

RANDOM TILINGS

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1 INTRODUCTION

2 TILINGS AND TILING GRAPHS

3 MIXING TIME

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INTRODUCTION

Families of tilings: dimers, balanced words, kagome tilings and others

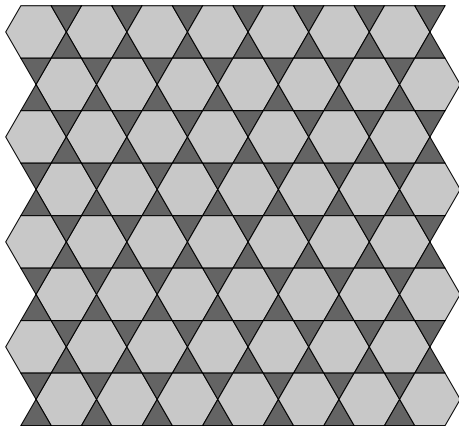
- Random generation
- Tiling graph properties

1 INTRODUCTION

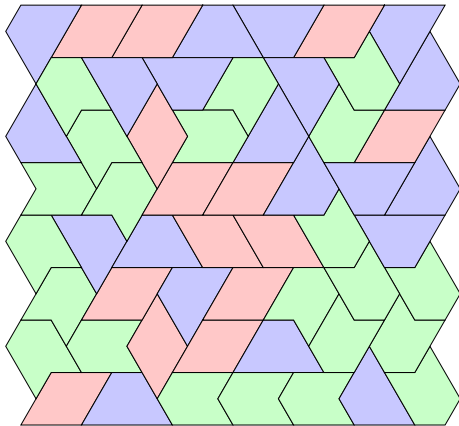
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KAGOME TILINGS

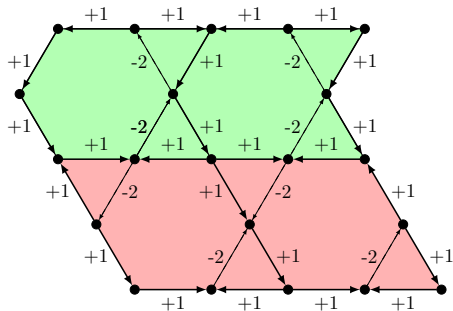


KAGOME TILINGS



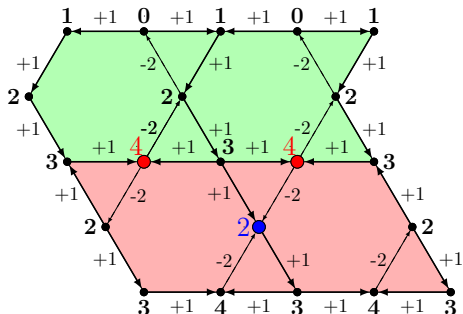
HEIGHT FUNCTION, FLIPS

- *Flip*: local min \leftrightarrow local max.
- Tiled region $R \rightsquigarrow$ tiling graph G_R (distributive lattice).



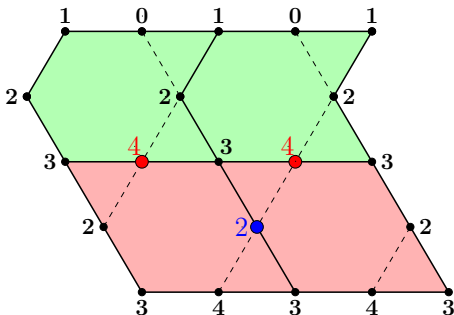
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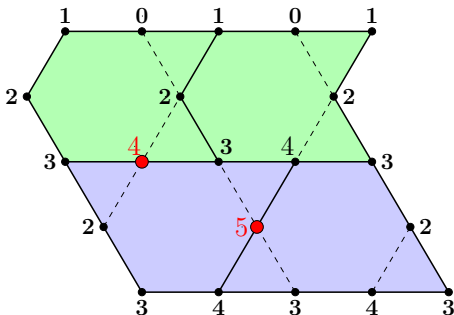
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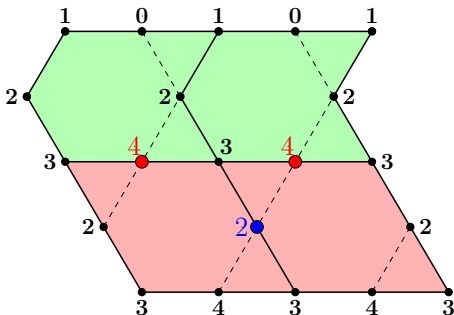
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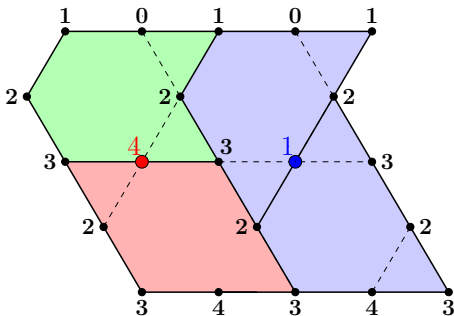
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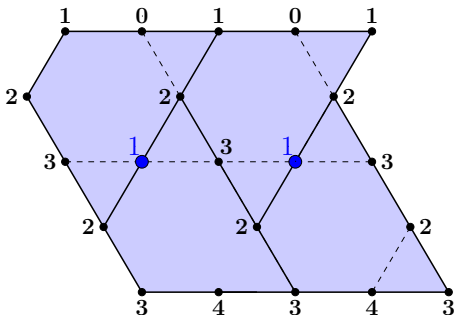
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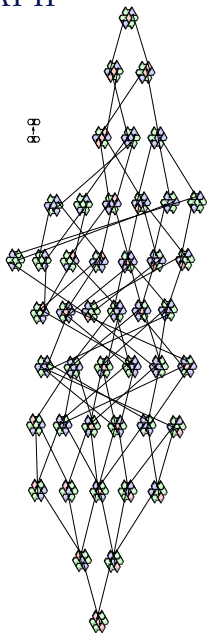
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KAGOME TILING GRAPH

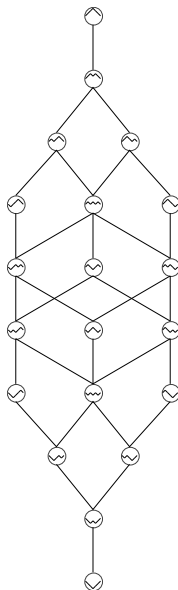
$$D = O(n^{3/2})$$



BALANCED WORDS ON $A = \{a, b\}$

flip: $ab \longleftrightarrow ba$

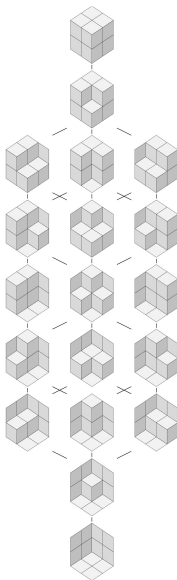
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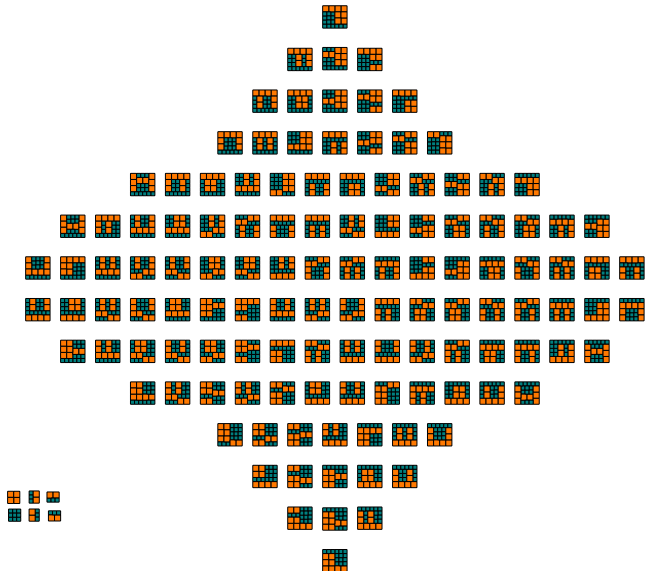
LOZANGE TILINGS

flip: add/remove a cube

$$D = O(n^{3/2})$$



SQUARE TILINGS 2×2 & 3×3



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MARKOV CHAIN ON THE TILING GRAPH

State space $\Omega = \{ \text{tilings of the chosen region } R \} = \{ \text{vertices of } G_R \}$

Starting at t_0 , repeat t times:

- 1 Pick inner vertex v of R
- 2 Pick a direction of a flip: "up" or "down"
- 3 Flip if possible, otherwise stay still.

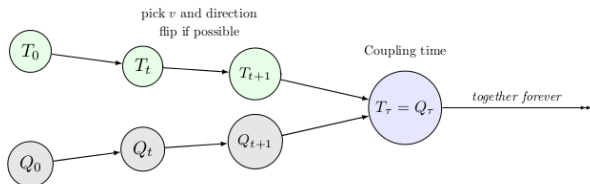
MC: aperiodic + irreducible = ergodic \Rightarrow unique stationary distribution

Goal: bound *Mixing time* τ

(time it takes for MC to be at distance $\varepsilon < \frac{1}{2}$ from stationary distrib)

COUPLING

- Coupling – MC on $\Omega \times \Omega$ defining $(T_t, Q_t)_{t=0}^{\infty}$:



- Theorem (Aldous): $\tau \leq CT_{coupling}$

PATH COUPLING, MODIFIED CHAINS (TOWER OF FLIPS)

- 2-letter words of size n : $\tau_{ab\text{-words}} = O(n^3)$
- k -colorings of a n -vertex graph: $\tau_{k\text{-colorings}} = O(n \ln(ne^{-1}))$
- Lozange tilings of the region of area n :
 $\tau_{\text{lozange}} = O(n^4)$ (Randall, Tetalli),
 $\tau_{\text{lozange}} = O(n^2 \ln n)$ (Wilson)
- Kagome tilings: $\tau_{\text{kagome}} = O(n^2) \leftarrow \text{Conjecture}$

MIXING TIME VIA CONDUCTANCE

$$\Phi = \min_{S:\pi(S)\leq\frac{1}{2}} \frac{\sum_{e\in\delta S} Q(e)}{\pi(S)},$$

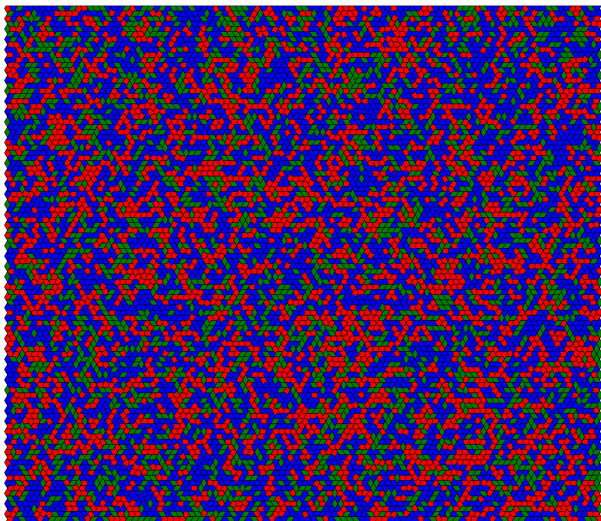
$$\frac{\Phi^2}{2} < \text{Gap} \leq 2\Phi,$$

$$\tau \leq \frac{\ln \pi^{-1}}{\Phi^2}.$$

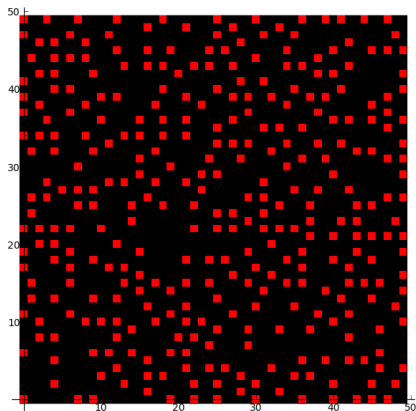
Can we prove that:

- $\Phi \geq \frac{1}{\text{poly}(n)}$,
- G has no bottleneck,
- G is a $(\frac{N}{2}, 1 + \epsilon)$ vertex expander,
- there are $\text{poly}(n)$ canonical paths passing through each edge ?

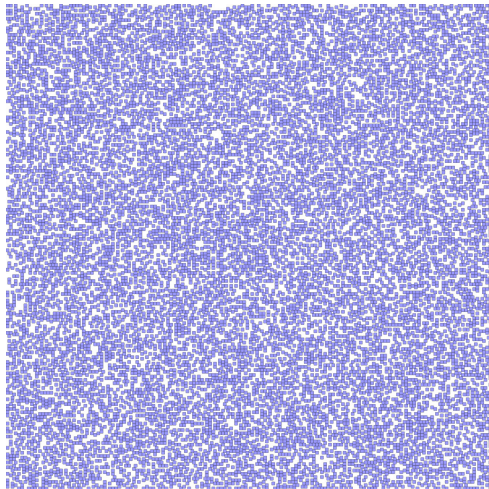
COUPLING FROM THE PAST FOR KAGOME



BOUNDING CHAIN FOR SQUARES 1×1 & 2×2



BOUNDING CHAIN FOR SQUARES 1×1 & 2×2



REFERENCES

- 1 M. Dyer, C. Greenhill, A More Rapidly Mixing Markov Chain for Graph Colorings, *Random Structures and Algorithms*, vol. 13, 1998, 285-317.
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- 6 D.B. Wilson, Mixing Times of Lozenge Tiling and Card Shuffling Markov Chains, *The Annals of Applied Probability*, 2004, Vol. 14, No. 1, 274–325.