

# Simulation between signal machines

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Partially funded by PHC Gundichapur n°909536E



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Laboratoire d'Informatique de l'École polytechnique



Algorithmic Questions in Dynamical Systems  
March 2018 — Toulouse

- 1 Introduction
- 2 Signal Machines (Introduction and Definition)
- 3 Relations to Models of Computation
- 4 Intrinsically Universal Family of Signal Machines
- 5 Non-determinism (work in progress)
- 6 conclusion

# Outline

## One model/dynamical system

- signal machines

## Relations to “usual” computational models

## Intrinsic universality

## Non determinism

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- Turing machines
- Linear Blum, Shub and Smale model

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- a family only

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## Intrinsic universality

- a family only

## Non determinism

- same power

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## Cellular Automata: signal use

## Firing Squad Synchronization (Goto, 1966)

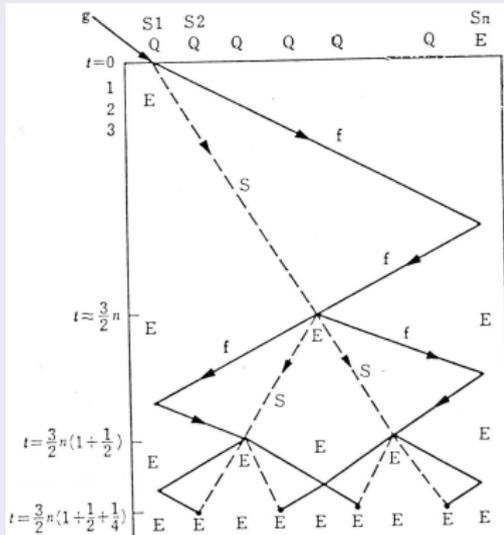
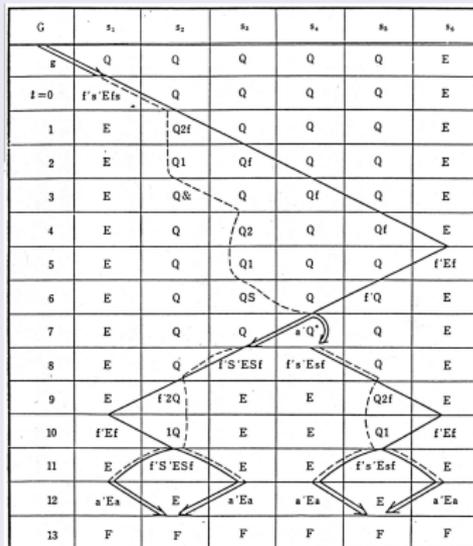
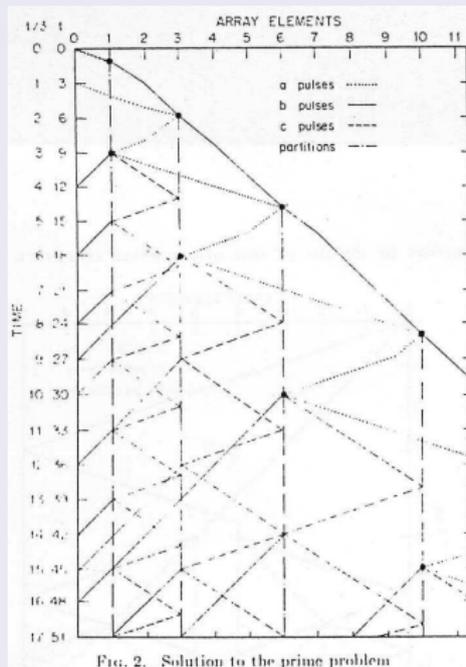


図 3-5 一斉射撃の問題 (連続近似)

図 3-6 一斉射撃解 ( $n=6$ )

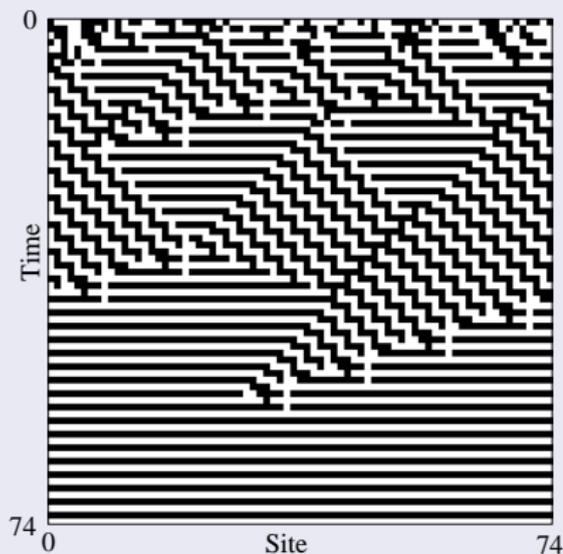
# CA: Conception with signals

Fischer (1965)

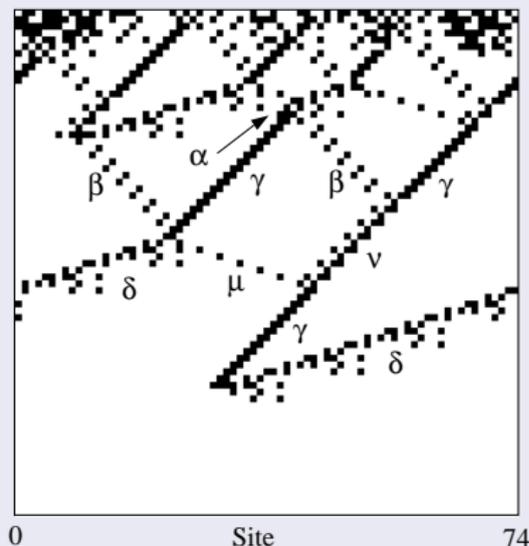


## CA: Analyzing with Signals

Das et al. (1995)

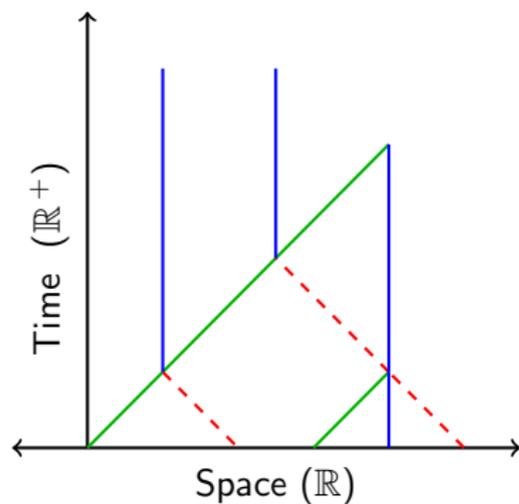
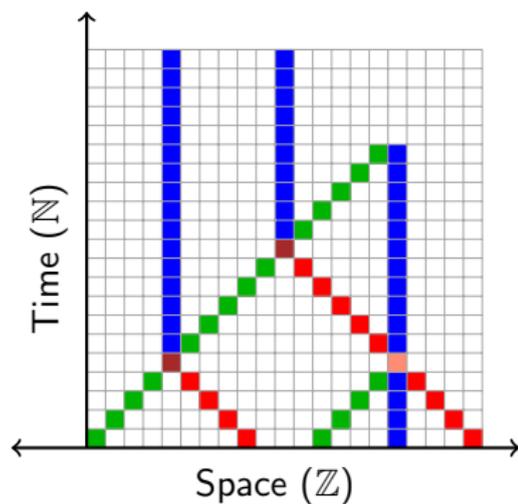


(a) Space-time diagram.



(b) Filtered space-time diagram.

# Signals



- Signal (meta-signal)
- Collision (rule)

# Vocabulary and Example: Find the Middle

M |

M |

Meta-signals (speed)

M (0)

Collision rules

# Vocabulary and Example: Find the Middle

div  M |

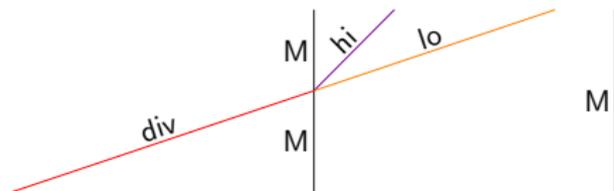
M |

## Meta-signals (speed)

M	(0)
div	(3)

## Collision rules

# Vocabulary and Example: Find the Middle



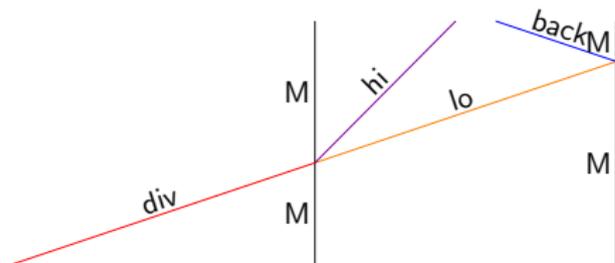
## Meta-signals (speed)

M	(0)
div	(3)
hi	(1)
lo	(3)

## Collision rules

$$\{ \text{div}, M \} \rightarrow \{ M, \text{hi}, \text{lo} \}$$

# Vocabulary and Example: Find the Middle



## Meta-signals (speed)

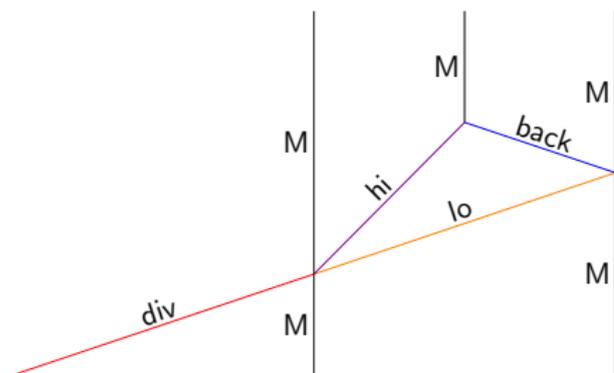
M	(0)
div	(3)
hi	(1)
lo	(3)
back	(-3)

## Collision rules

$$\{ \text{div}, M \} \rightarrow \{ M, \text{hi}, \text{lo} \}$$

$$\{ \text{lo}, M \} \rightarrow \{ \text{back}, M \}$$

# Vocabulary and Example: Find the Middle



## Meta-signals (speed)

M	(0)
div	(3)
hi	(1)
lo	(3)
back	(-3)

## Collision rules

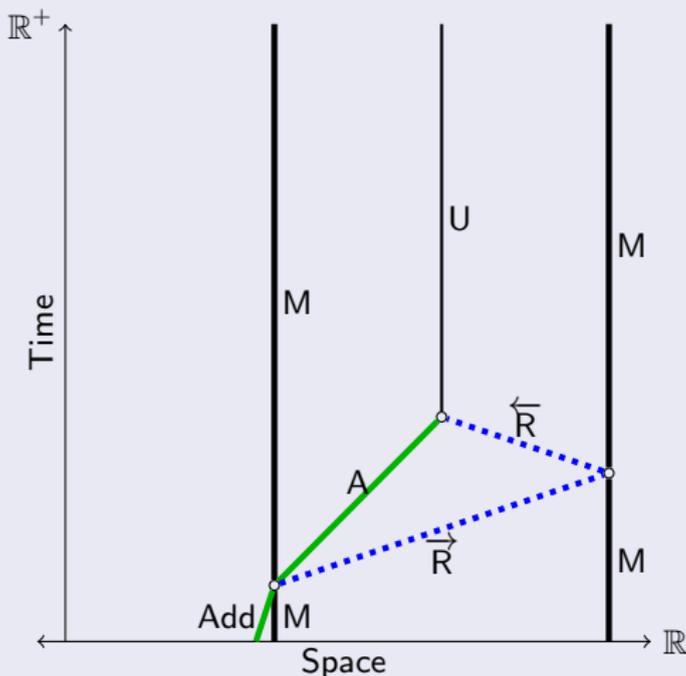
- $\{ \text{div}, M \} \rightarrow \{ M, \text{hi}, \text{lo} \}$
- $\{ \text{lo}, M \} \rightarrow \{ \text{back}, M \}$
- $\{ \text{hi}, \text{back} \} \rightarrow \{ M \}$

# Stack Implantation

Name	Speed
Add, Rem	1/3
A, E	1
U, M	0
$\vec{R}$	3
$\overleftarrow{R}$	-3

Collision rules

- $\{\text{Add}, M\} \rightarrow \{M, A, \vec{R}\}$
- $\{\vec{R}, M\} \rightarrow \{\overleftarrow{R}, M\}$
- $\{A, \overleftarrow{R}\} \rightarrow \{U\}$
- $\{\vec{R}, U\} \rightarrow \{\overleftarrow{R}, U\}$
- $\{\text{Rem}, M\} \rightarrow \{M, E\}$
- $\{E, U\} \rightarrow \{\}$

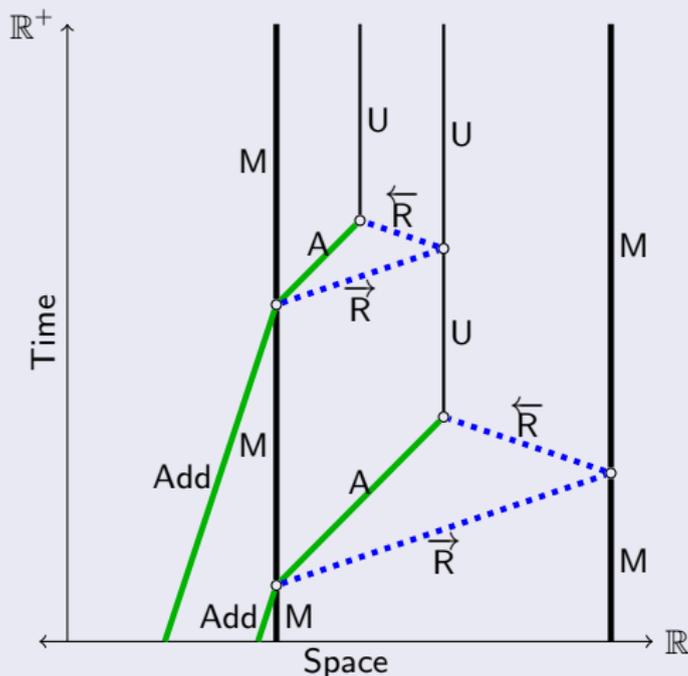


# Stack Implantation

Name	Speed
Add, Rem	1/3
A, E	1
U, M	0
$\vec{R}$	3
$\overleftarrow{R}$	-3

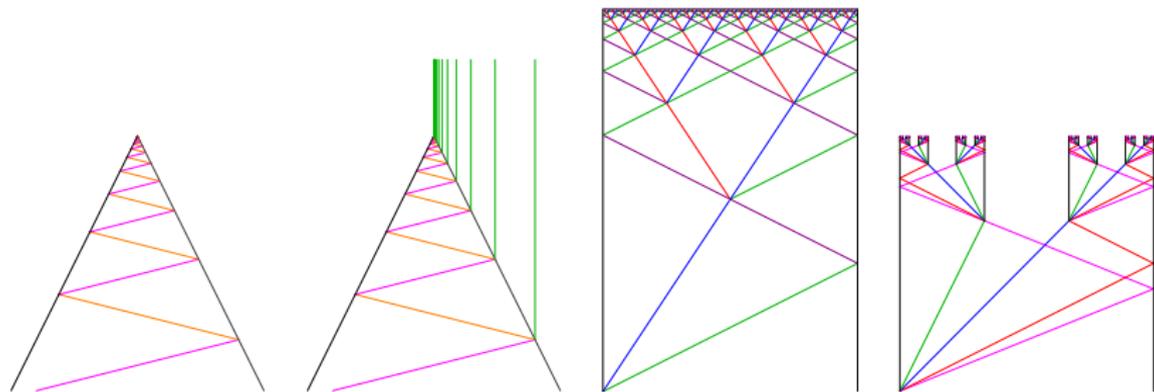
Collision rules

- $\{\text{Add}, M\} \rightarrow \{M, A, \vec{R}\}$
- $\{\vec{R}, M\} \rightarrow \{\overleftarrow{R}, M\}$
- $\{A, \overleftarrow{R}\} \rightarrow \{U\}$
- $\{\vec{R}, U\} \rightarrow \{\overleftarrow{R}, U\}$
- $\{\text{Rem}, M\} \rightarrow \{M, E\}$
- $\{E, U\} \rightarrow \{\}$

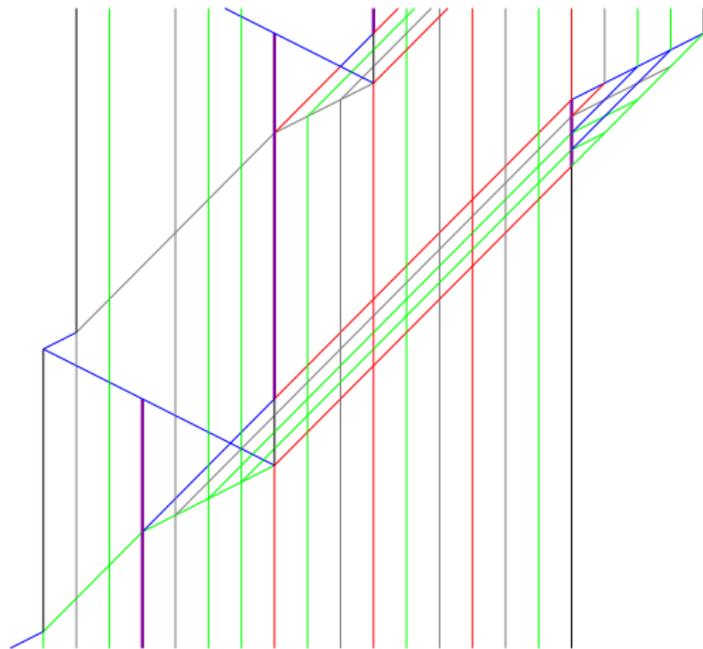




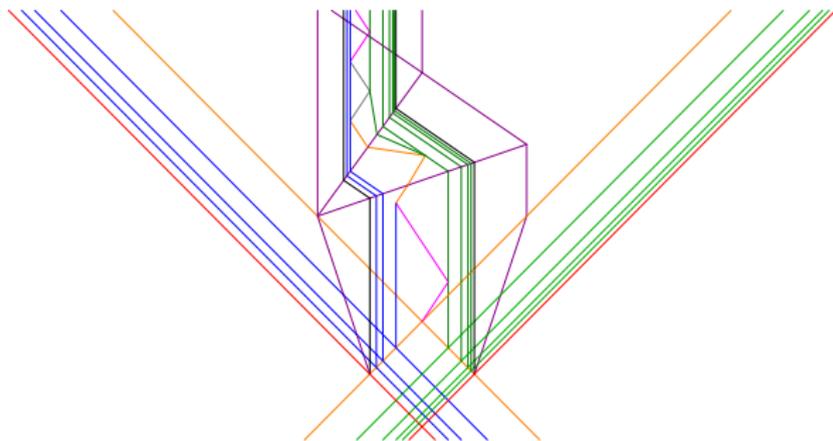
# Fractal Generation



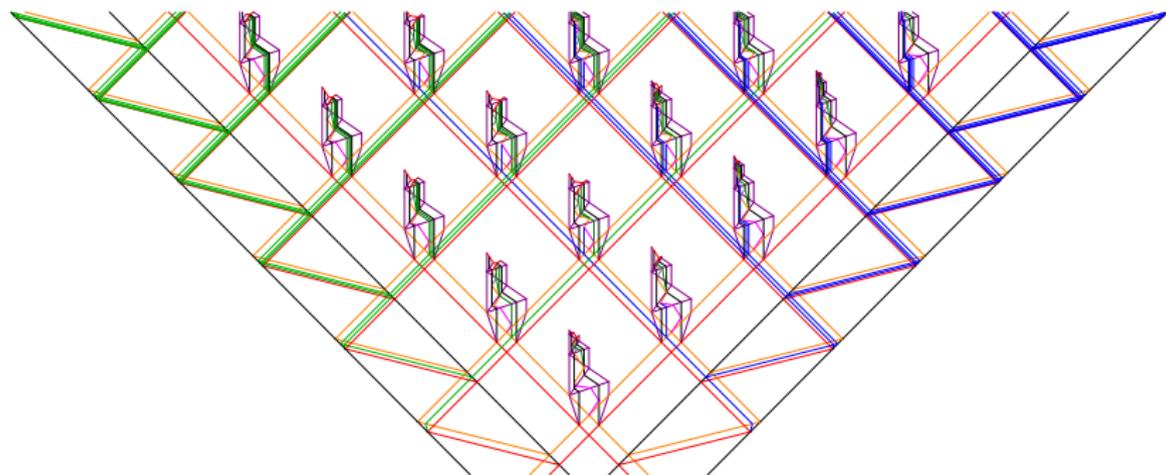
# Complex Dynamics



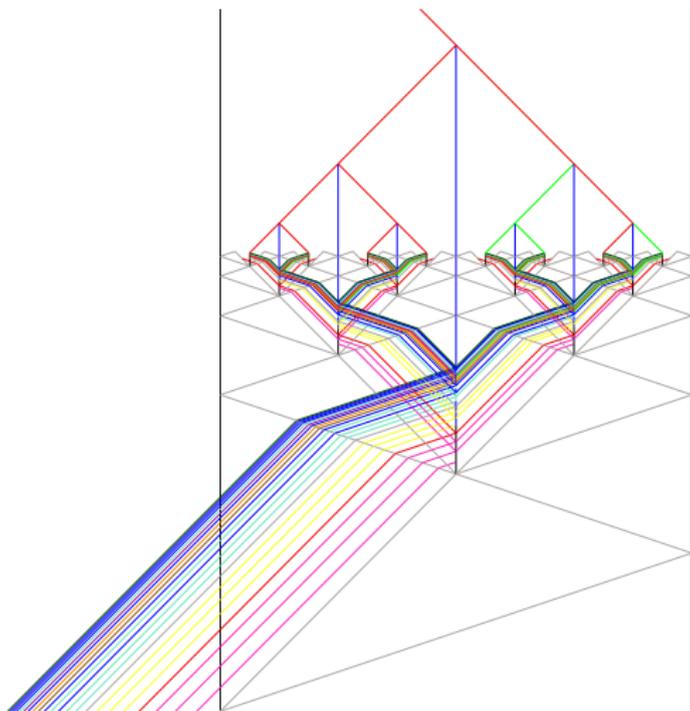
# Complex Dynamics



# Complex Dynamics



# Complex Dynamics



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  - Discrete computation: Turing Machines
  - Analog Computation: linear Blum, Shub and Smale
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# Turing-computation

## Turing Machine

 $q_f$   

	b	b	a	b	#
--	---	---	---	---	---

 $q_f$   

^	b	b	a	b	#
---	---	---	---	---	---

 $q_f$   

^	b	b	a	b	#
---	---	---	---	---	---

 $q_f$   

^	b	b	a	b	#
---	---	---	---	---	---

 $q_2$   

^	b	b	a	#
---	---	---	---	---

 $q_1$   

^	b	b	#
---	---	---	---

 $q_1$   

^	b	a	#
---	---	---	---

 $q_1$   

^	a	a	#
---	---	---	---

 $q_i$   

^	a	b	#
---	---	---	---

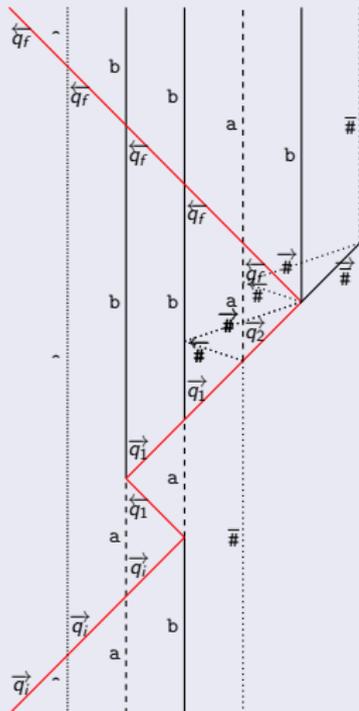
 $q_i$   

^	a	b	#
---	---	---	---

 $q_i$   

*	a	b	#
---	---	---	---

## Simulation



# Turing-computation

## Also with restrictions

- all different speed
- only  $2 \rightarrow 2$  rules (conservative)
- one-to-one rules (reversible)

## Any above and

- rational ( $\mathbb{Q}$ )

## Rational machines

- speeds  $\in \mathbb{Q}$
- initial positions  $\in \mathbb{Q}$
- $\Rightarrow$  collision coordinates  $\in \mathbb{Q}$
- exact simulation on computer/TM

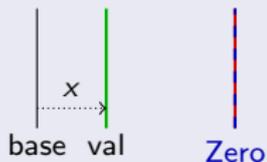
## Undecidability

- finite number de collisions
- meta-signal apperance
- use of a rule
- disappearing of all signals
- involvement of a signal in any collision
- extension on the side, etc.

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# Computing with Real Numbers

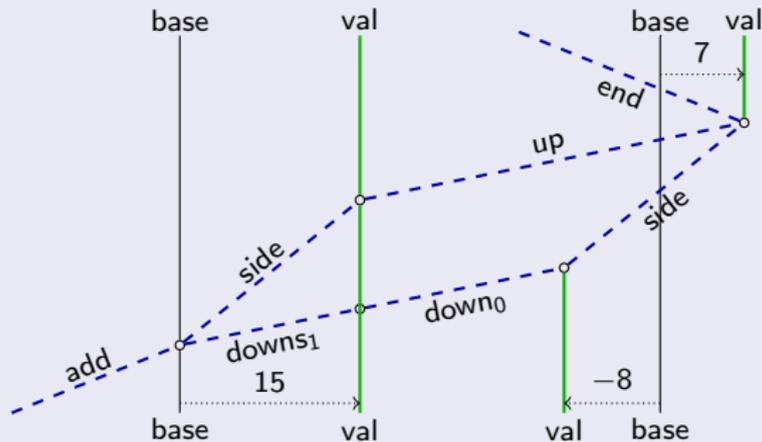
## Encoding



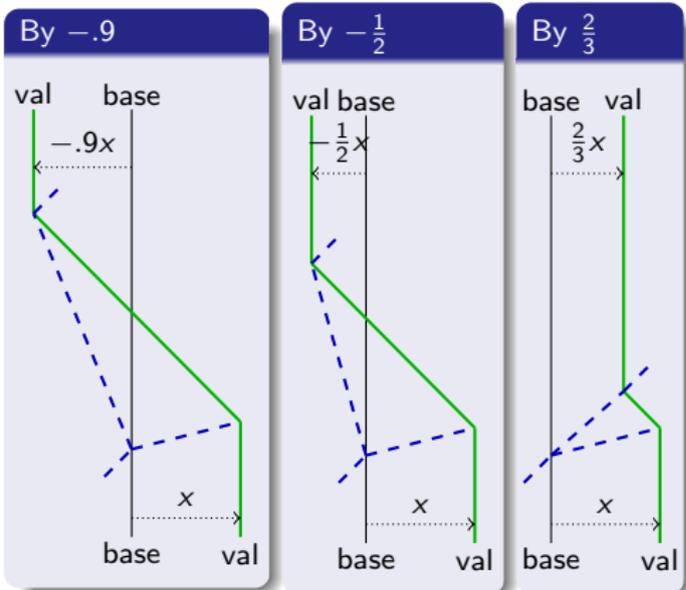
## Sign test



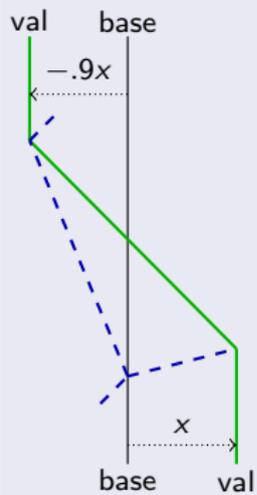
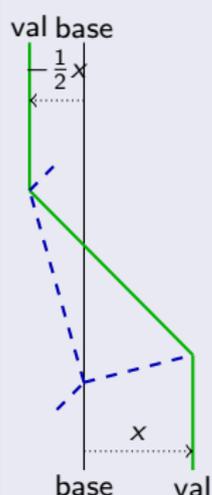
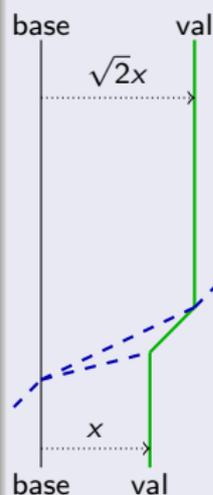
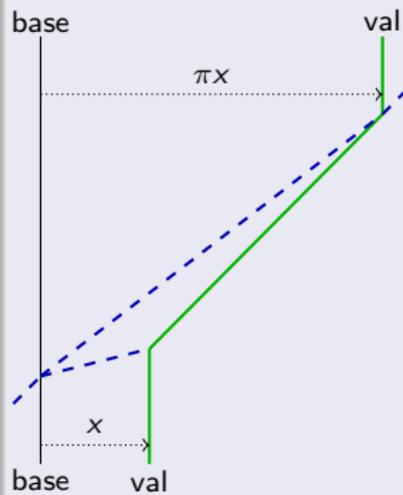
## Addition



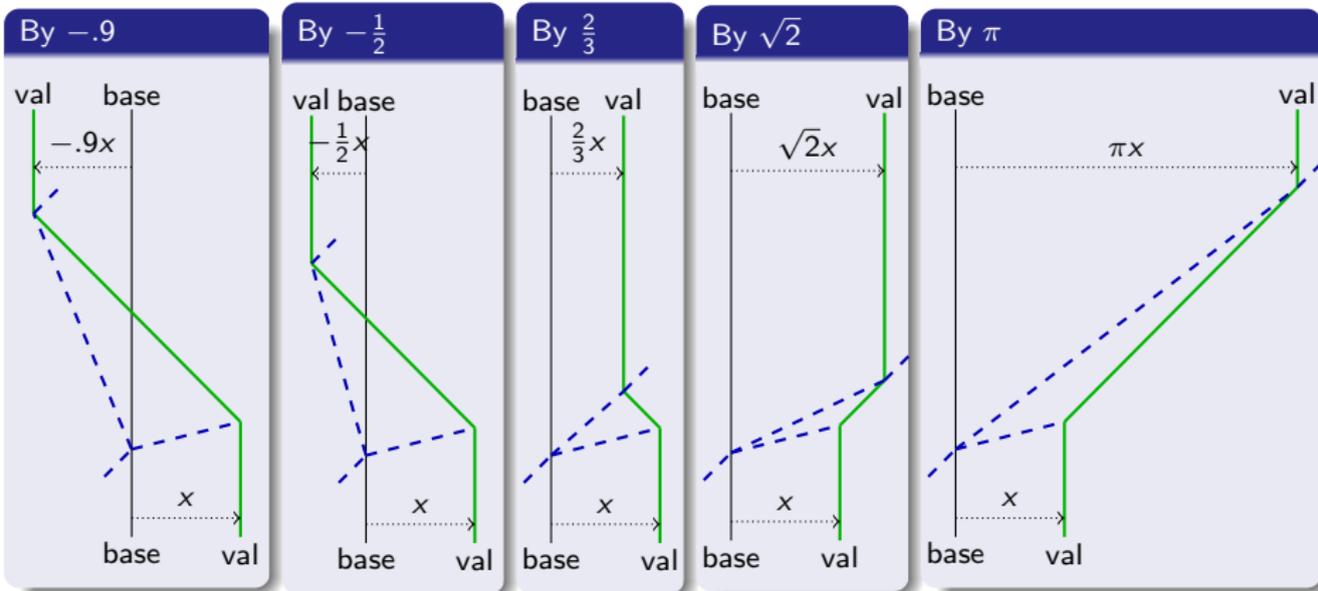
# Multiplication by a constant



# Multiplication by a constant

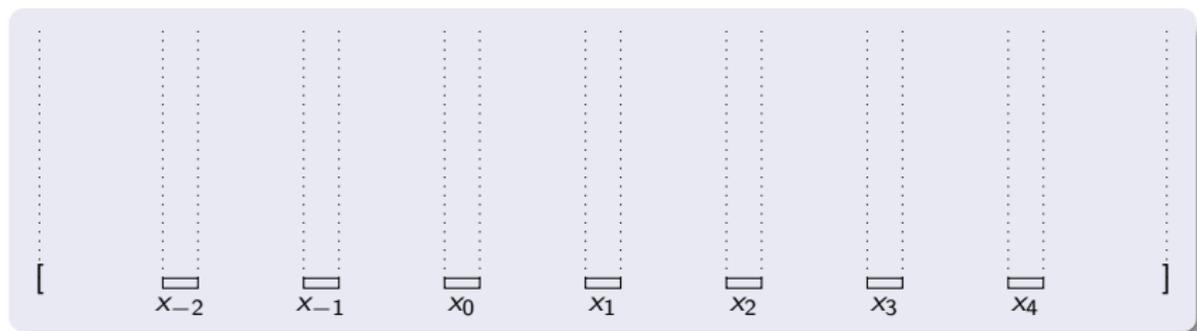
By  $-0.9$ By  $-\frac{1}{2}$ By  $\frac{2}{3}$ By  $\sqrt{2}$ By  $\pi$ 

# Multiplication by a constant



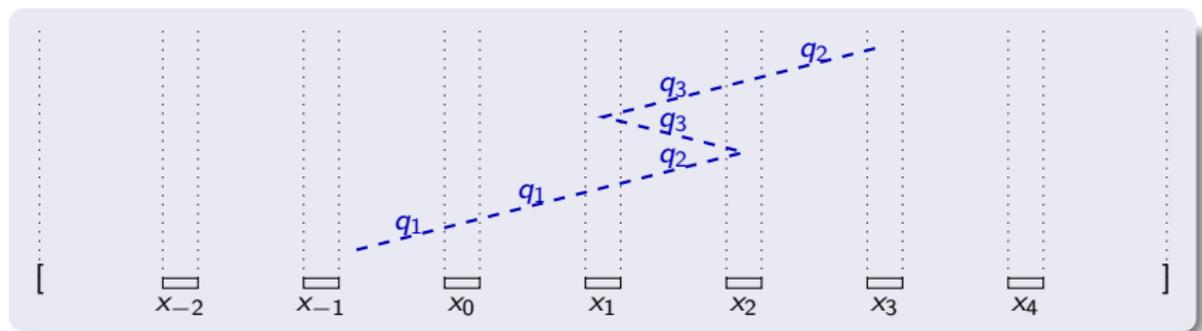
- Signal speeds are constants of the machine
- If  $x \leq 0$  then val is met before base

# Zooming out



Finite sequence of real numbers

# Zooming out



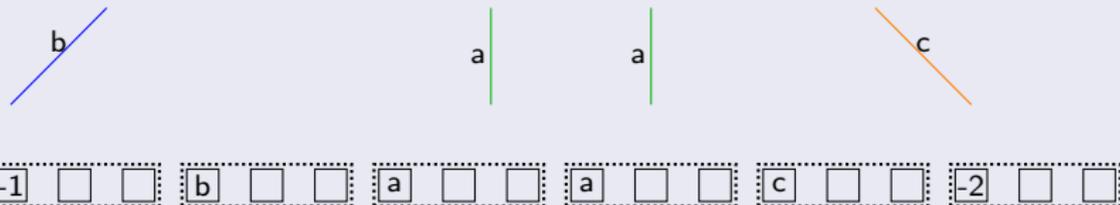
## Finite sequence of real numbers + Dynamics

- finite state automata
- sign test
- addition, multiplication by constant
- (set constant value)
- (enlarge the array)

Like a Turing machine with real numbers on the tape

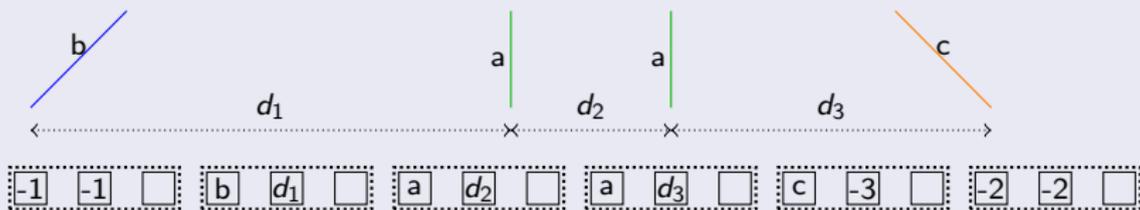
# Linear Blum, Shub and Smale with shift

## Encoding a configuration



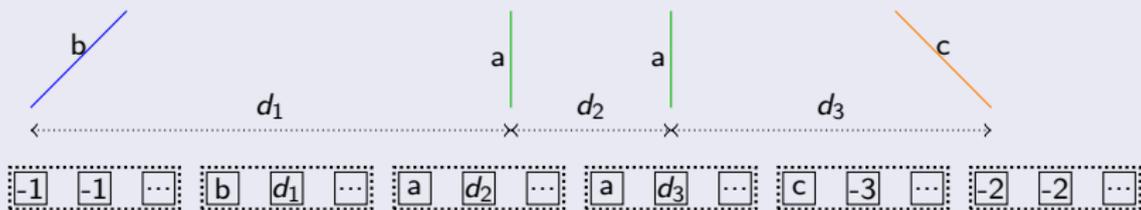
# Linear Blum, Shub and Smale with shift

## Encoding a configuration



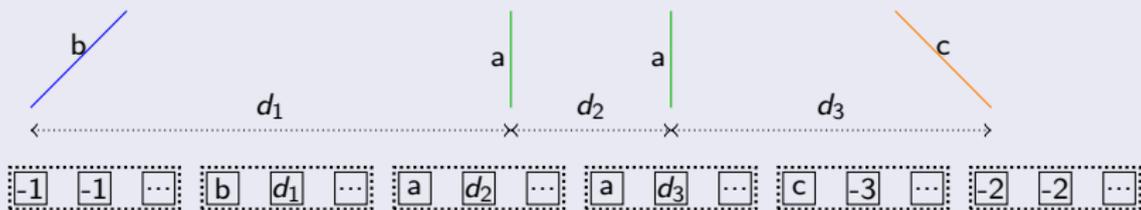
# Linear Blum, Shub and Smale with shift

## Encoding a configuration



# Linear Blum, Shub and Smale with shift

## Encoding a configuration



## Simulating a signal machine: loop

- 1 Compute the minimum time to a collision,  $\delta$
- 2 Advance time by  $\delta$  (update all distances)
- 3 Process collision(s)

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## Intrinsic Universality

Being able to *simulate* any other dynamical system of the its *class*.

## Cellular Automata

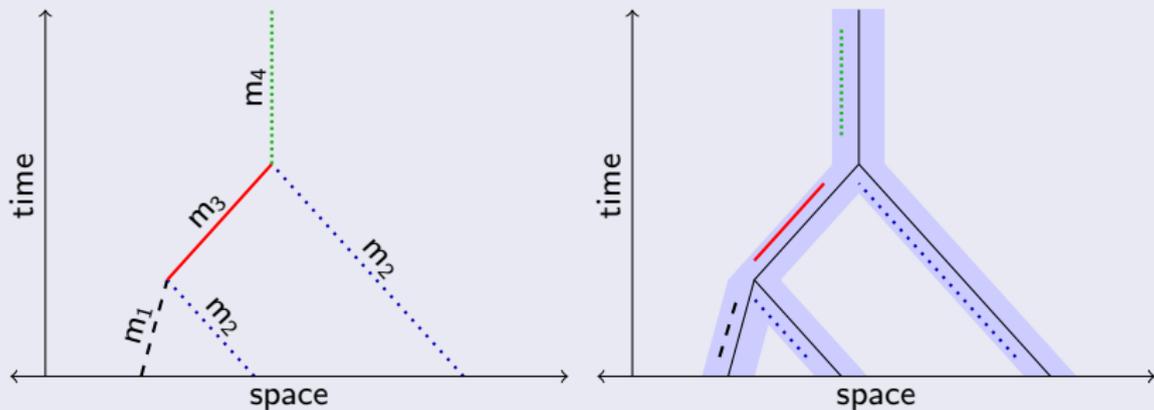
- regular (Albert and Čulik II, 1987; Mazoyer and Rapaport, 1998; Ollinger, 2001)
- reversible (Durand-Lose, 1997)
- freezing [Theyssier et Al.]

## Tile Assembly Systems

- possible at  $T=2$  and above (Woods, 2013)
- impossible at  $T=1$  (Meunier et al., 2014)

# Simulation for Signal Machines

## Space-Time Diagram Mimicking



## Signal Machine Simulation

$\mathcal{U}_S$  simulates  $\mathcal{A}$  if there is function from the configurations of  $\mathcal{A}$  to the ones of  $\mathcal{U}_S$  s.t. the space-time issued from the image always mimics the original one.

## Theorem

- For any finite set of real numbers  $\mathcal{S}$ , there is a signal machine  $\mathcal{U}_{\mathcal{S}}$ , that can simulate any machine whose speeds belong to  $\mathcal{S}$ .
- The set of  $\mathcal{U}_{\mathcal{S}}$  where  $\mathcal{S}$  ranges over finite sets of real numbers is an intrinsically universal family of signal machines.

## Rest of this section

Let  $\mathcal{S}$  be any finite set of real numbers,  
let  $\mathcal{A}$  be any signal machine whose speeds belongs to  $\mathcal{S}$ ,  
 $\mathcal{U}_{\mathcal{S}}$  is progressively constructed as simulation is presented.

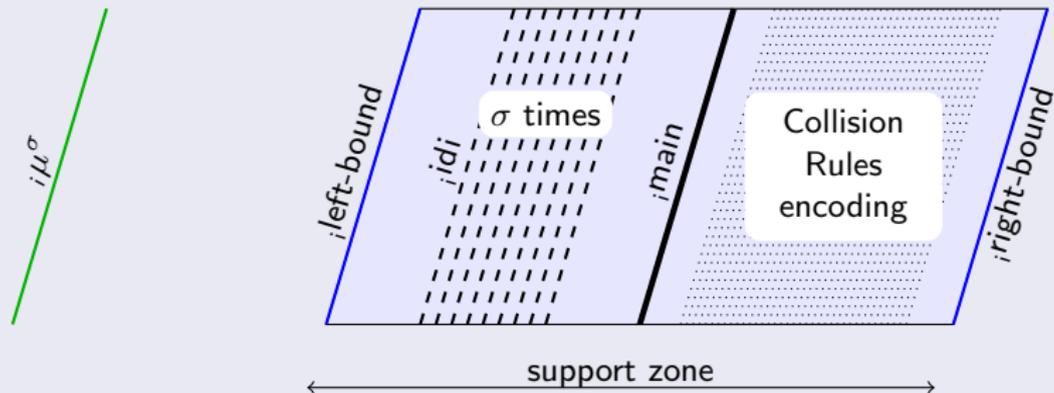
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# Macro-Signal

- Meta-signal of  $\mathcal{A}$  identified with numbers
- Unary encoding of numbers

## Structure

$i\mu^\sigma$ :  $\sigma$ th signal,  $i$ th speed



# Global scheme

## When Support Zones Meet

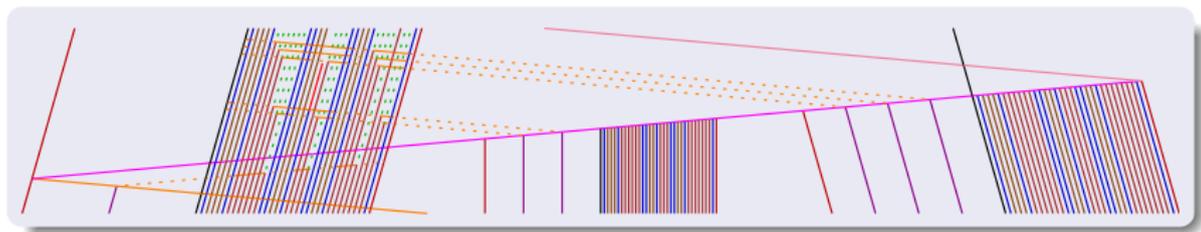
- 1 provide a delay
- 2 test if macro-collision is appropriate and what macro-signals are involved
- 3 if OK
  - start the macro-collision

## Hypotheses for macro-collision

- no other macro-signal nor macro-collision will interfere
- speed of involved macro-signals ranged  $[j, \dots, i]$  (included)
- their main <sup>$\emptyset$</sup>  signals intersect at a unique point

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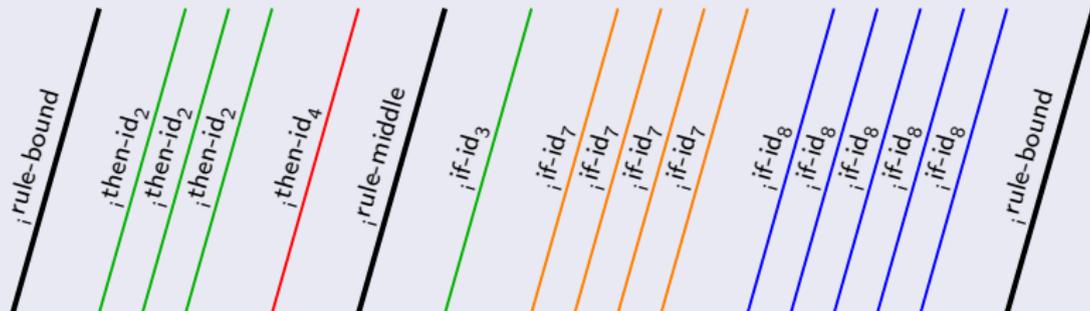
# Removing Unused Tables and Sending ids to Table



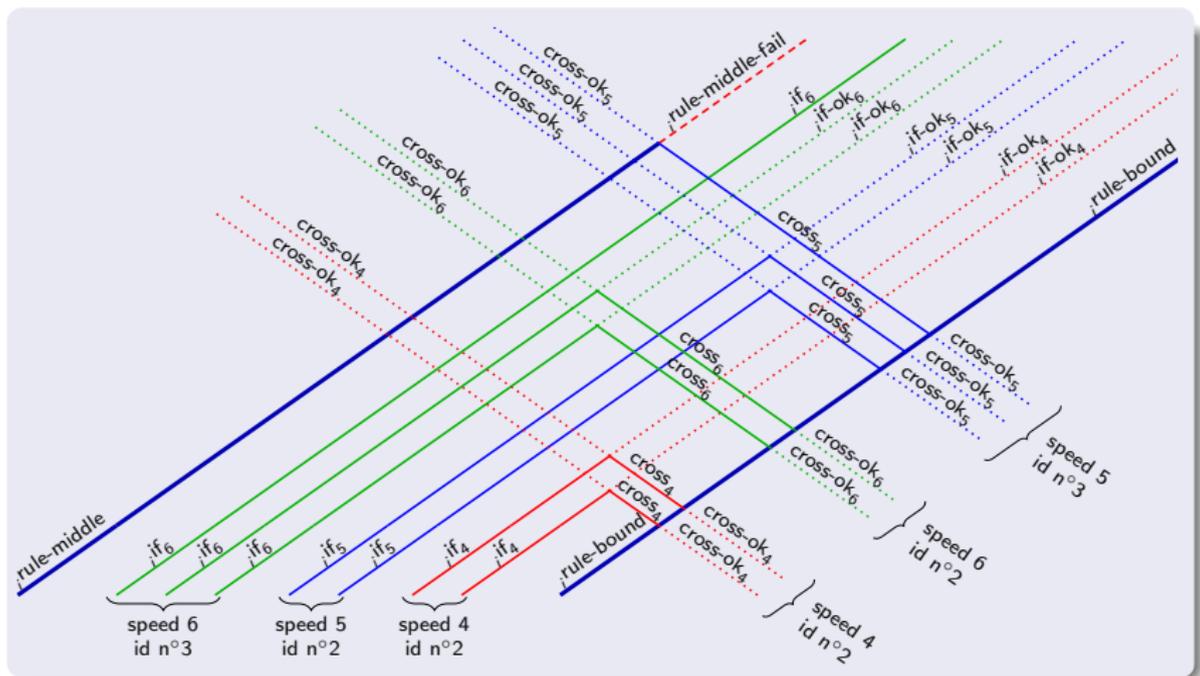
# Collision Rules Encoding

One rule after the other

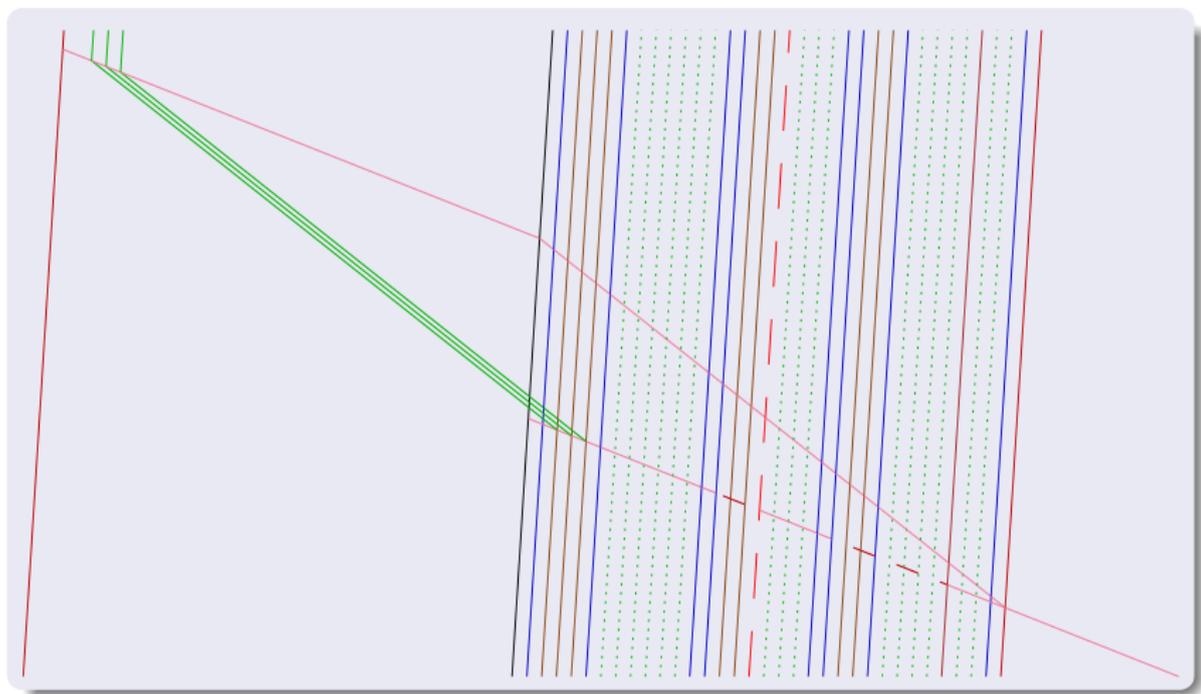
Encoding of  $\{3\mu^1, 7\mu^4, 8\mu^5\} \rightarrow \{2\mu^3, 4\mu^1\}$  in the direction  $i$ .



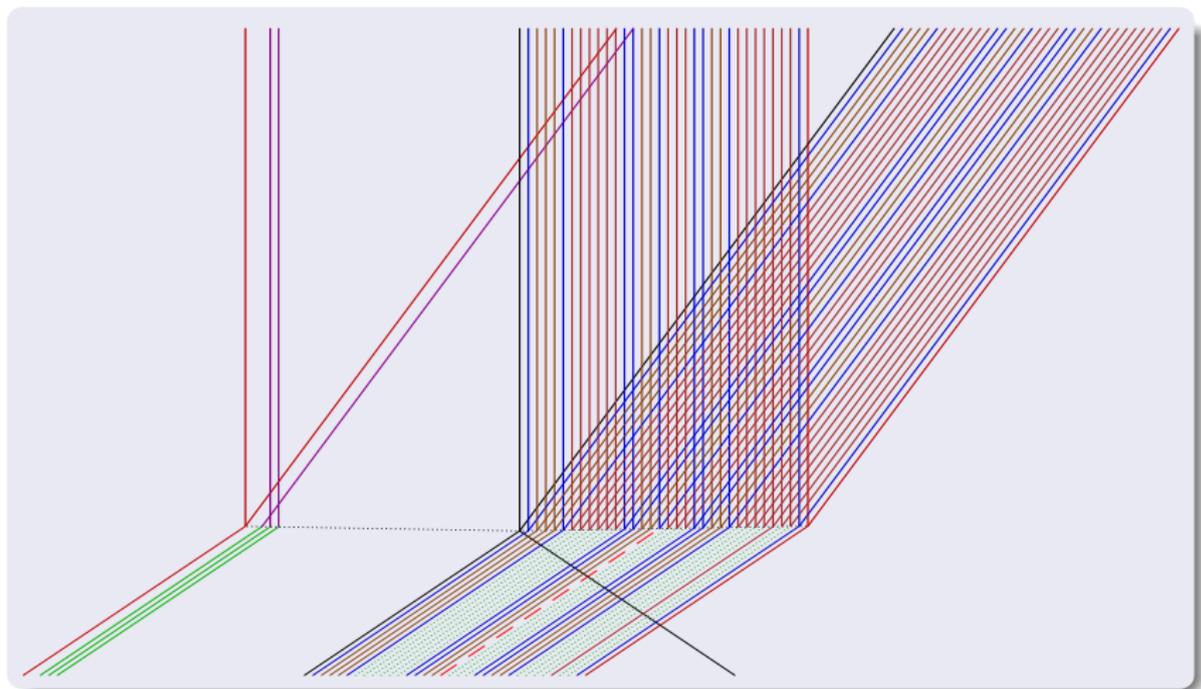
# Comparison of id's in the if-part of a Rule



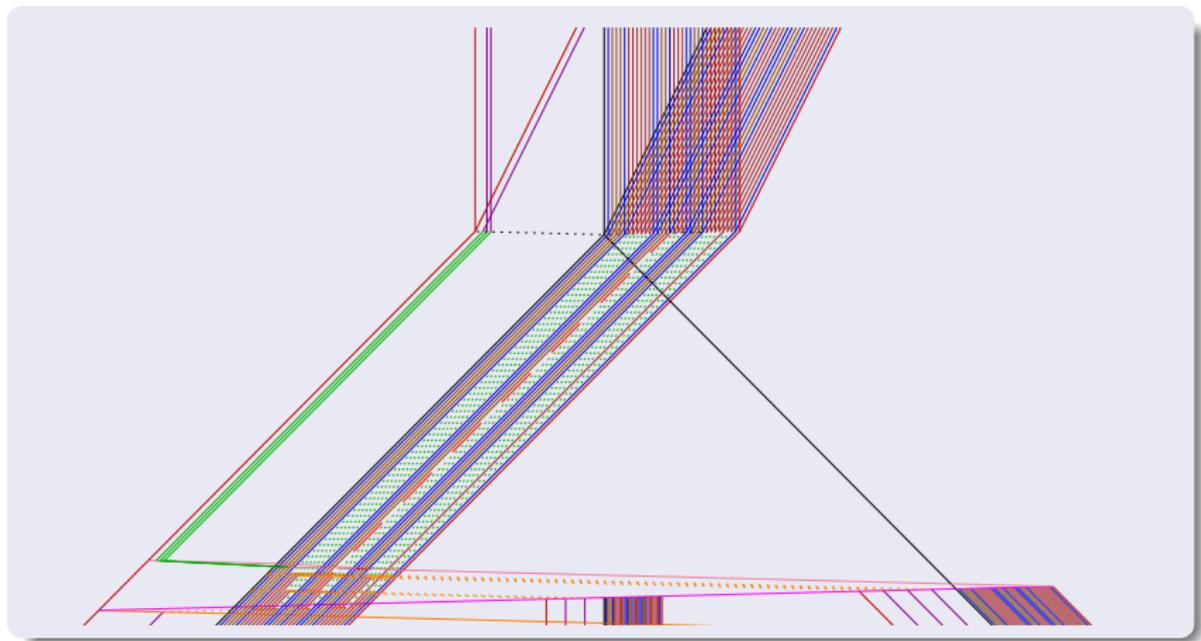
# Rule Selection



# Generating the Output

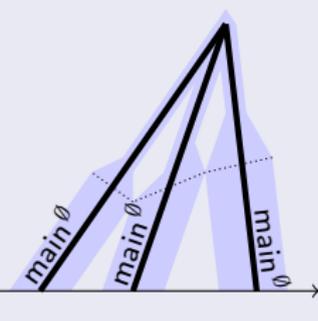
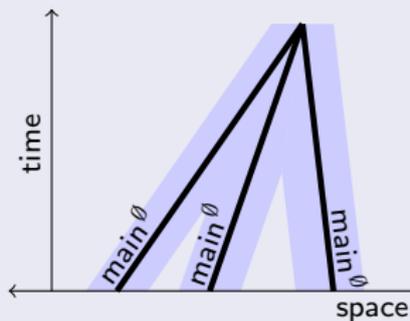
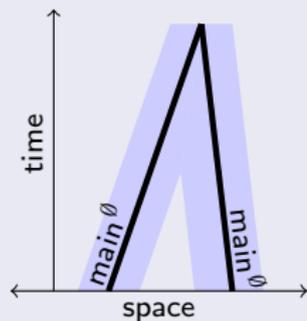


# Whole resolution

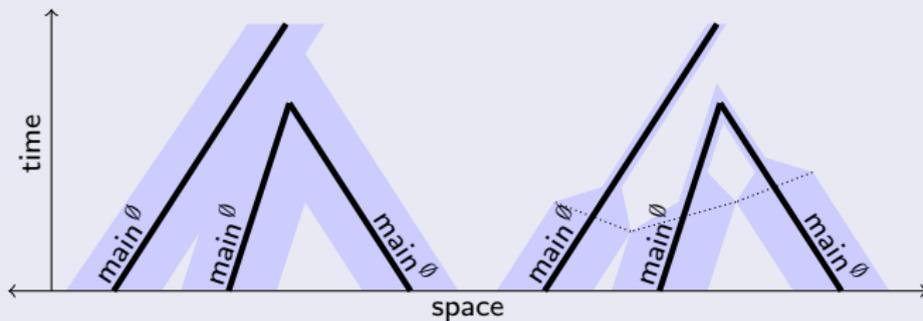


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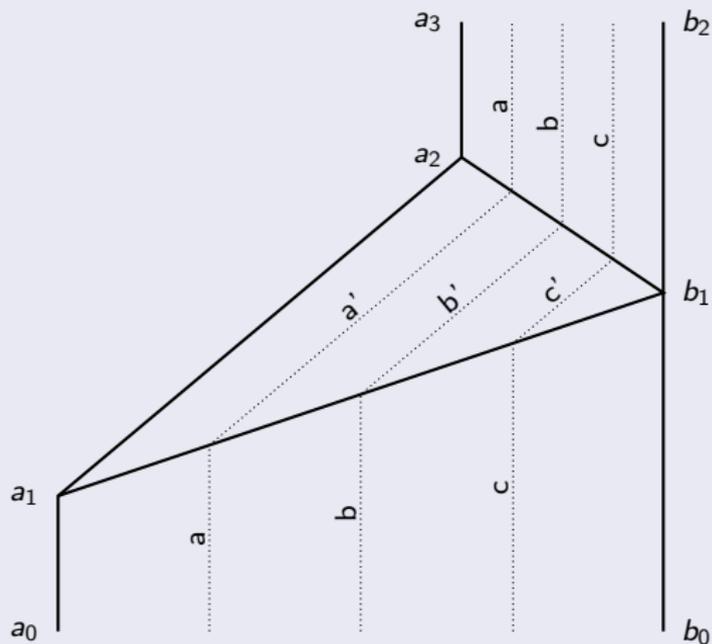
## Good Cases



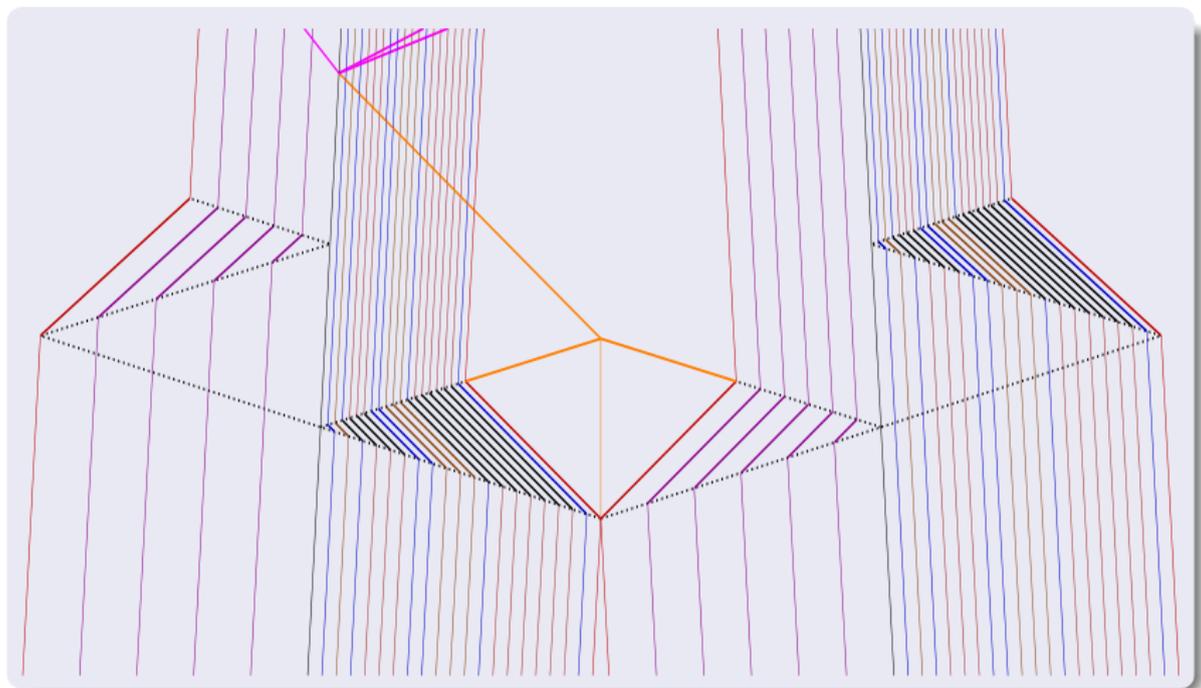
## Bad Case



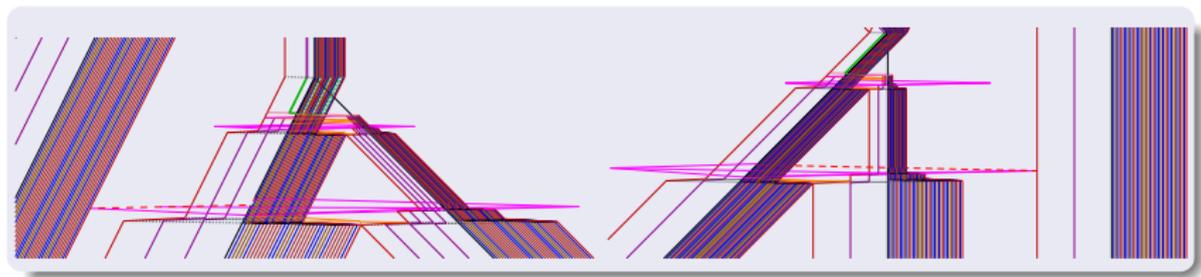
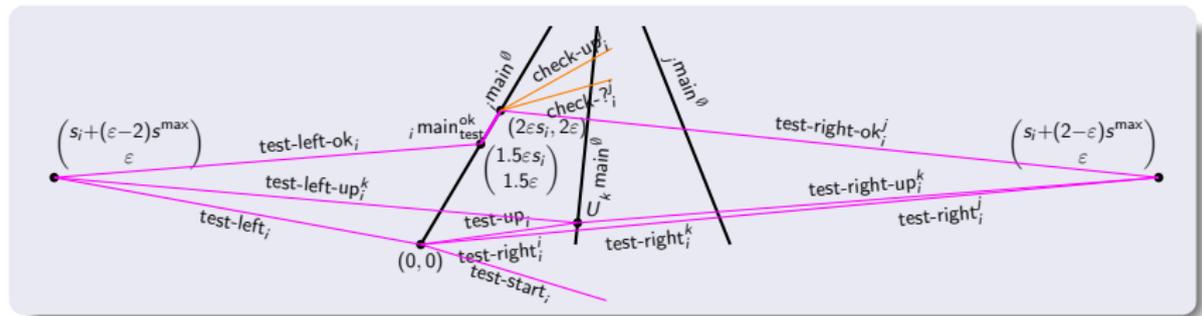
# Shrinking Unit



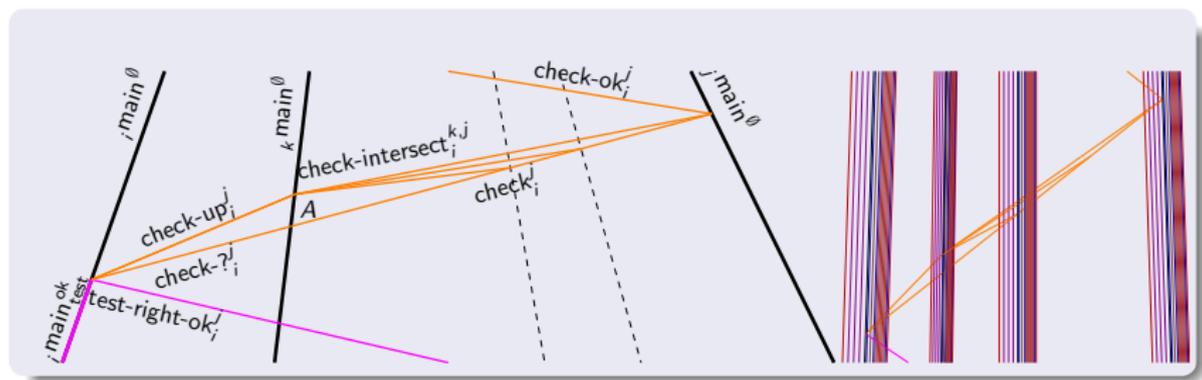
# Shrink



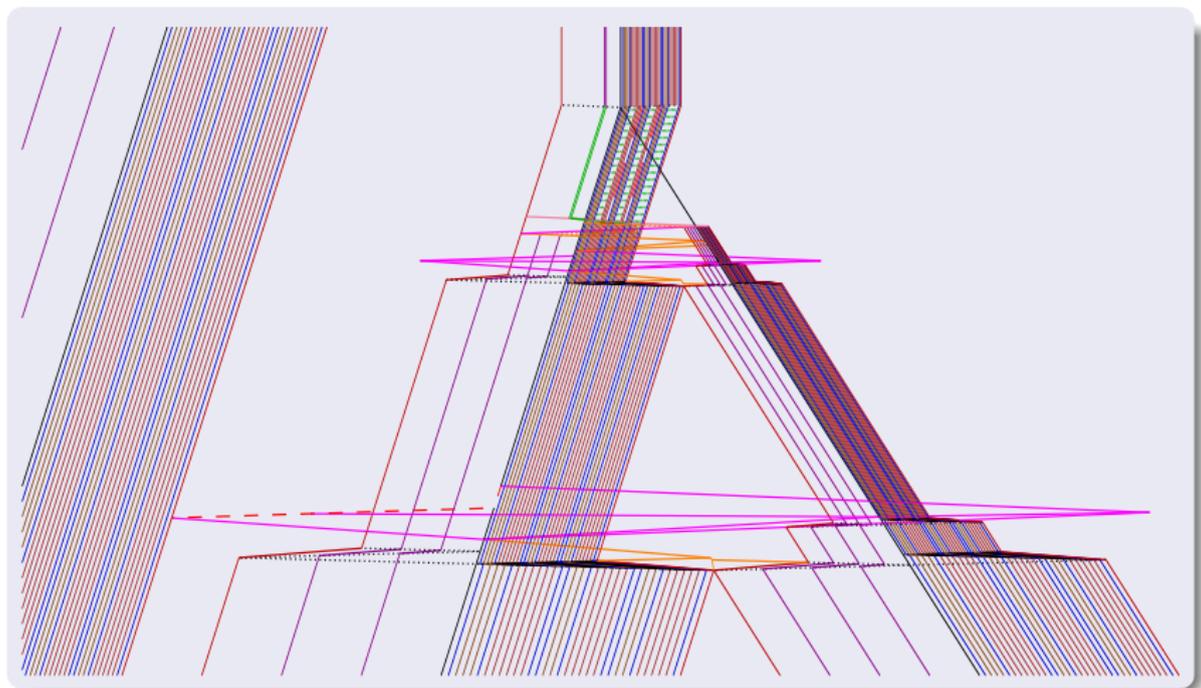
# Testing for Other main<sup>0</sup> Signals



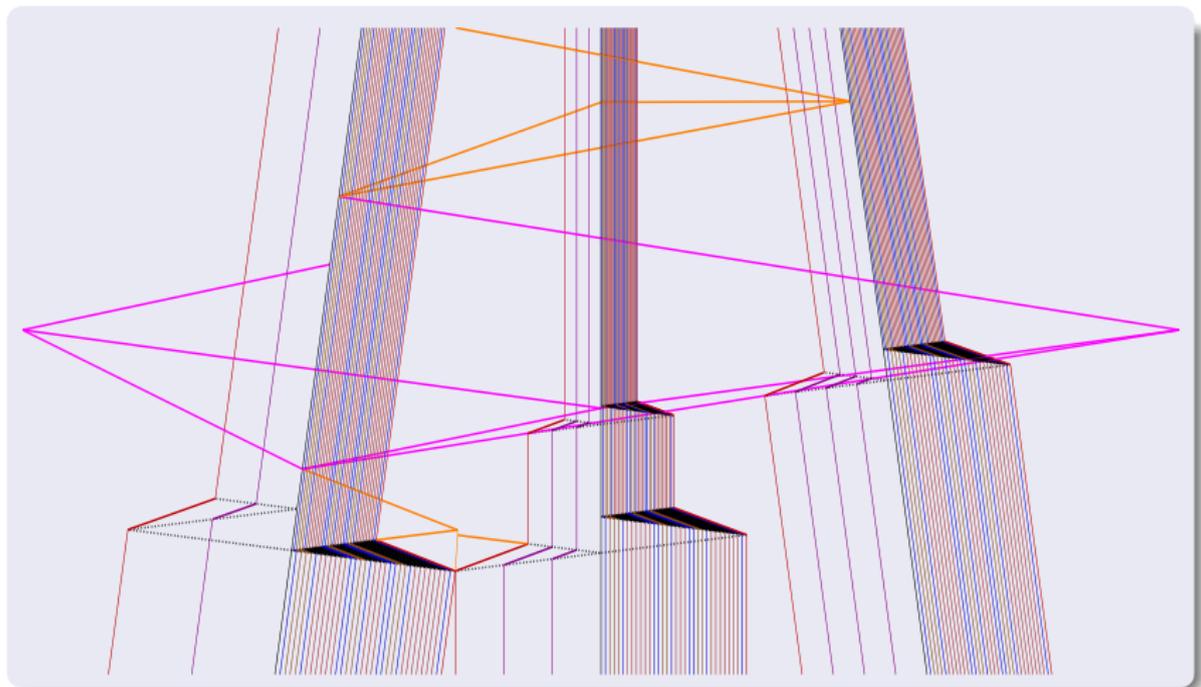
# Checking the right positioning of Other main<sup>0</sup> Signals



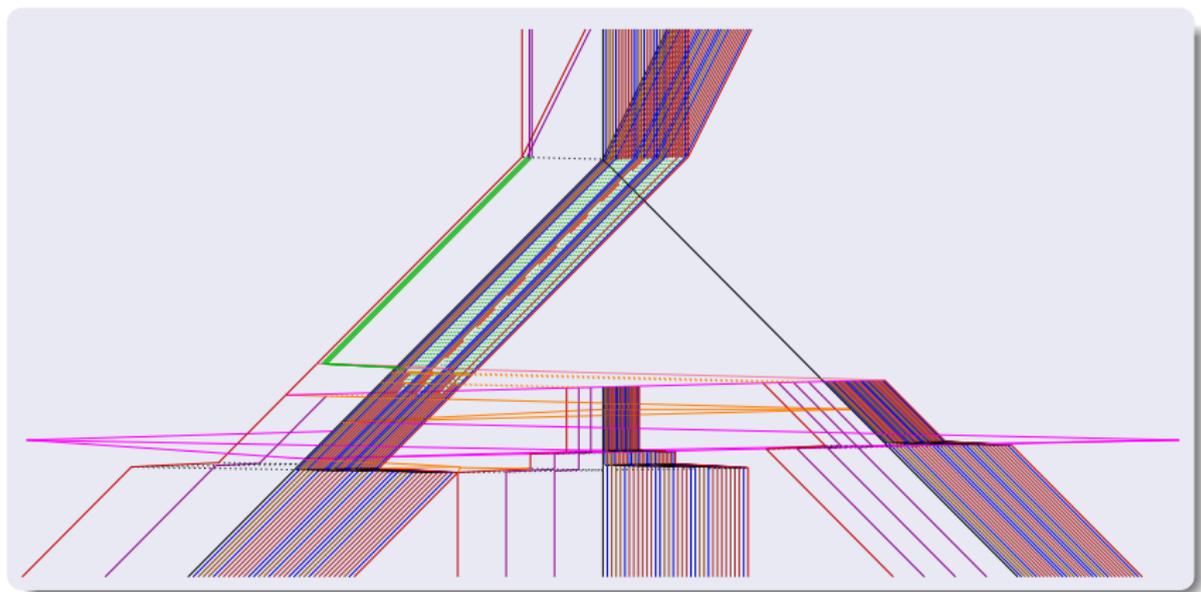
## Detecting Potential Overlaps



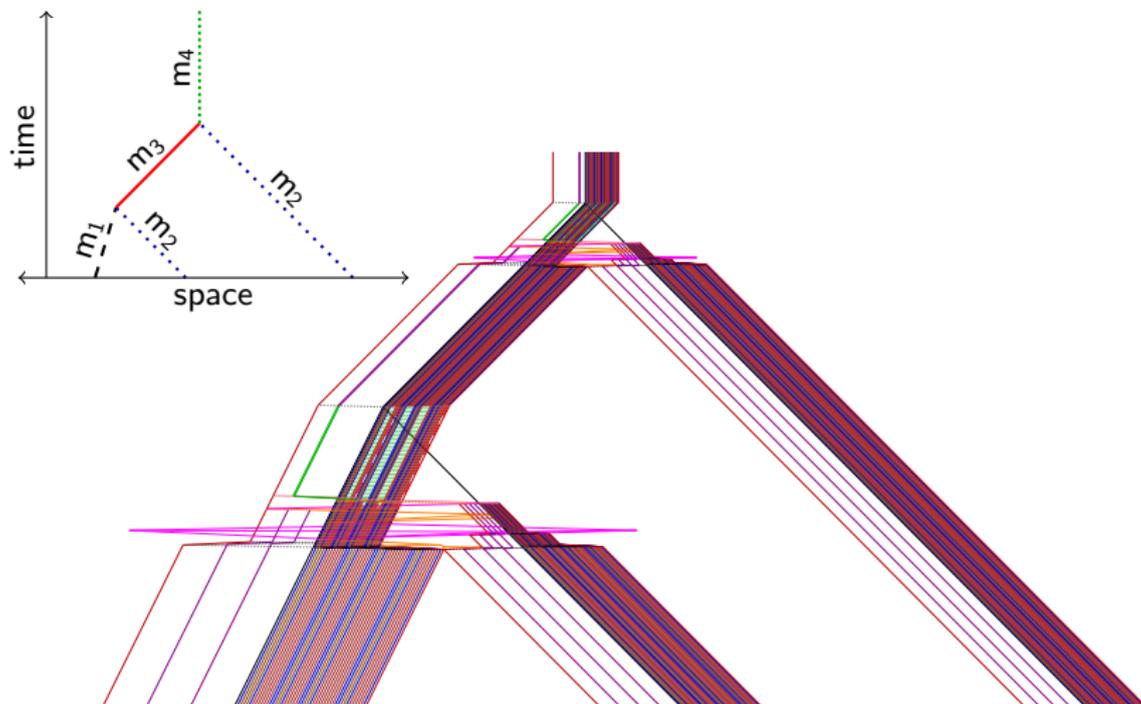
# Whole Preparation



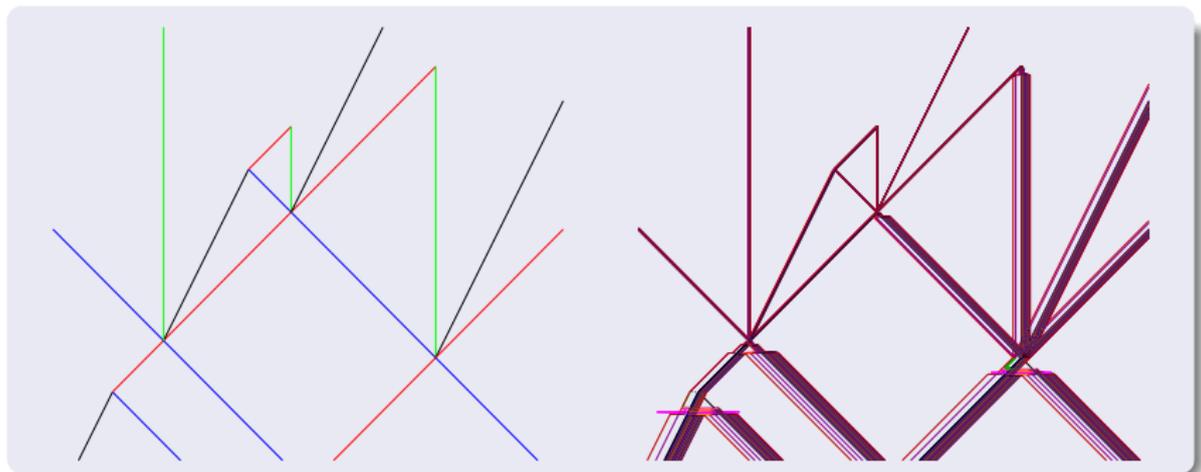
# Exact 3-signal collision



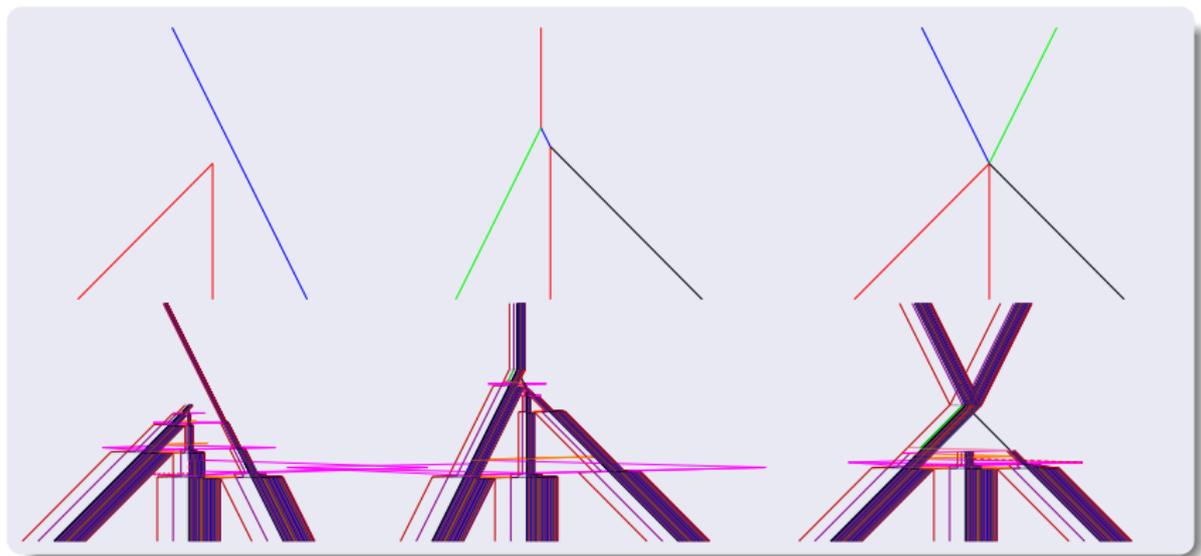
# Some examples



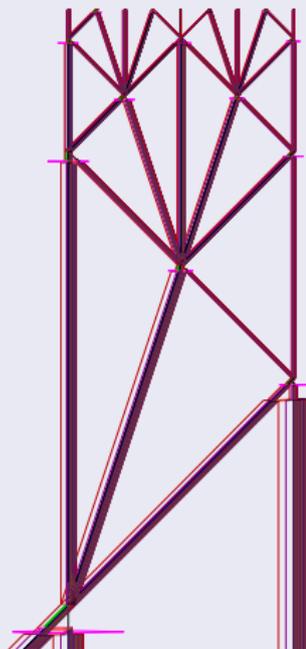
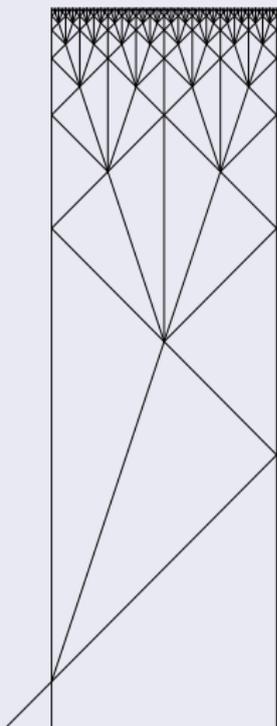
## Some examples



## Some examples



## Some examples



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# Non-determinism in rule output

## Meta-signals

a	0
b	-1
c	1

## Collision rules

$\{ a, b \} \rightarrow \{ a, c \}$   
 $\{ a, b \} \rightarrow \{ b \}$   
 $\{ c, a \} \rightarrow \{ b \}$   
 $\{ c, a \} \rightarrow \{ a \}$

$\{ a, b \} \rightarrow \{ a \}$

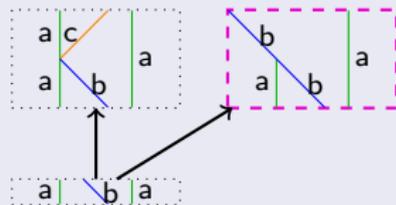
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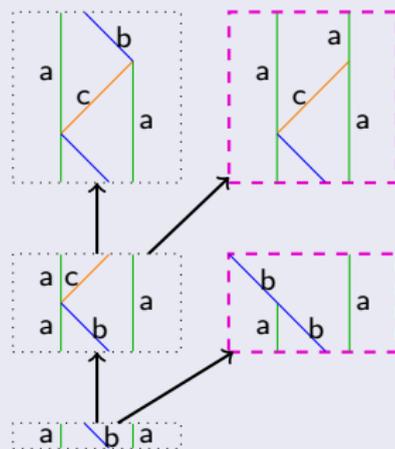
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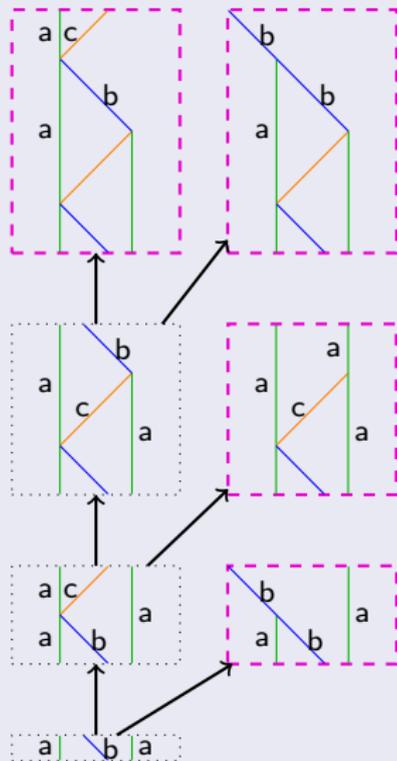
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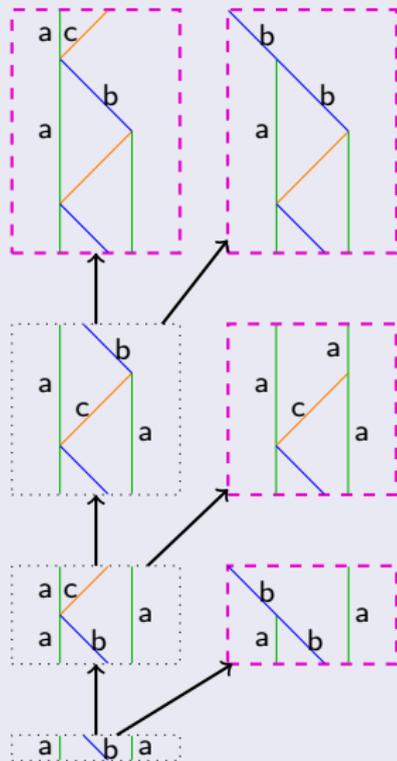
# Non-determinism in rule output

Meta-signals

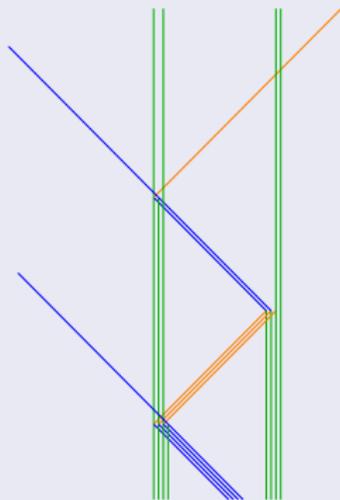
a	0
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c	1

Collision rules

- $\{a, b\} \rightarrow \{a, c\}$
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- $\{c, a\} \rightarrow \{b\}$
- $\{c, a\} \rightarrow \{a\}$



Shifted superposition



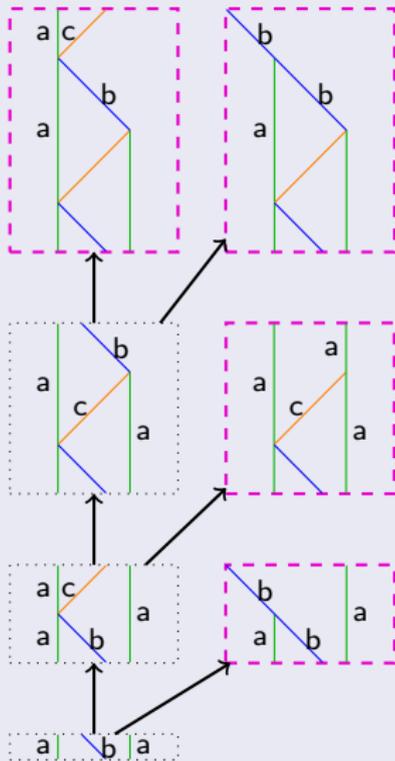
# Non-determinism in rule output

Meta-signals

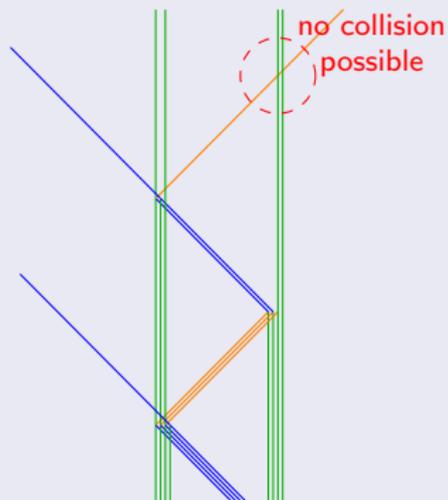
a	0
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Shifted superposition



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## All possible universes

$$\mathcal{U} = \left\{ \left[ \begin{array}{c} \text{diagram 1} \\ \text{diagram 2} \\ \text{diagram 3} \\ \text{diagram 4} \end{array} \right] \right\}$$

## Unbounded signals

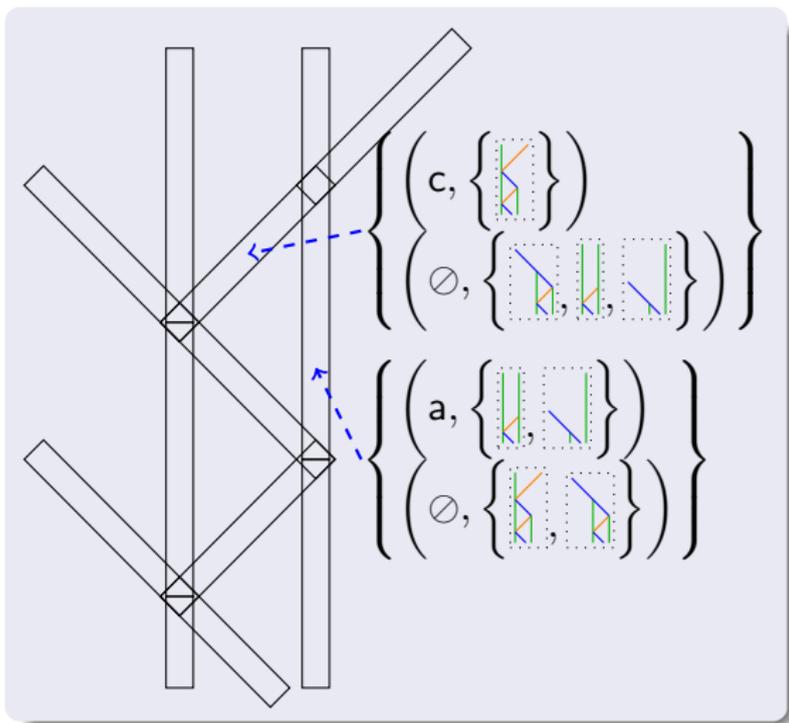
Information held:

$$\{(v_\alpha, u_\alpha)\}_\alpha$$

such that:

$$v_\alpha \in \{\emptyset, a, b, c\}$$

$$\biguplus_\alpha u_\alpha = \mathcal{U}$$



## All possible universes

$$\mathcal{U} = \left\{ \left[ \begin{array}{c} \text{diag 1} \\ \text{diag 2} \\ \text{diag 3} \\ \text{diag 4} \end{array} \right], \left[ \begin{array}{c} \text{diag 5} \\ \text{diag 6} \\ \text{diag 7} \\ \text{diag 8} \end{array} \right], \left[ \begin{array}{c} \text{diag 9} \\ \text{diag 10} \\ \text{diag 11} \\ \text{diag 12} \end{array} \right], \left[ \begin{array}{c} \text{diag 13} \\ \text{diag 14} \\ \text{diag 15} \\ \text{diag 16} \end{array} \right] \right\}$$

## Unbounded signals

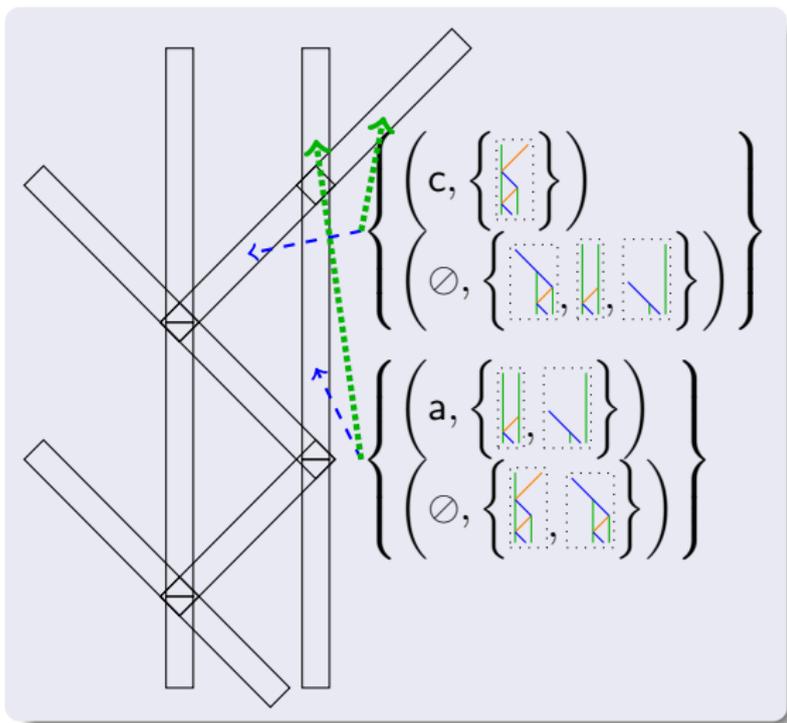
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$$\mathcal{U} = \left\{ \left( \begin{array}{c} \text{[diagram 1]} \\ \text{[diagram 2]} \\ \text{[diagram 3]} \\ \text{[diagram 4]} \end{array} \right) \right\}$$

## Unbounded signals

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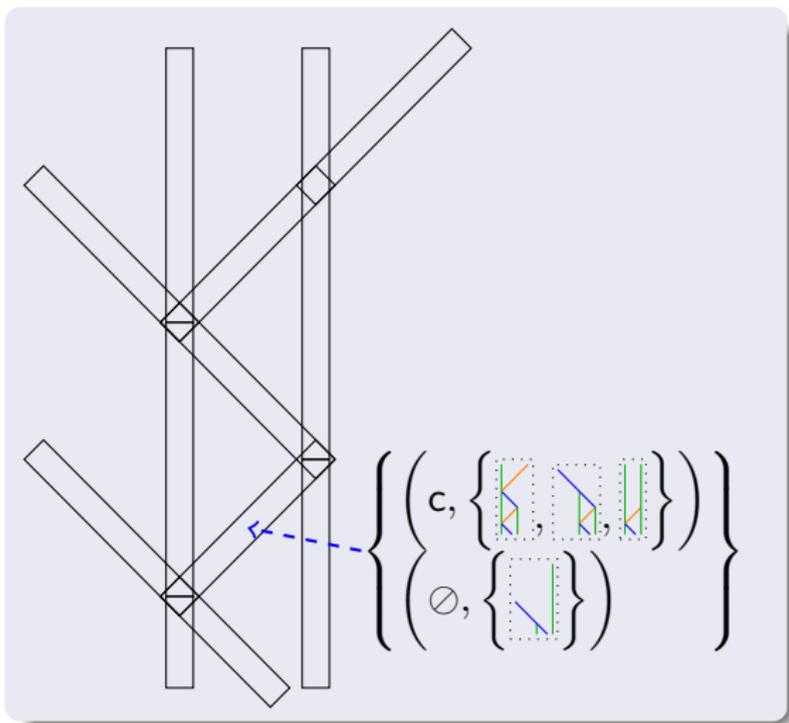
$$\{(v_\alpha, u_\alpha)\}_\alpha$$

such that:

$$v_\alpha \in \{\emptyset, a, b, c\}$$

$$\biguplus_\alpha u_\alpha = \mathcal{U}$$

## Future unknown



## All possible universes

$$\mathcal{U} = \left\{ \left\{ \begin{array}{c} \text{K} \\ \text{L} \end{array} \right\}, \left\{ \begin{array}{c} \text{L} \\ \text{K} \end{array} \right\}, \left\{ \begin{array}{c} \text{K} \\ \text{L} \end{array} \right\}, \left\{ \begin{array}{c} \text{L} \\ \text{K} \end{array} \right\} \right\}$$

## Unbounded signals

Information held:

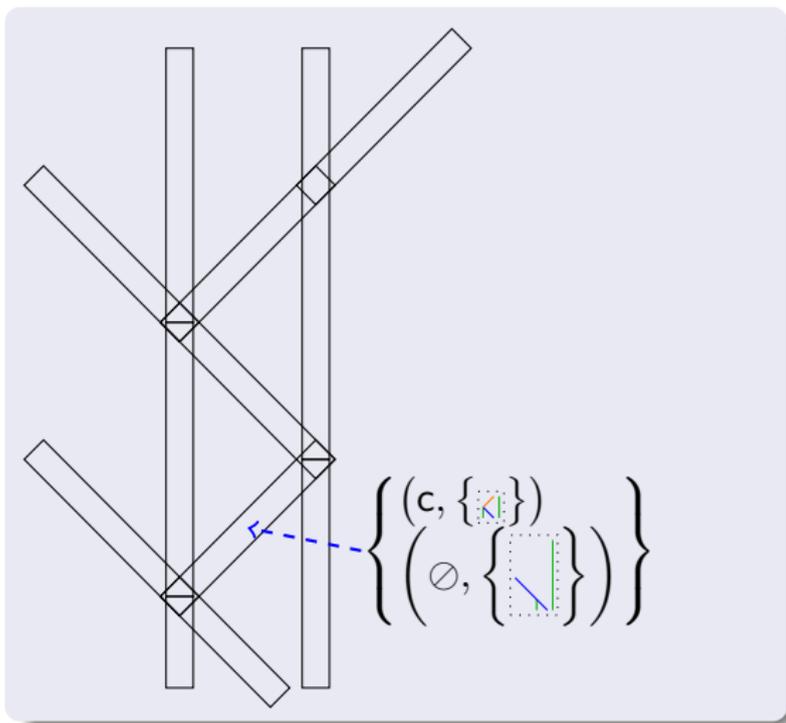
$$\{(v_\alpha, u_\alpha)\}_\alpha$$

such that:

$$v_\alpha \in \{\emptyset, a, b, c\}$$

$$\biguplus_\alpha u_\alpha = \mathcal{U}$$

## Future unknown



## All possible universes

$$\mathcal{U} = \left\{ \begin{array}{c} \text{[Diagram 1]} \\ \text{[Diagram 2]} \\ \text{[Diagram 3]} \\ \text{[Diagram 4]} \end{array} \right\}$$

## Unbounded signals

Information held:

$$\{(v_\alpha, u_\alpha)\}_\alpha$$

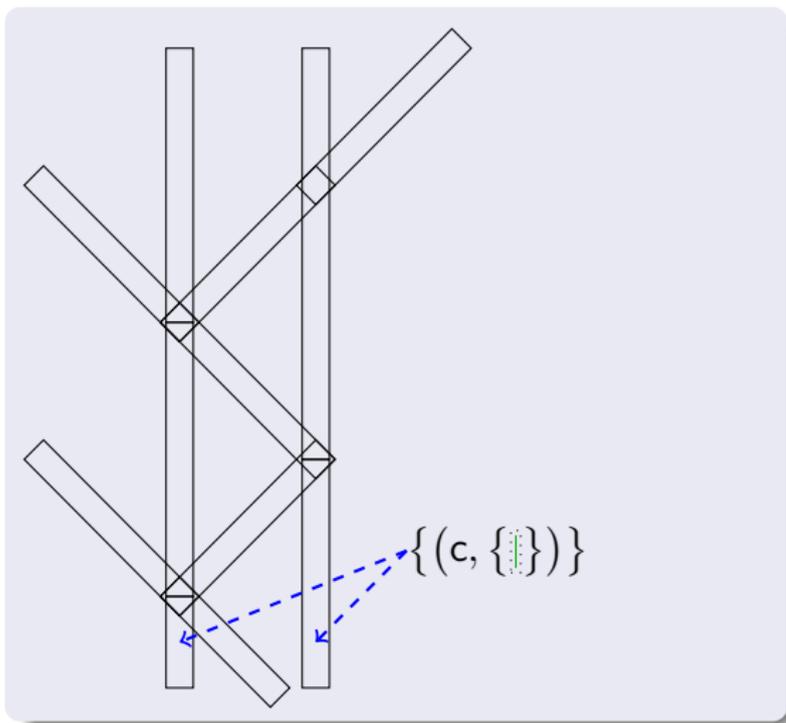
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$$\bigsqcup_\alpha u_\alpha = \mathcal{U}$$

Future unknown

Distinguish



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## Unbounded signals

Information held:

$$\{(v_\alpha, u_\alpha)\}_\alpha$$

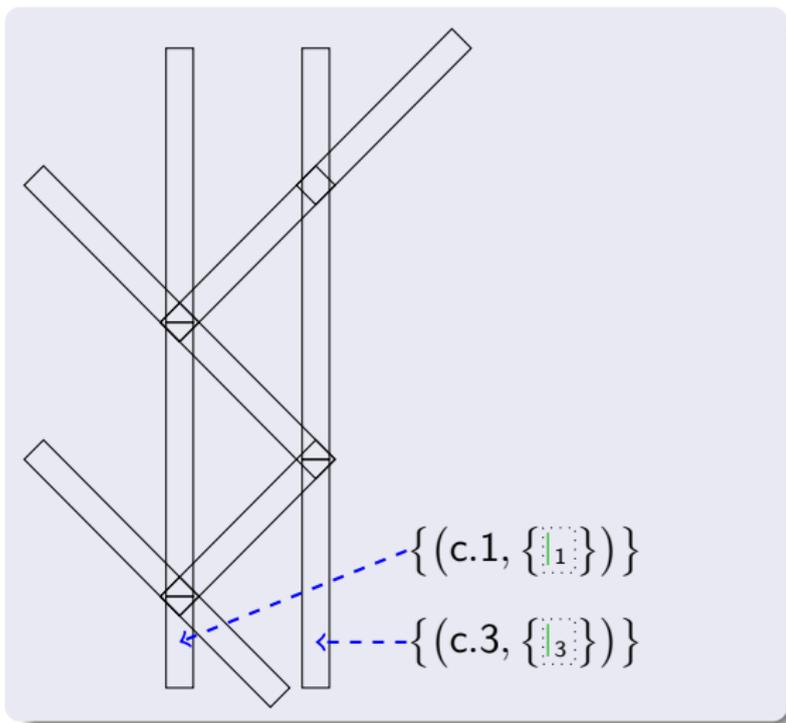
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Future unknown

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## Unbounded signals

Information held:

$$\{(v_\alpha, u_\alpha)\}_\alpha$$

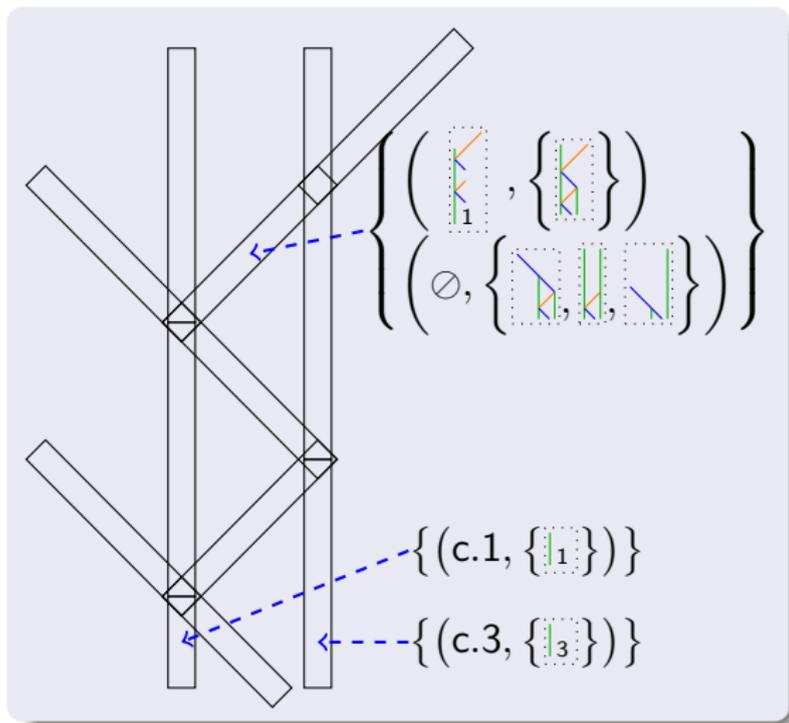
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Future unknown

Distinguish



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# Macro-collision

$\{ (v_\alpha, u_\alpha) \}_\alpha$  meets  $\{ (v'_\beta, u'_\beta) \}_\beta$

$$\textcircled{1} \quad E_1 = \left\{ \left( (v_\alpha, v'_\beta), u_\alpha \wedge u'_\beta \right) \right\}_{\alpha, \beta} \quad \mathcal{U} = \bigsqcup_{\alpha, \beta} u_\alpha \wedge u'_\beta$$

$$\textcircled{2} \quad E_2 = \{ ((v, v'), u \wedge u') \in E_1 \mid u \wedge u' \neq \emptyset \}$$

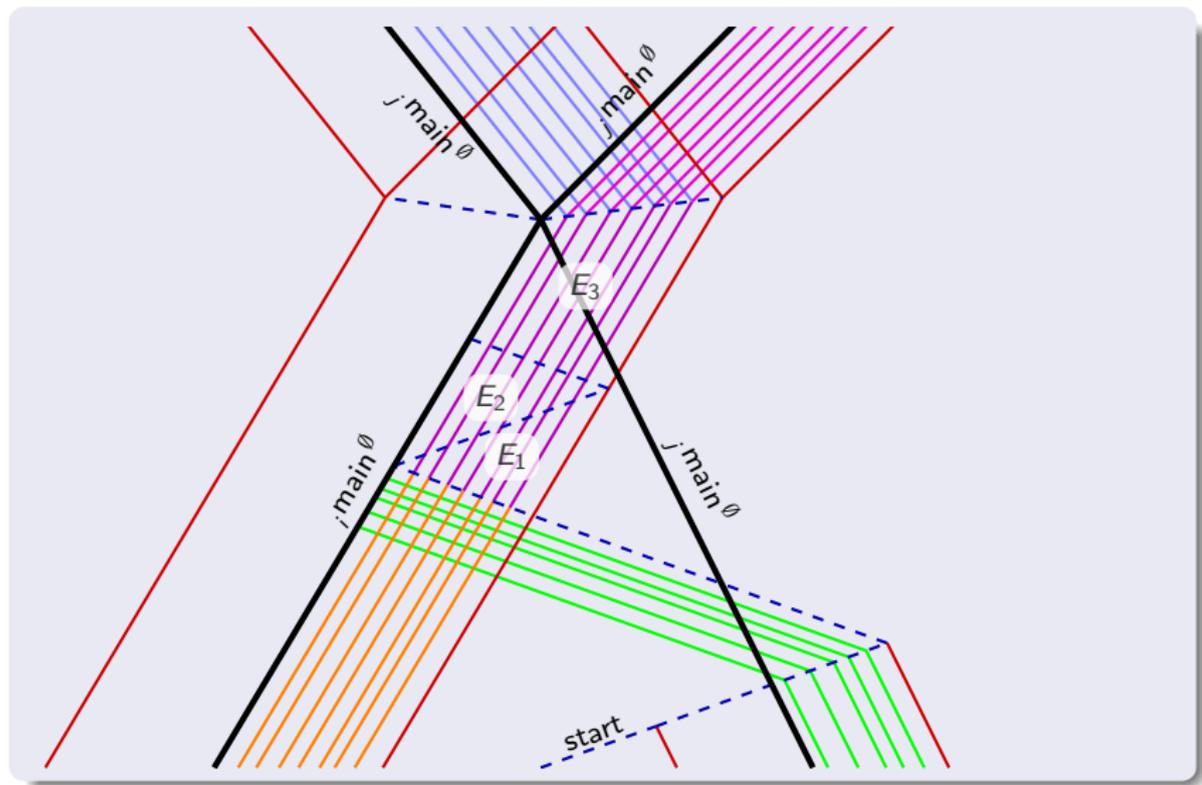
$$\textcircled{3} \quad E_3 = \{ (\emptyset, w) \mid ((\emptyset, \emptyset), w) \in E_2 \} \\ \cup \{ (\{\mu\}, w) \mid ((\mu, \emptyset), w) \in E_2 \vee ((\emptyset, \mu), w) \in E_2 \} \\ \cup \{ (\rho^+.\mu, \rho(\mu, \nu) \wedge w) \mid ((\mu, \nu), w) \in E_2 \wedge \rho^- = \{\mu, \nu\} \}$$

$$\textcircled{4} \quad \text{Out\_Speed} = \{ \text{Speed}(\mu) \mid \exists \mu, (F, w) \in E_3, \mu \in F \}$$

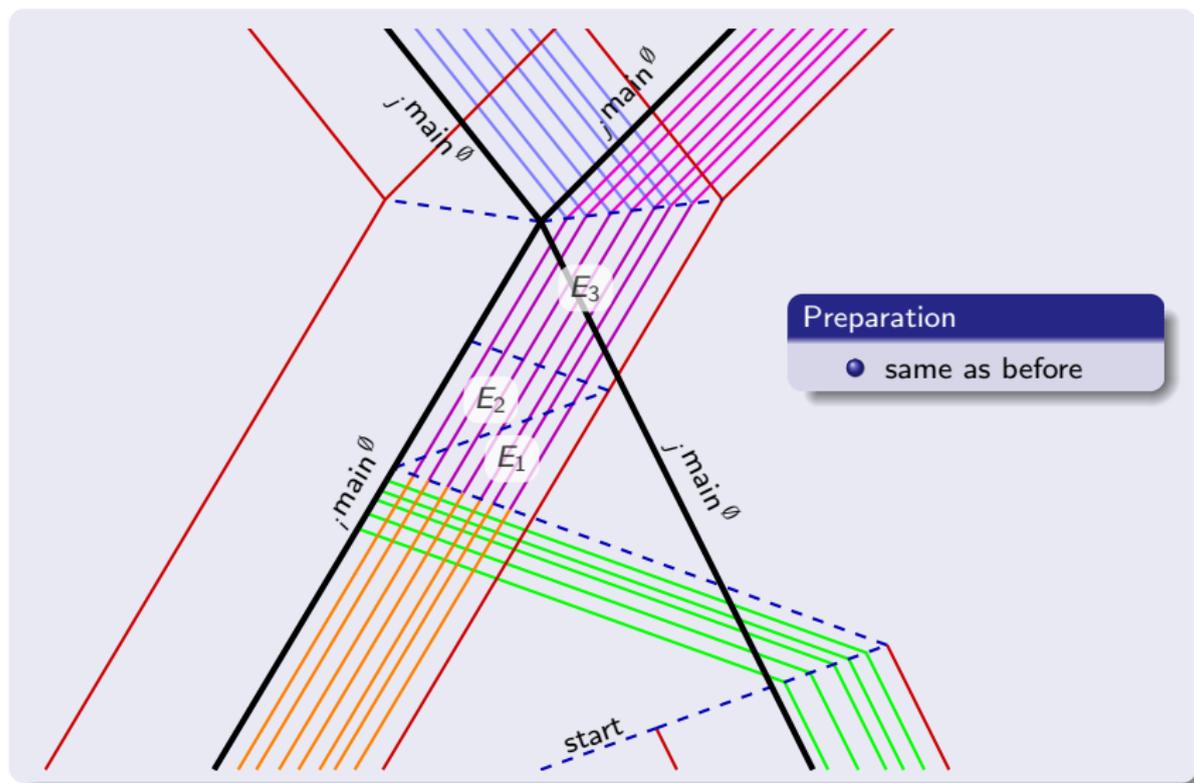
$$\textcircled{5} \quad \forall s \in \text{out}, \\ \text{Out}_s = \{ (\mu, w) \mid \exists F, (F, w) \in E_3 \wedge \mu \in F, \text{Speed}(\mu) = s \} \\ \cup \{ (\emptyset, w) \mid \exists F, (F, w) \in E_3 \wedge \forall \mu \in F, \text{Speed}(\mu) \neq s \}$$

- Compatible string encodings

# Displaying the operations



# Displaying the operations



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- Very rich setting
- Intrinsically universal family of signal machines
- Non-deterministic signal machine are not “more powerful”

- Very rich setting
- Intrinsically universal family of signal machines
- What is the complexity?
- Non-deterministic signal machine are not “more powerful”
- How to extract “result”?
- What is the complexity?

- Very rich setting
- Intrinsically universal family of signal machines
- What is the complexity?
- Non-deterministic signal machine are not “more powerful”
- How to extract “result”?
- What is the complexity?
- Augmented signal machines

That's all folks!

Thank you for your attention

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