

Abstract geometrical computation: small Turing universal signal machines

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IW Complexity of Simple Programs

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- 1 Introduction
- 2 Turing machines
- 3 Cellular automata
- 4 Cyclic tag systems
- 5 Conclusion

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Context

Collision based computing

Idealization

- continuous space
- continuous time
- dimensionless particles/signals

Abstract geometrical computation

Signal machines

- meta-signals (finitely many)
- their speed/velocity
- collision rules

Signals, e.g.

- red (with speed 1) at position xx
- blue (with speed -1) at position yy

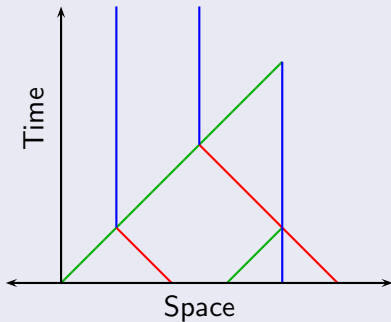
Collision, e.g.

- rule $\{\text{green, red}\} \rightarrow \{\text{blue}\}$

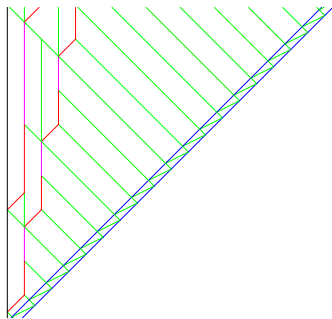
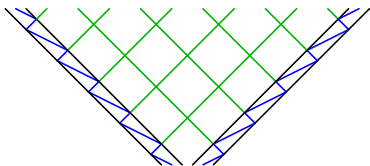
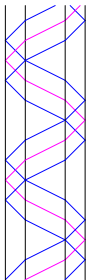
- application



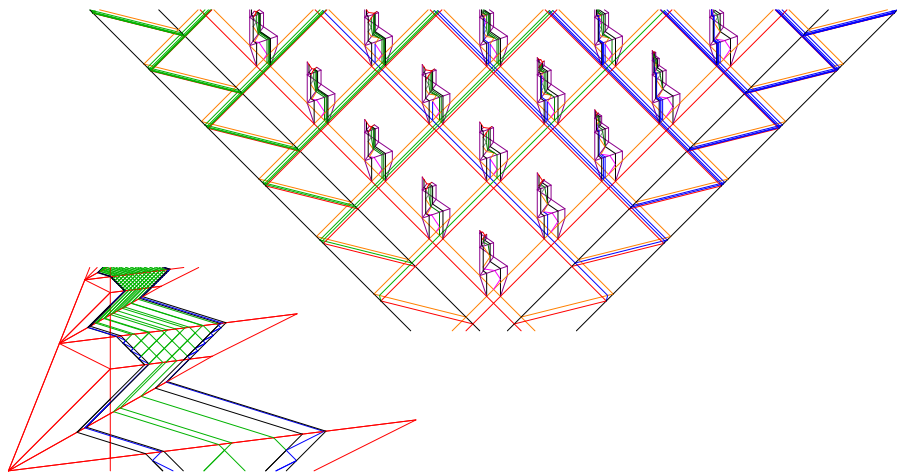
An example



More examples

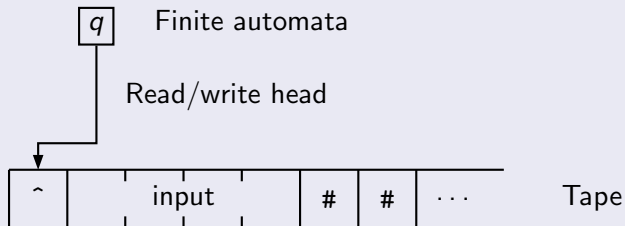


More complex examples



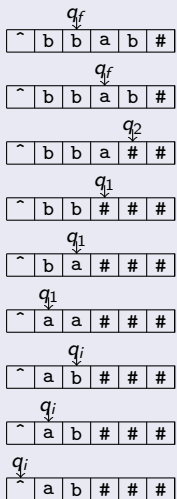
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Turing machines?

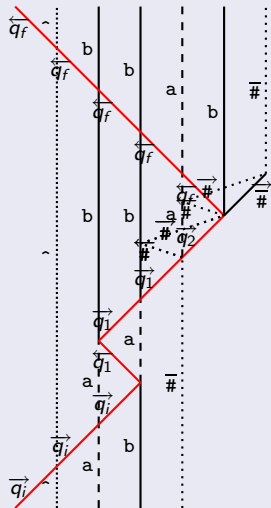


Simulation

Iterations of a Turing machine



Corresponding signal machine



How many meta-signals?

		Meta-signal
1 symbol	\rightsquigarrow	1
1 state	\rightsquigarrow	2
<i>for finitness</i>		
$\overline{\#}, \overleftarrow{\#}, \overrightarrow{\#}, \overleftrightarrow{\#}$	\rightsquigarrow	4

$$|\Gamma| + 2|Q| + 4$$

Results

universal
18 (Woods and Neary, 2007)

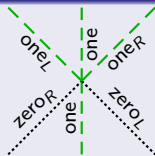
semi-universal
7 (Smith, 2007)

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Cellular automata

Rule 110 and one transition implementation

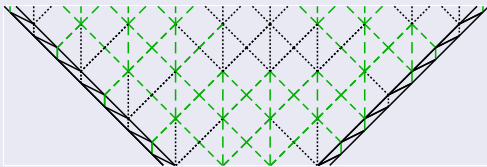
Output	0	1	1	0	1	1	1	0
Input	111	110	101	100	011	010	001	000



Evolution and simulation on 11 framed by $\omega(10)$ and $(011)^\omega$

```

0111001100
1011000111
 10111110
   101101
  
```



How many meta-signals?

		Meta-signal
1 state	\rightsquigarrow	3
<i>for finitness</i>		
Regular pattern on both side	\rightsquigarrow	expensive

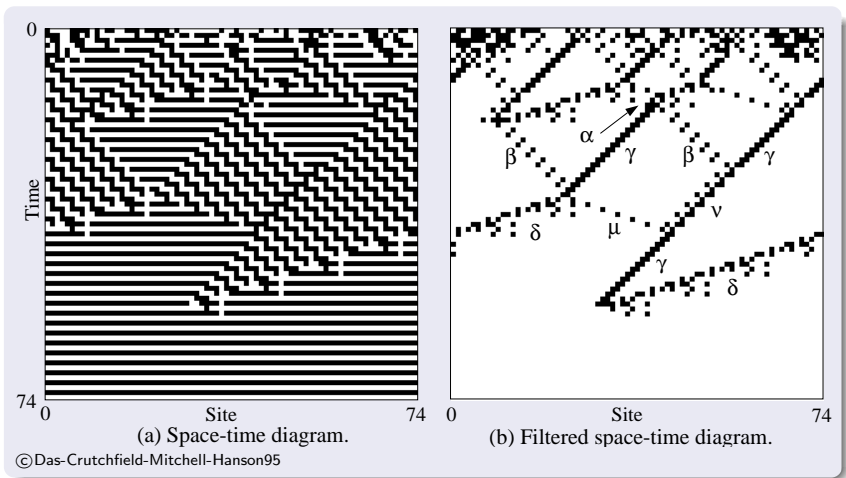
$$3|Q|+???$$

Results

universal
not interesting

semi-universal
6 (Cook, 2004)

Link CA-ACG



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Cyclic tag system?

Definition

- a binary word
- a circular list

Dynamics

101

011 :: h :: 0110 :: 01011

Halt

- empty word
- halt appendant (here h)
- cycle (too expensive to test)

Cyclic tag system?

Definition

- a binary word
- a circular list

Dynamics

~~h~~01011

~~011~~ :: h :: 0110 :: 01011 :: 011

Halt

- empty word
- halt appendant (here h)
- cycle (too expensive to test)

Cyclic tag system?

Definition

- a binary word
- a circular list

Dynamics

~~1~~01011

~~0~~11 :: ~~1~~ :: 0110 :: 01011 :: 011 :: h

Halt

- empty word
- halt appendant (here h)
- cycle (too expensive to test)

Cyclic tag system?

Definition

- a binary word
- a circular list

Dynamics

~~101~~011**0110**

011 :: ~~h~~ :: ~~0110~~ :: 01011 :: 011 :: h :: 0110

Halt

- empty word
- halt appendant (here h)
- cycle (too expensive to test)

Cyclic tag system?

Definition

- a binary word
- a circular list

Dynamics

~~1~~010110110

011 :: h :: 0110 :: 01011 :: 011 :: h :: 0110 :: 01011

Halt

- empty word
- halt appendant (here h)
- cycle (too expensive to test)

Cyclic tag system?

Definition

- a binary word
- a circular list

Dynamics

~~10101101100~~**11**

011 :: h :: 0110 :: 01011 :: 011 :: h :: 0110 :: 01011 :: 011

Halt

- empty word
- halt appendant (here h)
- cycle (too expensive to test)

Cyclic tag system?

Definition

- a binary word
- a circular list

Dynamics

~~101011~~**0110011**

011 :: h :: 0110 :: 01011 :: 011 :: h :: 0110 :: 01011:: 011 :: h

Halt

- empty word
- halt appendant (here h)
- cycle (too expensive to test)

How many meta-signals?

Universality

- 13 meta-signals
- 21 non-blank rules

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Conclusion

Bounds

universal

13

semi-universal

6

Future work

- lower this bound
- bounds for
 - accumulation (3 or 4)
 - black hole implementation (13 ? 14 ? more)
 - analog computation (BSS or computable analysis)