ►
2|2 ►
0|2 ◄
1|0 ◄
p
d
0|2 ►
1|0 ►
2|0 ►
2|0 ◄
q
1|1 ◄
2|2 ◄

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