Packing, Cutting, Shifting, and Twisting

space-time diagrams of Cellular Automata

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Cellular Automata

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• The global rule applies δ uniformly:

$$G_{\mathcal{A}}(C)_i = \delta\left(C_{i+v_1}, \dots, C_{i+v_\nu}\right)$$

where $N = \{v_1, ..., v_{\nu}\}.$











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Theorem[Hedlund 69]. A map $G : S^{\mathbb{Z}^d} \to S^{\mathbb{Z}^d}$ is the global rule of a *d*-CA if and only if it is continue and commute with the shifts.

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- ♦ Select an appropriate rule (*e.g.* some LGCA).
- Choose a set of initial configurations.
- Construct their space-time diagrams.
- Investigate the diagrams for properties.

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Some properties depend on the granularity of the model!

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Classical Simulations

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Any Nil CA can be simulated by the trivial one



Simulation and Transformations

Idea. A CA \mathcal{A} simulates another CA \mathcal{B} if, up to geometrical transformations, any space-time diagram from \mathcal{B} is a space-time diagram from \mathcal{A} .



Good transformations





 $o^m \circ G_{\mathcal{A}} \circ o^{-m}$

 $G^n_{\mathcal{A}}$

spatial organization temporal organization

Good transformations (2)



 $G_{\mathcal{A}} \circ \sigma_k$

 $\prod_i G_{\mathcal{A}_i}$

information mixing

independent

Generalizing Transformations



 $\varphi:\mathbb{N}\times\mathbb{Z}\to 2^{\{1,\ldots,k\}\times\mathbb{N}\times\mathbb{Z}}$

The new CA must be *completely* defined.

Geometrical Characterization

Theorem. There exist no transformation but compositions of the 4 good previous ones.





PCST transformations

Definition. The $\prod \langle m_i, n_i, k_i \rangle$ regular P⁻CST transformation of a CA \mathcal{A} is the CA $\mathcal{A}^{\prod \langle m_i, n_i, k_i \rangle}$ where

$$\mathcal{A}^{\prod \langle m_i, n_i, k_i \rangle} = \prod_i o^{m_i} \circ \mathcal{A}^{m_i n_i} \circ o^{-m_i} \circ \sigma^{k_i}.$$

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Theorem. The PCST relation of simulation is a quasi-order with a maximal equivalence class.

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Definition of CA rules from elementary rules by an algebraic closure.

The semi-lattice

Theorem. PCST induces a sup semi-lattice with the natural operation $\mathcal{A} \times \mathcal{B}$ as a sup operation.

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Ideals capture interesting notions:

- reversibility (principal),
- ultimately periodic,
- simple signal,
- naive non-chaoticity.