

Subdivision of oriented cycles in oriented graphs

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Orientation and colouring

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Theorem (Bondy)

D is strong, $\chi(D) = k \Rightarrow D$ has a directed cycle with at least k vertices.

Blocks

Definition

In an oriented path or cycle, a *block* is a maximal directed sub-path.

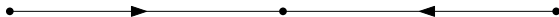


Figure : A path with 2 blocks.

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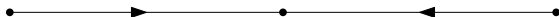


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Theorem (L.Addario-Berry, F.Havet and S. Thomassé)

If $\chi(D) = k$, then D contains every path with 2 blocks on k vertices.

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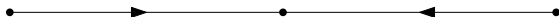


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Theorem (L.Addario-Berry, F.Havet and S. Thomassé)

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Lemma (Gyárfás, Thomassen)

There exists oriented graph D_i such that $\chi(D_i) \geq i$ and every cycle in D_i has at least four blocks.

Conjecture (L. Addario-Berry et al.)

For all k and l , there exists an integer $\alpha(k, l)$ such that :
 $\chi(D) > \alpha(k, l)$ and D **strong** $\Rightarrow D$ contains an oriented cycle with
two blocks, one of length at least k , and the other at least l .

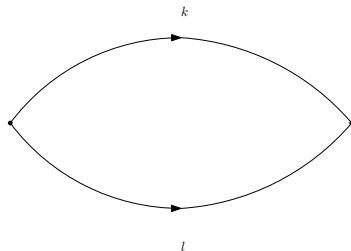


Figure : Cycle with two blocks.

This can be seen as finding a subdivision of $C(k, l)$ the cycle with one block of size k and one block of size l .

Results

We showed the following:

Theorem

$\alpha(1, 2) = \alpha(2, 2) = \alpha(1, 3) = 4$, $\alpha(2, 3) = 5$ and $\alpha(3, 3) \leq 7$.

Where $\alpha(k, l)$ is the smallest integer such that:

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Theorem

Let D be a 2-strong digraph. If $\chi(D) \geq (k + l - 2)(k - 1) + 2$, then D contains a subdivision of $C(k, l)$.

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A *handle decomposition* of D starting at $v \in V(D)$ is a triple $(v, (h_i)_{1 \leq i \leq p}, (D_i)_{0 \leq i \leq p})$, where $(D_i)_{0 \leq i \leq p}$ is a sequence of strongly connected digraphs and $(h_i)_{1 \leq i \leq p}$ is a sequence of handles such that:

- $V(D_0) = \{v\}$,
- for $1 \leq i \leq p$, h_i is a handle of D_i and D_i is the (arc-disjoint) union of D_{i-1} and h_i , and
- $D = D_p$.

Handle decomposition



v_0

Figure : $D_0 = \{v_0\}$

Handle decomposition

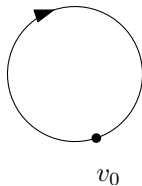


Figure : $D_1 = D_0 \cup h_1$

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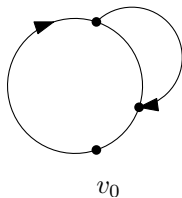


Figure : $D_2 = D_1 \cup h_2$

Handle decomposition

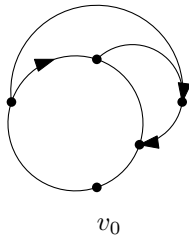
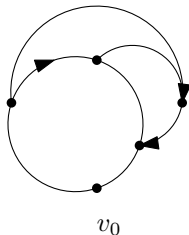


Figure : $D_3 = D_2 \cup h_3$

Handle decomposition



Theorem

Every strong digraph admits a handle decomposition.

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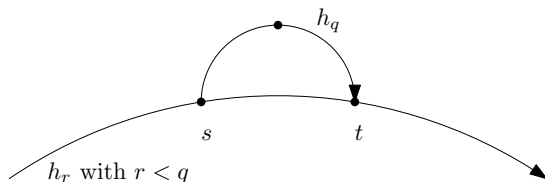


Figure : h_q is an (s, t) dipath

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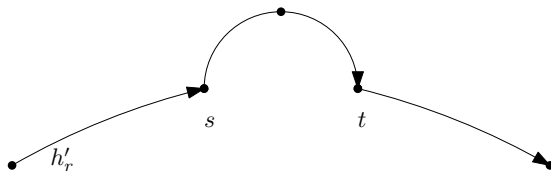


Figure : h'_r is longer than h_r

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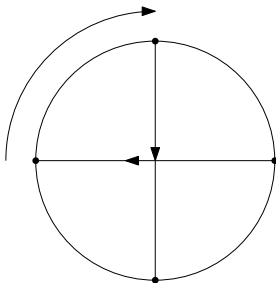


Figure : 2 crossing chords

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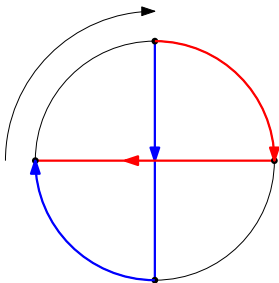


Figure : A subdivision of $C(2, 2)$

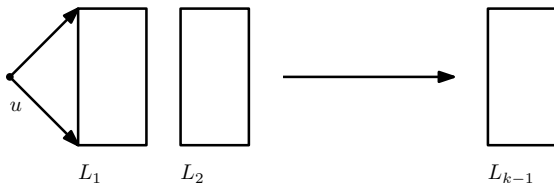
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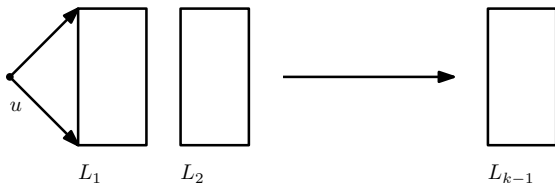
Assume $k > l$ and pick u any vertex. Let $L_i = \{v \mid \text{dist}_D(u, v) = i\}$.
 If $v \in L_k$ we have a subdivision of $C(k, k)$. If $L_k = \emptyset$,
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$\chi(D) \leq 1 + \sum_i \chi(L_i)$, so there exists i , with $\chi(L_i) \geq (k + l - 1)$.

Because $\chi(L_i) \geq (k + l - 1)$, there exists a path with two blocks, one of length k , the other of length $l - 1$.

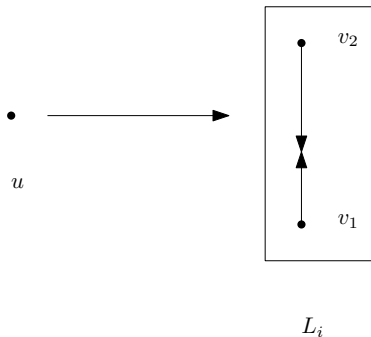


Figure : Path of two blocks in L_i

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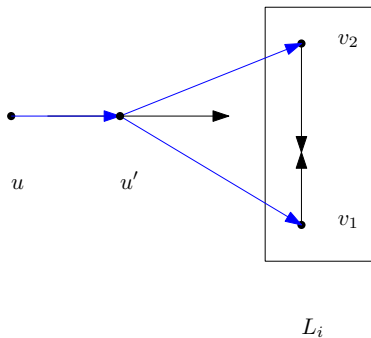


Figure : Subdivision of $C(k, l)$

Open questions

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C has more than two blocks

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- *Which connectivity condition is necessary.*
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Proposition

$\chi(D) > 12$ and D strong $\Rightarrow D$ contains a subdivision of a cycle with 4 blocks.

Thank you!